

CS 111 Midterm Exam

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TOTAL POINTS

84 / 100

QUESTION 1

1 Pages and page frames 8 / 10

- 0 pts Correct

- 10 pts No answer.

✓ - 2 pts Did not explain that pages get placed into page frame.

- 1 pts Said page is mapped to page frame

- 1 pts No mention of location of page & page frame in system

- 4 pts Incorrect Reasoning

- 3 pts Not accurate enough

- 1 pts No mention of virtual address

- 5 pts Page frame contains exactly one page.

- 2 pts Did not mention that page and page frame are the same size.

QUESTION 2

2 ABIs 9 / 10

- 0 pts Correct

- 10 pts No answer

✓ - 1 pts Did not mention operating system

- 4 pts Said applications don't need to worry about hardware differences

- 2 pts Off-topic

- 3 pts Incorrect reasoning

- 2 pts Did not mention software distribution

- 3 pts Did not mention hardware and OS

- 2 pts Left out hardware

QUESTION 3

3 Information hiding 3 / 10

- 0 pts Correct

- 10 pts No answer.

- 6 pts Primary benefit is to avoid bugs arising from improper dependencies among modules.

- 3 pts Major element of benefit is that it allows changes in module implementation.

✓ - 7 pts Really about hiding details of OS modules' implementation, not about concealing processes' address spaces from each other.

- 2 pts Nothing to do with open vs. closed source.

- 9 pts Primarily an issue involving abstraction.

QUESTION 4

4 Context switches 4 / 10

- 0 pts Correct

- 2 pts Not mentioning general registers

✓ - 2 pts Not mentioning PC

✓ - 2 pts Not mentioning Stack ptr

✓ - 1 pts Not mentioning PSW

- 2 pts No discussion of memory mapping data

- 2 pts No need to explicitly save data, since it's already sitting in memory.

- 3 pts Generally nothing goes to disk on a context switch.

- 1 pts What about memory needs to be saved?

- 2 pts Much of this stuff need not be saved, since it's already in memory. OS just needs to be sure it can be found again when process is switched back in.

- 2 pts The PCB is an OS data structure that exists as long as the process is around, so it need not be saved on a context switch.

- 1 pts File size has nothing to do with a context switch.

- 1 pts I have no idea what the flag you're talking about is.

✓ - 1 pts "state of the process" is vague.

- 2 pts Caches aren't saved.

- 2 pts No need to update a file descriptor during a context switch.

QUESTION 5

5 Trap tables 10 / 10

✓ - 0 pts Correct

- 10 pts No answer
- 3 pts Answer incomplete, should mention trap table is used to specify what code to run when trap occurs.
- 8 pts Wrong answer.
- 3 pts Answer incomplete.
- 2 pts User process has no thing to do with trap?

QUESTION 6

6 Race conditions 10 / 10

✓ - 0 pts Correct

- 10 pts No answer
- 5 pts Answer incomplete.
- 8 pts Answer incorrect.
- 2 pts Missing some details.

QUESTION 7

7 Blocking and threads 10 / 10

✓ - 0 pts Correct

- 10 pts No answer or Wrong answer
- 5 pts Missing: User-mode threads block other threads of the same process.
- 5 pts Missing: Kernel-mode threads do not block other threads of the same process, as other threads can be scheduled to run on the same or another core.

QUESTION 8

8 STCF 10 / 10

✓ - 0 pts Correct

- 10 pts No answer or Wrong answer
- 5 pts Missing: interrupt the running one OR switch to the newly-added shorter ones.
- 8 pts Missing mention of new processes that might have shorter time to completion.

QUESTION 9

9 Fork and exec 10 / 10

✓ - 0 pts Correct

- 10 pts No answer

- 4 pts Not mentioning code replacement.
- 3 pts Not mentioning stack replacement.
- 3 pts Not mentioning heap replacement.
- 9 pts The question was about what happens after the exec, not the fork.
- 8 pts What resources are replaced by the exec?
- 7 pts Stack and code are changed by exec.
- 5 pts Fork/exec work with processes, not threads.
- 2 pts Even any data written after fork gets replaced by exec.
- 2 pts The old stack is totally overwritten.
- 10 pts Totally wrong. Nothing to do with multithreading and multicore.
- 6 pts So, what resources are replaced?

QUESTION 10

10 Fragmentation for memory management schemes 10 / 10

✓ - 0 pts Correct

- 10 pts No answer
- 5 pts Not identifying internal fragmentation for pages.
- 5 pts Paged segments suffer 1/2 page fragmentation.
- 5 pts Fixed segments suffer 1/2 internal segment fragmentation.
- 3 pts On average 50%
- 2 pts The 1.5% was a particular example. It will be 1/2 page, on average.
- 2 pts Internal fragmentation has little to do with how long the system runs, unlike external.
- 2 pts Paging and fixed size partitions never experience external fragmentation.
- 2 pts The paging form of fragmentation you describe is internal fragmentation.
- 2 pts Paging doesn't use binary buddy. It allocates in fixed size pages.
- 4 pts Fixed segments are likely to waste more memory on internal fragmentation than paging, not less.
- 3 pts No external fragmentation with paging.

- **5 pts** No answer on segmented system.
- **2 pts** Calling internal fragmentation "external".
- **1 pts** This form of fragmentation is called "internal."
- **4 pts** Paged segment fragmentation only occurs in the last page.