Lab 1: Performance measurements and Amdahl's Law

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Outline

- Performance equation
- Amdahl's Law
- Programming assignment get familiar with C/C++

Q3-Q13: CPU Performance Equation

$$Performance = \frac{1}{Execution \ Time}$$

$$Execution \ Time = \frac{Instructions}{Program} \times \frac{Cycles}{Instruction} \times \frac{Seconds}{Cycle}$$

$$ET = IC \times CPI \times CT$$

- IC (Instruction Count) Q3 Q4
 - ISA, Compiler, algorithm, programming language, programmer
- CPI (Cycles Per Instruction)— Q7 Q13
 - Machine Implementation, microarchitecture, compiler, application, algorithm, programming language, programmer
- Cycle Time (Seconds Per Cycle)— Q5 Q6
 - Process Technology, microarchitecture, programmer

Q4: Performance equation (round 2)

 Consider the following c code snippet and x86 instructions implement the code snippet

Comparing the case where count equal to 1,000,000,000 and 2,000,000,000, what factor in performance equation would change?

```
A. IC
B. CPI
C. CT
```

Q4: Speedup per element?

How much time do we have to spend on each s += a[i];

Should be the same — theoretically

The reality is ...



0.5

1.0

1.5

2.0

2.5

0.5

1.0

1.5

size

2.0

2.5

0.5

1.0

1.5

size

2.0

2.5

0.5

1.0

2.0

2.5

Q14—Q16: Practicing Amdahl's Law

$$Speedup_{enhanced}(f, s) = \frac{1}{(1 - f) + \frac{f}{s}}$$

f is the fraction of "execution time" — neither of the IC, CPI or CT

• Q14 — Q16: How to find the f in our program?

```
int main() {
    A();
    B();
    C();
    return 0;
}
```

function	time	fraction?	
main	10		
Α	5	50%	
В	2	20%	
С	3	30%	

Q17—Q19: Throughput and latency

- Throughput is the amount of work/instructions/data that the system/machine/processor can deliver within a period of time
- It does not mean that the processor will deliver work at the "average" latency
- Throughputs reflect to performance relative well on large amount of work, but poorly on small amount of work

What's a "floating point operation"

- An arithmetic operation with at least one operand as "floating point" data types
- Memory access does not count as a FLOP
- How many floating point operations in the following statement?

```
uint32_t *A, size;
float *B;
for(int i = 0; i< size; i++) {
    B[i] = B[i] + A[i]*2;
}

floating point
    integer mul
    add</pre>
```

Let's measure the FLOPS of matrix multiplications

```
double **a, **b, **c;
for(i = 0; i < ARRAY_SIZE; i++) {
  for(j = 0; j < ARRAY_SIZE; j++) {
    for(k = 0; k < ARRAY_SIZE; k++) {
      c[i][j] += a[i][k]*b[k][j];
    }
  }
}</pre>
```

Floating point operations per second (FLOP"S"):

$$i \times j \times k \times 2$$

Given
$$i = j = k = 2048$$

$$2^{3 \times 11} \times 2 = 2^{34}$$
 FLOPs in total

$$FLOPS = \frac{i \times j \times k \times 2}{ET_{seconds}}$$

What if a and c are integers?

```
uint64_t **a, **b;
double **c;
for(i = 0; i < ARRAY_SIZE; i++) {
  for(j = 0; j < ARRAY_SIZE; j++) {
    for(k = 0; k < ARRAY_SIZE; k++) {
      c[i][j] += a[i][k]*b[k][j];
    }
}</pre>
```

Floating point operations per second (FLOP"S"):

$$i \times j \times k \times 1$$

Given
$$i = j = k = 2048$$

$$2^{3\times11} = 2^{33}$$

FLOPs in total

$$FLOPS = \frac{i \times j \times k \times 1}{ET_{seconds}}$$

Demo: matmul on GPU

Size	Total FLOPs	Latency	Relative Latency	Throughput (Output Numbers Per	Relative Throughput
16x16x16	(16) ³ x2				
32x32x32	(32) ³ x2				
64x64x64	(64) ³ x2				

Larger throughput doesn't mean shorter latency!



Demo: matmul on GPU

Size	Total FLOPs	Latency	Relative Latency	Throughput (Output Numbers Per Second)	Relative Throughput
16x16x16	(16) ³ x2	~ 0.09ms	1	0.09ms/8192	1
32x32x32	(32) ³ x2	~ 0.09ms	1	0.09ms/65536	8
64x64x64	(64) ³ x2	~ 0.09ms	1	0.09ms/524288	64

Larger throughput doesn't mean shorter latency!

Programming assignment



Programming assignments

- Each lab will have a programming assignment in C/C++ for you to practice your skills of code optimizations
- escalab.org/datahub provides
 - VSCode server if you're more familiar with it
 - You may also use the editors in jupyterhub for code development
- The performance & grading are based on "gradescope's" server, not our servers.



Why C/C++ programming?

The only pathway to performance programming

Use "gdb" as a debugging tool

- compile your program with "-g"
- gdb "binary_execution"
 gdb >> break hello2.cpp:"line_number"
 gdb >> run arguments
- gdb "binary_execution" gdb >> run arguments crashed! gdb >> bt

Lab 1: the main function & basic I/O

```
#include <fstream>
#include <iostream>
int main(int argc, char *argv[])
  std::ofstream ofs ("hello.txt", std::ofstream::out);
  ofs << "Hello CSE142L!\n";
  ofs.close();
  std::cout << "Execution Complete" << std::endl;</pre>
  return 0;
```

Hints to Lab 1 PA

- You need to manipulate the argv array
- You need to find out how to convert "ASCII" based characters into integers

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Questions for Lab 1 and PA 1?

Announcement

- Lab report 1 and PA 1 due 8/10 midnight through gradescope
- Lab 2 will be released on Sunday



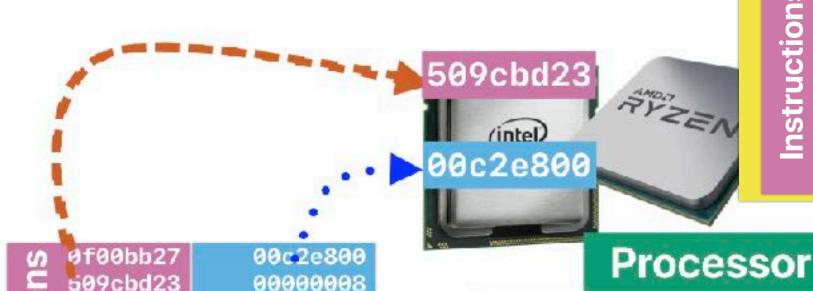
Computer Science & Engineering

1421





How my "C code" becomes a "program" Objects, Libraries



00c2f000

00000008

00c2f800

00000008

00c30000

00000008

cafebabe 00000033 001d0a00 06000f09 00100011 0800120a 00130014 07001507 00c2e800 00000008 00c2f000 00000008 00c2f800 00000008 00c30000 00000008

Source Code

Linker

Compiler (e.g., gcc)

00c2e800



Program

9f00bb27 509cbd23 00005d24 0000bd24 2ca422a0 130020e4 00003d24 2ca4e2b3

00000008 0002f000 00000008 0002f800

00000008 00c30000 00000008

Memory

00005d24

0000bd24

2ca422a0

130020e4

00003d24

2ca4e2b3

Program

0f00bb27 00c2e800 00000008 509cbd23 00005d24 00c2f000 0000bd24 00000008 2ca422a0 00c2f800 130020e4 00000008 00003d24 00c30000 2ca4e2b3 00000008

Storage