A0 Submission Report

Ву

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Q1

We have profiled the programs

- 1. Bubble Sort
- 2. Selection Sort
- 3. Merge Sort
- 4. Quick Sort

belonging to Sorting Family for:

<u>Note:</u> Instrumentation code written (p1-p4.cpp) and Trace obtained for each sorting program are present in folder "q1"

a) Instruction count, Instruction Address Trace, Memory Reference Trace

```
dibyadarshan@hota:~/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ ../.../.pin -t obj-intel64/p1.so -- ~/Desktop/HPC/hpc-assignments/a0/q1/s
.out
Selection sort, sorted array:
0.500000 1.000000 2.300000 2.500000 6.400000 25.5000000
dibyadarshan@hota:~/Desktop/pin-3.11-97999-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ head p1.out
Image /home/dibyadarshan/Desktop/HPC/hpc-assignments/a0/q1/a.out has 271 instructions
Image /lib64/ld-linux-x86-64.so.2 has 29322 instructions
Image /lib64/ld-linux-x86-64.so.2 has 29322 instructions
Image /lib64/ld-linux-x86-64.so.2 has 29322 instructions
Image source /linux/source/tools/manualExamples$
Image source /linux/source/tools/manualExamples$

=> 0x00007f6b38ee9003
Write 0x00007ff6b38ee9003
Write 0x00007ff6b38ee9ea0
Write 0x00007ff6b38ee9ea0

=> 0x00007f6b38ee9ea1

=> 0x00007f6b38ee9ea4
dibyadarshan@hota:-/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$
```

```
dibyadarshan@hota:-/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ ../../.pin -t obj-intel64/p1.so -- ~/Desktop/HPC/hpc-assignments/a0/q1/a.out
Werge sort, sorted array:
0.500000 1.000000 2.300000 2.500000 6.400000 25.500000 25.600000
dibyadarshan@hota:-/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ head p1.out
Inage /home/dibyadarshan/Desktop/HPC/hpc-assignments/a0/q1/a.out has 443 instructions
Inage /lib64/ld-linux-x86-64.so.2 has 29322 instructions
Inage /lib64/ld-linux-x86-64.so.2 has 29322 instructions
==> 0x00007fc25409e090
```

```
dibyadarshan@hota:~/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ ../.././pin -t obj-intel64/p1.so -- ~/Desktop/HPC/hpc-assignments/a0/q1/a .out
Outck sort, sorted array:
0.500000 1.000000 2.300000 2.300000 2.5000000 6.400000 25.5000000
dibyadarshan@hota:~/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ head p1.out
Image /home/dibyadarshan/Desktop/HPC/hpc-assignments/a0/q1/a.out has 317 instructions
Image /loft64/di-linux-x86-64.so.2 has 29322 instructions
Image [vdso] has 426 instructions
==> 0x0000779e05730090
Write 0x0000779e05730090
Write 0x0000779e05730090
Write 0x0000779e05730000
Write 0x0000779e05730000
Write 0x0000779e05730000
Hrite 0x0000779e057300001
==> 0x0000779e057300001
==> 0x0000779e057300001
==> 0x0000779e057300001
==> 0x0000779e057300001
```

Analysis:

We have obtained the instruction count using the static image of the object file, thus counting only the total instruction present, for the 4 programs as shown in the screenshots above. The instruction address along with memory reference trace (for reads and writes to a memory address belonging to the instruction logged before it) is obtained placing instrumentation code before each executed instruction.

b) Instruction mix with the total number of dynamic instructions, integer, floating-point, load, store, branch.

Analysis:

We obtained the dynamic instruction count by placing the instrumentation code before every instruction executed. For integer and floating operation we compared the instruction category of the instruction with appropriate ones. Load, stores are detected using for each instruction using Pin's inspection API's. Branch instructions also can be determined by using one of the API.

c) Total branches that are taken and total forward branches that are taken

```
dibyadarshan@hota:~/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ ../.../.pin -t obj-intel64/p3.so -- ~/Desktop/HPC/hpc-assignments/a0/q1/a.out
Selection sort, sorted array:
6.500000 1.000000 2.300000 2.500000 6.400000 25.5000000
dibyadarshan@hota:~/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ head p3.out
Total Taken Branches: 20382
Total Taken Forward Branches: 9039
dibyadarshan@hota:~/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ |
```

```
dibyadarshan@hota:-/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ ../../.pin -t obj-intel64/p3.so -- ~/Desktop/HPC/hpc-assignments/a0/q1/a
.out
Bubble sort, sorted array:
0.500000 1.000000 2.300000 2.500000 6.400000 25.500000 25.500000 dibyadarshan@hota:-/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ head p3.out
Total Taken Branches: 20374
Total Taken Franches: p034
dibyadarshan@hota:-/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ ../../.pin -t obj-intel64/p3.so -- ~/Desktop/HPC/hpc-assignments/a0/q1/a
.out
Merge sort, sorted array:
0.500000 1.000000 2.300000 2.500000 6.400000 25.500000 25.600000
dibyadarshan@hota:-/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ head p3.out
Total Taken Branches: 20429
Total Taken Branches: 20429
Total Taken Branches: 9077
dibyadarshan@hota:-/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ |

dibyadarshan@hota:-/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ |

dibyadarshan@hota:-/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ |

dibyadarshan@hota:-/Desktop/pin-3.11-97998-g7ecce2dac-gcc-linux/source/tools/ManualExamples$ head p3.out
Total Taken Branches: 20372
```

Analysis:

To detect taken branches we placed the instrumentation function to be present when a branch is taken in the executing program. For identifying forward branches we compare the current branch instruction address to the address taken by the branch.

d) Read-After-Write(RAW), Write-After-Write(WAW) and Write-After-Read(WAR) distribution in the dynamic instruction stream.

```
dibyadarshan@hota:-/Desktop/pin-3.11-97998-gPecce2dac-gcc-linux/source/tools/ManualExamples$ ../../.pin -t obj-intel64/p4.so -- ~/Desktop/HPC/hpc-assignments/a0/q1/a out sources of the second of the second of the second out sources of the second out second out sources of the second out sources of the second out sources out second out second out sources o
```

Analysis:

To help keep track of all read and writes to a memory address, we keep a map indexed by memory address with each entry storing count of instruction in which the memory address was read and written to. For any address encountered next, if this address is present in the map we identify its total read and write dependencies using the map and increment the respective dependency count. Basically, the map acts as a DP approach to get dependencies rather than iterating through the instruction trace every time.

Q2

We wrote a program to find the highest eigenvalue of an NxN real symmetric matrix using the Power iteration algorithm.

Note: Valgrind output for each case for both sections along with the power iteration algorithm (with case-specific code commented) code are present in folder "q2"

a) Valgrind output showing cache and branch statistics for the two possibilities of matrix multiplication is shown below.

Case 1: b[x] += M[x][y] * v[y]

```
---
7-- warning: L3 cache found, using its data for the LL simulation.
   7== I refs:
7== I1 misses:
7== LLi misses:
7== LLi miss rate:
7== LLi miss rate:
7=-
                                 1,222
  7==
17== D refs:
17== D1 misses:
17== LLd misses:
17== D1 miss rate:
17== LLd miss rate:
                        1,874,839,099
21,513,963
21,149,917
    /==
/== LL refs:
    == LL misses:
== LL miss rate:
```

Case 2: b[x] += M[y][x] * v[y]

```
ayadarshan@hota:-/Desktop/HPC/hpc-assignments/a0/q2$ valgrind --tool=cachegrind --cache-sim=yes --branch-sim=yes --log-file=Case2 ./a.out
gest Eigen Value 1.031383
gyadarshan@hota:-/Desktop/HPC/hpc-assignments/a0/q2$ cat Case2
3322== Cacheqprind, a cache and branch-prediction profiler
3322== Copyright (C) 2002-2017, and CNU GPL'd, by Nicholas Nethercote et al.
3322== Using Valgrind-3.13.0 and LibVEX; rerun with -h for copyright info
3322== Parent PID: 12711
               warning: L3 cache found, using its data for the LL simulation.
              : II misses:
: LLi misses:
: II miss rate:
: LLi miss rate:
             = D refs:
= D1 misses:
= LLd misses:
= D1 miss rate:
= LLd miss rate:
               LL misses:
LL miss rate:
```

Analysis

Case 1 has considerably lower L1 - D cache misses than Case 2. This is because Case 1 accesses elements in a row-wise manner and thus expected to have better cache performance due to the matrix being stored in row order format and hence elements in the locality of a referenced location would possibly be in the cache.

b) Timewise profile and hotspots are identified

Case 1: b[x] += M[x][y] * v[y]

```
addition += m[i][j] * v[j];
                 while(1) {
    double new_lambda;
                        double add = 0;
for(int i = 0; i < n; ++i) {
    add += b[i] * b[i];
}
           add += b[t] * b[t];
}
new_lambda = sqrtl(add);
=> /build/glibc-0TSEL5/glibc-2.27/math/w_sqrtl_compat.c:sqrtl (12x)
2,460
5,400
                        for(int i = 0; i < n; ++i) {
   v[i] = b[i] / new_lambda;</pre>
                        for(int i = 0; i < n; ++i) {
2,460
1,200
                              double addition = 0;
                              for(int j = 0; j < n; ++j) {
    // addition += m[j][i] * v[j];
    addition += m[i][j] * v[j];
}</pre>
23,000
40.000
3,000
                        if(fabs(new_lambda - cur_lambda) < 1e-6) {
    break;</pre>
                       }
else {
cur_lambda = new_lambda;
          printf("Largest Eigen Value %lf\n", cur_lambda);
=> /build/glibc-0TsEL5/glibc-2.27/stdio-common/printf.c:printf (1x)
                  return 0;
```

Case 2: b[x] += M[y][x] * v[y]

```
| Section | Sect
```

Analysis:

The hotspot in the power iteration program is the matrix multiplication in the iterative part of the code. Case 1 and Case 2 do not have any differences in the hotspot location as there is a difference only in the way the matrix is accessed and it does not change the number of times a part of code is accessed, unlike cache where locality makes a difference.

Q3

We wrote programs using i) recursion and ii) dynamic programming techniques for matrix chain multiplication. The output of perf having

- Task clock
- CPU cycles
- Instructions count
- Total cache (all levels) references and misses
- L1 D-cache loads, load-misses and stores
- L1 I-Cache load-misses

on these programs are as follows:

<u>Note:</u> Perf output along with the code for matrix chain multiplication using recursion and DP are present in folder "q3"

Recursion

```
dibyadarshan@hota:-/Desktop/HPC/hpc-assignments/a0/q3$ gcc matrixChainRecursive.c
dibyadarshan@hota:-/Desktop/HPC/hpc-assignments/a0/q3$ sudo perf stat -e task-clock,cycles,instructions,cache-references,cache-misses ./a.out
Operations required: 796659

Performance counter stats for './a.out':

15.117860 task-clock (msec) # 0.949 CPUs utilized
2,55,69,939 cycles # 1.691 GHz
6,86,57,533 instructions # 2.69 insn per cycle
6,86,57,533 instructions # 4.283 M/sec
27,946 cache-references # 4.283 M/sec
27,946 cache-misses # 43.161 % of all cache refs

0.015933323 seconds time elapsed

dibyadarshan@hota:-/Desktop/HPC/hpc-assignments/a0/q3$ sudo perf stat -e L1-dcache-load-misses,L1-dcache-stores,L1-icache-load-misses ./a.out
Operations required: 796659

Performance counter stats for './a.out':

2,68,66,412 L1-dcache-loads
15,085 L1-dcache-load-misses # 0.06% of all L1-dcache hits
1,29,46,155 L1-dcache-stores
28,046 L1-icache-load-misses
0.020607413 seconds time elapsed

dibyadarshan@hota:-/Desktop/HPC/hpc-assignments/a0/q3$
```

DP

Analysis

We can see that the recursive solution has considerably higher cycle count, instructions, L1 D-cache loads and stores, etc. as recomputing overlapping subproblems leads to greater accesses to memory than the DP solution. Although the L1 D-cache load miss rate of DP is higher than the recursive solution. This is because of memoization we reference the DP matrix to check whether this solution has been computed earlier or not, this reference doesn't exploit spatial or temporal locality and possibly leading to cache miss.

<u>Q4</u>

We wrote a recursive TSP program running on 12 nodes. Running gprof we obtained the following flat profile and call graphs

Note: gprof output along with the code for tsp using recursion are present in folder "q4"

Flat Profile

```
dibyadarshan@hota:-/Desktop/HPC/hpc-assignments/a0/q4$ cc. -Wall -pg tsp.c
dibyadarshan@hota:-/Desktop/HPC/hpc-assignments/a0/q4$ ./a.out
Number of Cittes: 12
0 84 87 78 16 94 93 68 79 35 50 22 63
08 90 16 06 42 74 12 77 33 71 269
68 30 0 83 31 63 24 68 36 30 3 23
59 70 68 0 94 57 12 43 30 74 22 20
85 38 99 25 0 16 71 14 27 92 81 57
74 63 71 97 82 0 6 26 85 28 37 6
47 30 14 58 25 96 0 83 46 15 60 35
65 44 51 88 9 77 79 0 89 85 4 52
55 100 33 61 77 69 40 13 0 27 87 95
40 96 71 35 79 68 2 98 3 0 18 93
53 57 2 81 87 42 66 90 45 20 0 41
30 32 18 98 72 82 76 10 28 68 57 0
cost: 183
dibyadarshan@hota:-/Desktop/HPC/hpc-assignments/a0/q4$ gprof a.out gmon.out > prof_output
dibyadarshan@hota:-/Desktop/HPC/hpc-assignments/a0/q4$ cat prof_output
Flat profile:

Each sample counts as 0.01 seconds.
% cumulative self self total
time seconds calls s/call s/call name
98.93 3.88 3.88 1 3.88 travellingSalesmanProblem
1.28 3.93 0.05 frame_dummy
anin
```

Time spent in the travellingSalesmanProblem function is shown above.

Call graph

For the *travellingSalesmanProblem function* is called recursively by itself 108505111 times and that function is called by *main* once.