## **EXAMINATION SCHEDULING**

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## **OUTLINE**

Problem

Modeling The Problem

Next Steps

# **PROBLEM**

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Find a good examination schedule for the exam period of the TUM.

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- Rooms for an exam are minimized
- Time between exams is maximized

# MODELING THE PROBLEM

#### **VARIABLES**

$$x_{i,k,l} := \left\{ \begin{array}{l} 1, & \text{if exam $i$ is written in period $l$ in room $k$} \\ 0, & \text{otherwise} \end{array} \right.$$

#### **VARIABLES**

$$\begin{split} x_{i,k,l} &:= \left\{ \begin{array}{l} 1, & \text{if exam $i$ is written in period $l$ in room $k$} \\ 0, & \text{otherwise} \end{array} \right. \\ y_{i,l} &:= \left\{ \begin{array}{l} 1, & \text{if exam $i$ is written in period $l$} \\ 0, & \text{otherwise} \end{array} \right. \end{split}$$

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$$(1) \qquad \sum_{\text{rooms }k} x_{i,k,l} \qquad \leq y_{i,l} \cdot M \qquad \forall \text{ exams } i, \ \forall \text{ periods } l$$

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(2) 
$$\sum_{\text{rooms } k} x_{i,k,l} \geq y_{i,l} \qquad \forall \text{ exams } i, \forall \text{ periods } l$$

• Each exam is planned in exactly one period:

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(3) 
$$\sum_{\text{periods } l} y_{i,l} = 1 \quad \forall \text{ exams } i$$

• There are enough seats for the students in the exam rooms:

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$$(4) \qquad \sum_{\substack{\text{periods } l, \\ \text{rooms } k}} c_k \cdot x_{i,k,l} \qquad \geq s_i \qquad \forall \text{ exams } i$$

 $s_i := \#$  students taking exam i  $c_k := \#$  seats in room k

• In every room there is at most one exam at a given time:

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(5) 
$$\sum_{\text{exams i}} x_{i,k,l} \leq 1 \quad \forall \text{ rooms } k, \forall \text{ periods } l$$

• There are no conflicts for students taking multiple exams:

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Idea: 
$$y_{i,l} = 1 \Rightarrow \sum_{j \text{ conflicts } i} y_{j,l} = 0 \quad \forall \text{ periods } l$$

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(6) 
$$\sum_{\substack{\text{exams j,} \\ \text{j conflicts i}}} y_{j,l} \leq (1 - y_{i,l}) \cdot M \quad \forall \text{exams i, } \forall \text{periods l}$$

• Minimize the total number of rooms:

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 $\min \sum_{\substack{\text{exams } i, \\ \text{rooms } k, \\ \text{periods } l}} x_{i,k,l}$ 

• Maximize the time between two exams:

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$$\max \sum_{\text{exams } i} \min_{\substack{\text{exam } j \\ \text{conflicts with } i}} d_{i,j}$$

 $d_{i,j} := distance between exams i and j$ 

• Combine the previous two objective functions using a weighting factor  $\gamma > 0$ :

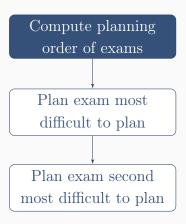
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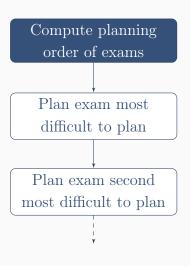
$$(a,b) \qquad \min \sum_{\substack{\text{exams } i, \\ \text{rooms } k, \\ \text{periods } l}} x_{i,k,l} \ - \ \gamma \cdot \sum_{\substack{\text{exams } i \\ \text{conflicts with } i}} \min_{\substack{\text{exam } j \\ \text{conflicts with } i}} d_{i,j}$$

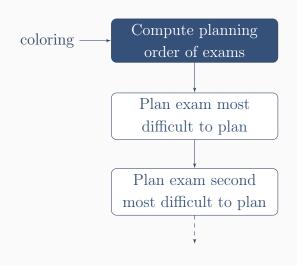
# **NEXT STEPS**

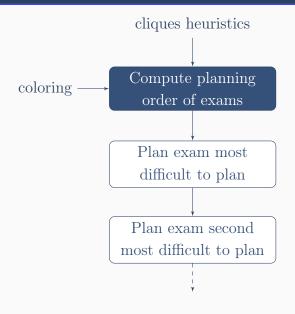
Compute planning order of exams

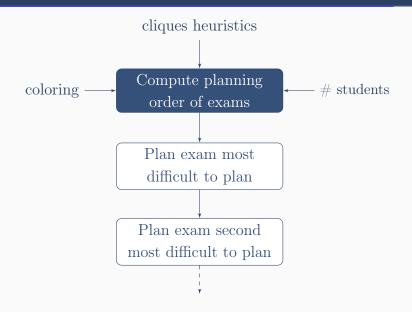












(7) 
$$\sum_{i \text{ in clique}} y_{i,l} \leq 1 \quad \forall \text{ cliques, } \forall \text{periods } l$$

• Add clique constraints:

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- Remove absolute value in the objective function

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- Improve running time using heuristics, pre-solving, etc.
- Use a path-based model and column generation

## DATA ACQUISITION



Figure 1: Server, hopefully somewhere in Germany