

Relation between RE and FA

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Conversion from DFA to RE

- Transitive Closure Method
 - Rather theoretical approach.
 - Sketch of the method:
 - Let $Q = \{q_1, q_2, \dots, q_m\}$ be the set of all automaton states.
 - Suppose that regular expression R_{ij} represents the set of all strings that transition the automaton from q_i to q_j .
 - Wanted regular expression will be the union of all R_{sf} , where q_s is the starting state and q_f is one of the final states.
 - • The main problem is how to construct R_{ij} for all states q_i, q_j .

How to construct R_{ij} ?

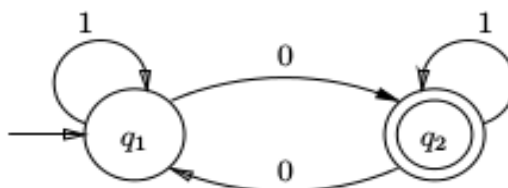
R_{ij}^k is recursively defined as:

$$R_{ij}^k = R_{ij}^{k-1} + R_{ik}^{k-1} (R_{kk}^{k-1})^* R_{kj}^{k-1}$$

Assuming we have initialized R_{ij}^0 to be:

$$R_{ij}^0 = \begin{cases} r & \text{if } i \neq j \text{ and } r \text{ transitions NFA from } q_i \text{ to } q_j \\ r + \varepsilon & \text{if } i = j \text{ and } r \text{ transitions NFA from } q_i \text{ to } q_j \\ \emptyset & \text{otherwise} \end{cases}$$

Example



$$r_{11}^0 = 1 + \epsilon$$

$$r_{12}^0 = 0$$

$$\begin{aligned} r_{12}^1 &= r_{12}^0 + r_{11}^0 (r_{11}^0)^* r_{12}^0 \\ &= 0 + (1 + \epsilon)^+ 0 \end{aligned}$$

$$r_{22}^0 = 1 + \epsilon$$

$$r_{21}^0 = 0$$

$$\begin{aligned} r_{22}^1 &= r_{22}^0 + r_{21}^0 (r_{11}^0)^* r_{12}^0 \\ &= (1 + \epsilon) + 0(1 + \epsilon)^* 0 \end{aligned}$$

$$r_{12}^2 = r_{12}^1 + r_{12}^1 (r_{22}^1)^* r_{22}^1$$

$$= (0 + (1 + \epsilon)^+ 0) + (0 + (1 + \epsilon)^+ 0)((1 + \epsilon) + 0(1 + \epsilon)^* 0)^+$$

$$= (0 + (1 + \epsilon)^+ 0)(1 + \epsilon + 0(1 + \epsilon)^* 0)^*$$

$$= (1 + \epsilon)^* 0(1 + \epsilon + 01^* 0)^*$$

$$= 1^* 0(1 + 01^* 0)^*$$

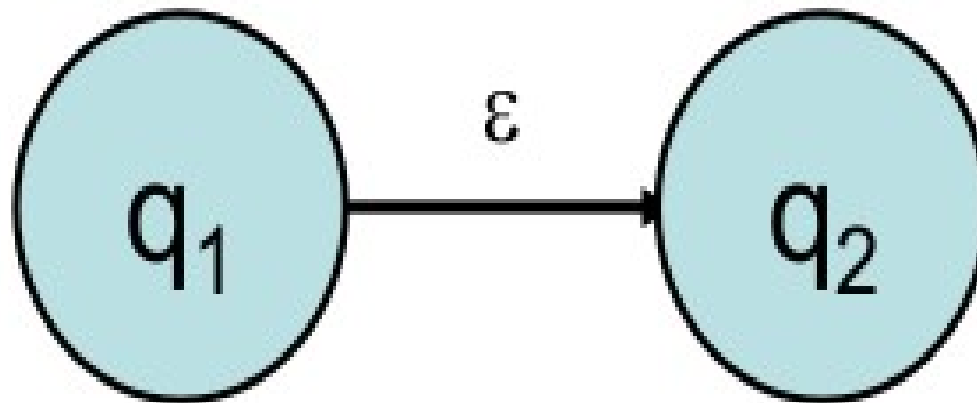
$$L(M) = L(r_{12}^2)$$

Next Part:

Conversion from RE to NFA

First of all.... ϵ -NFA

- transitions made “for free”, without “consuming” any input symbols.



Finite Automata and Regular Expression

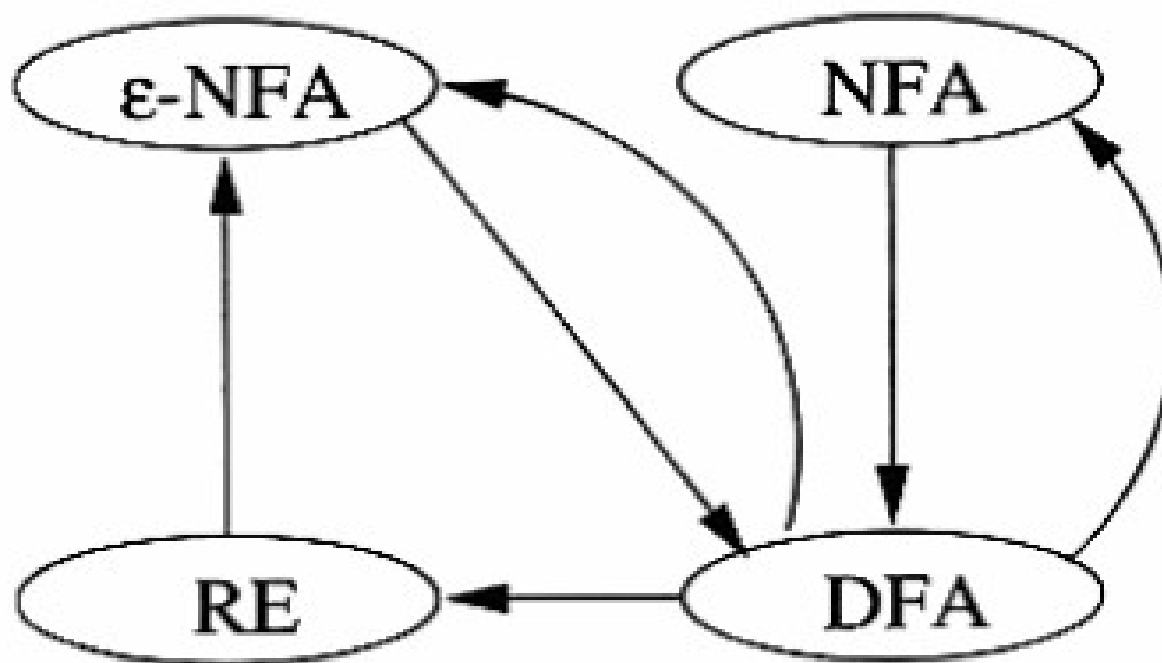
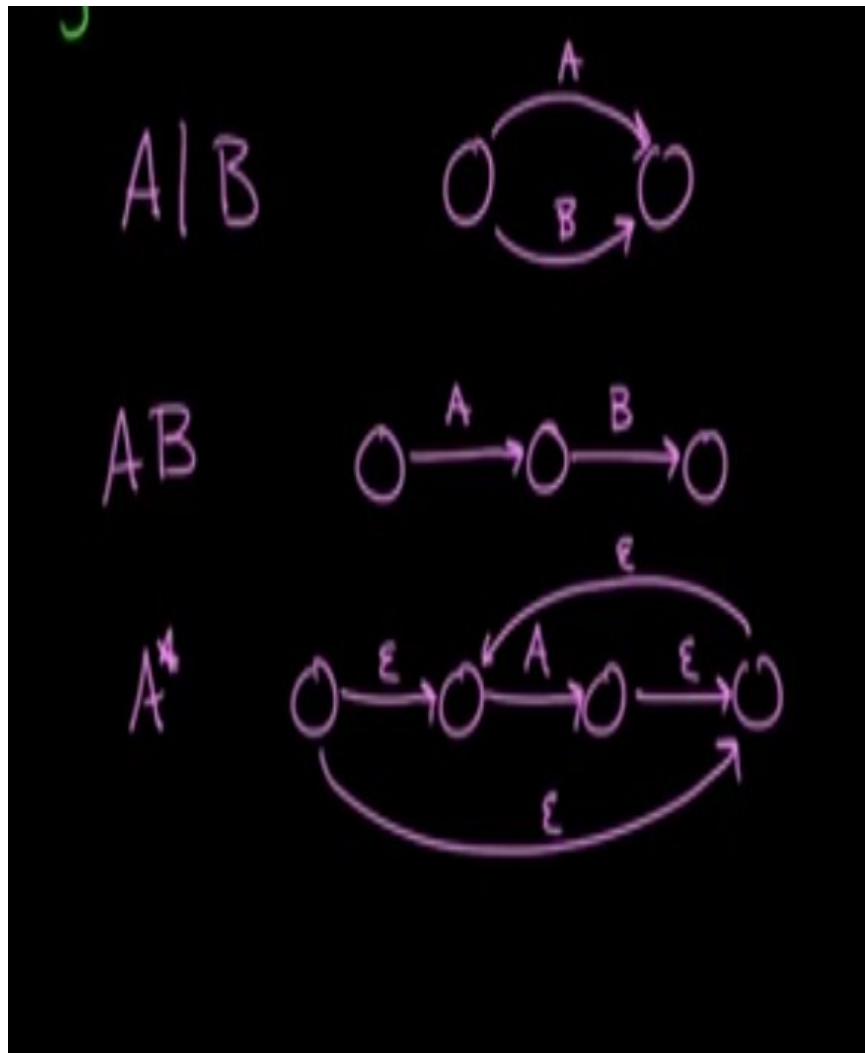
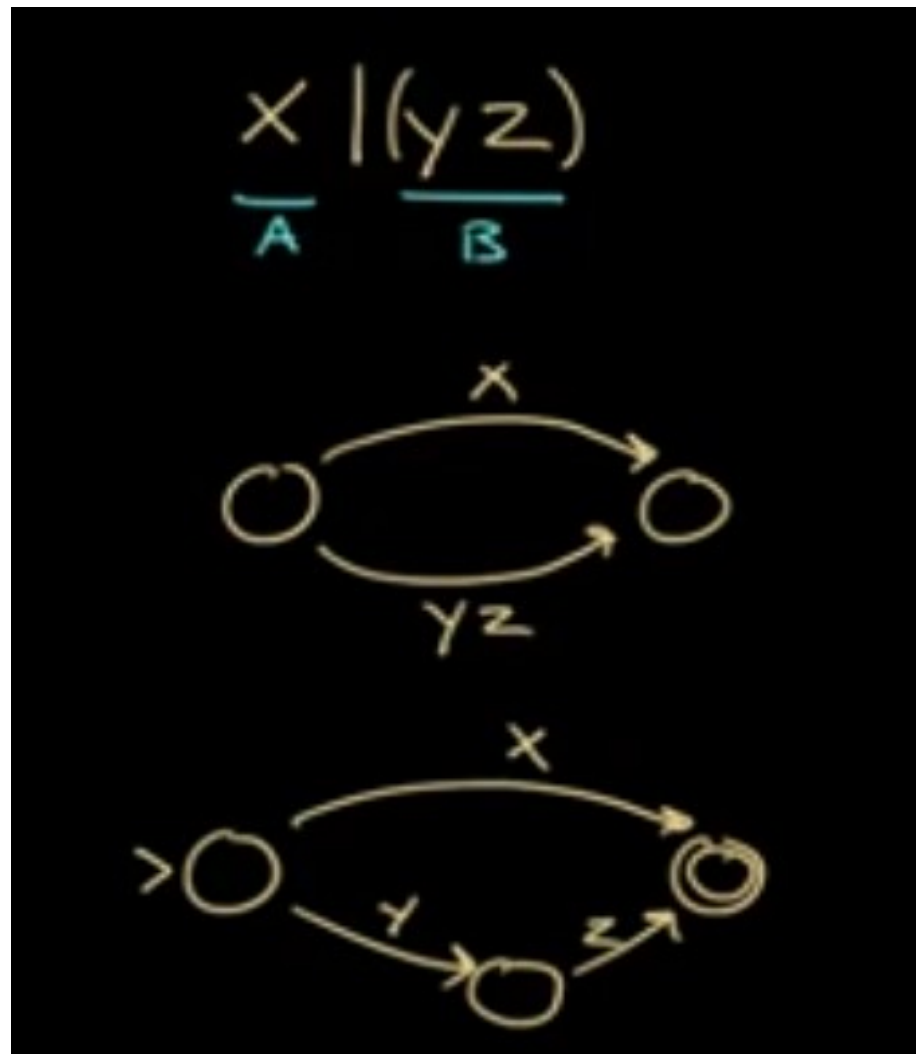
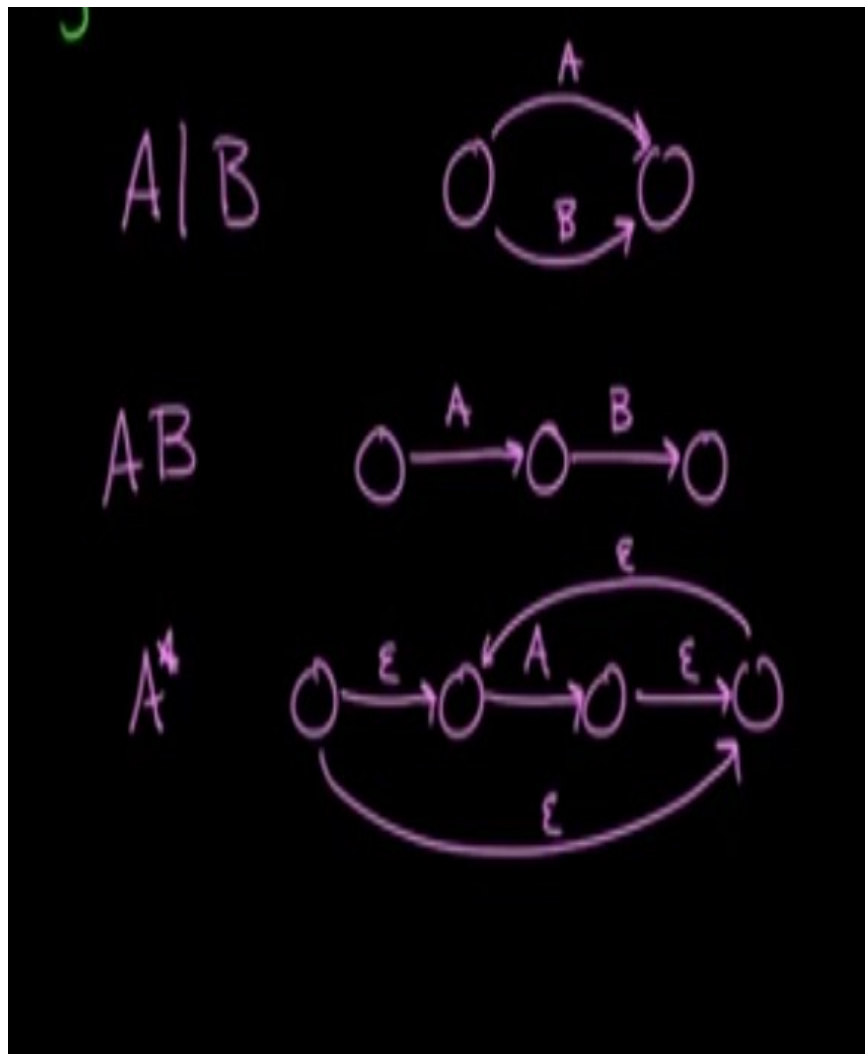


Figure 3.1: Plan for showing the equivalence of four different notations for regular languages

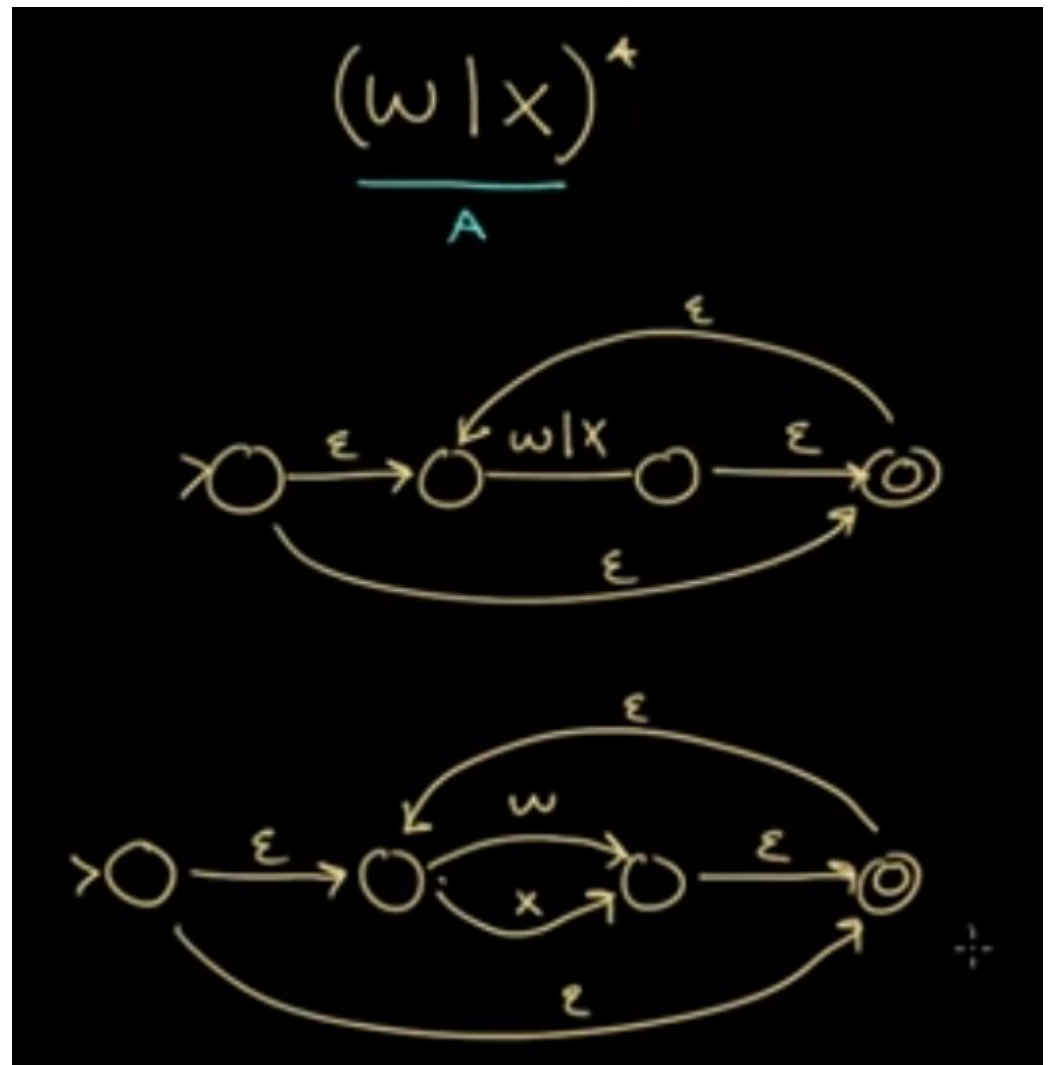
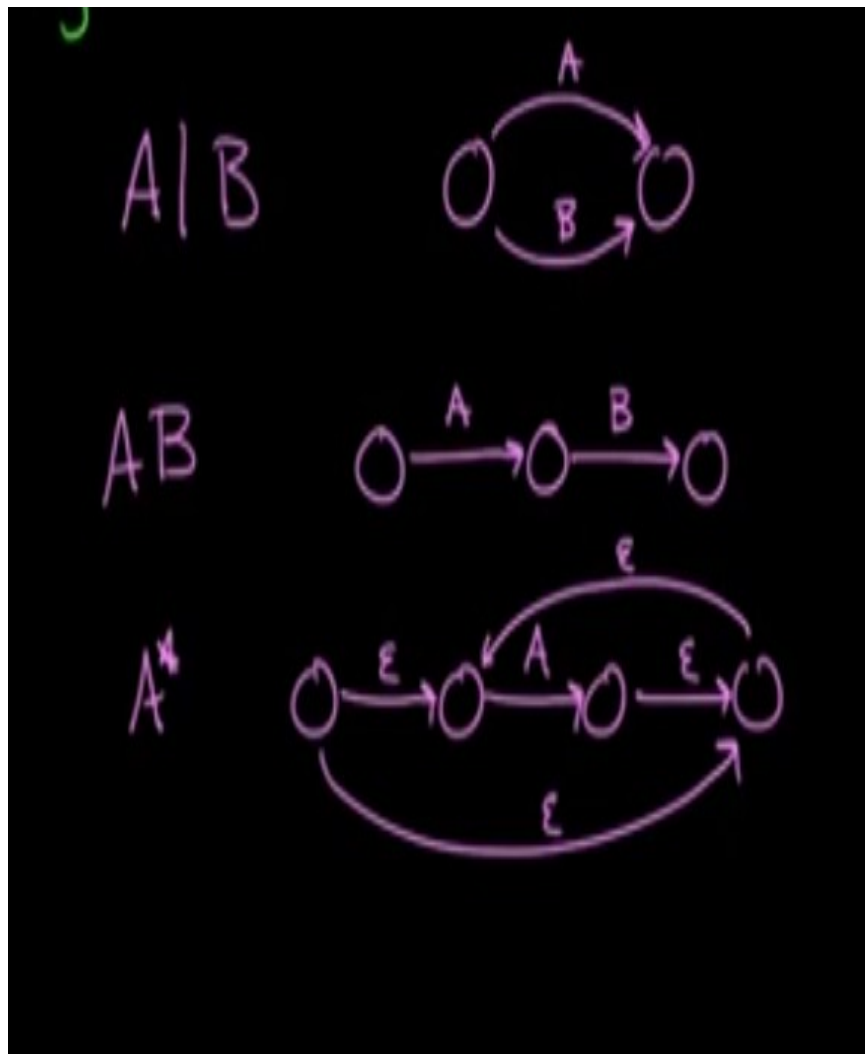
Conversion from RE to NFA



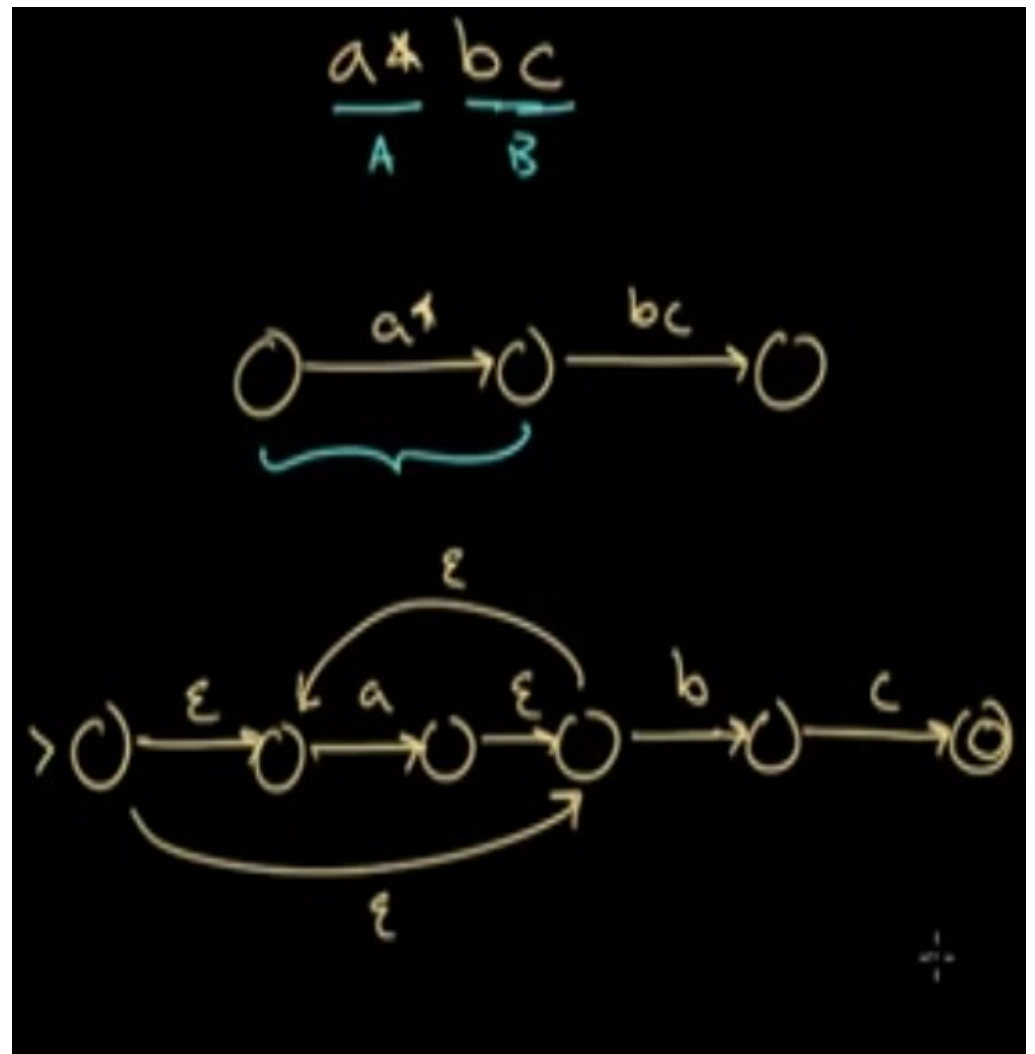
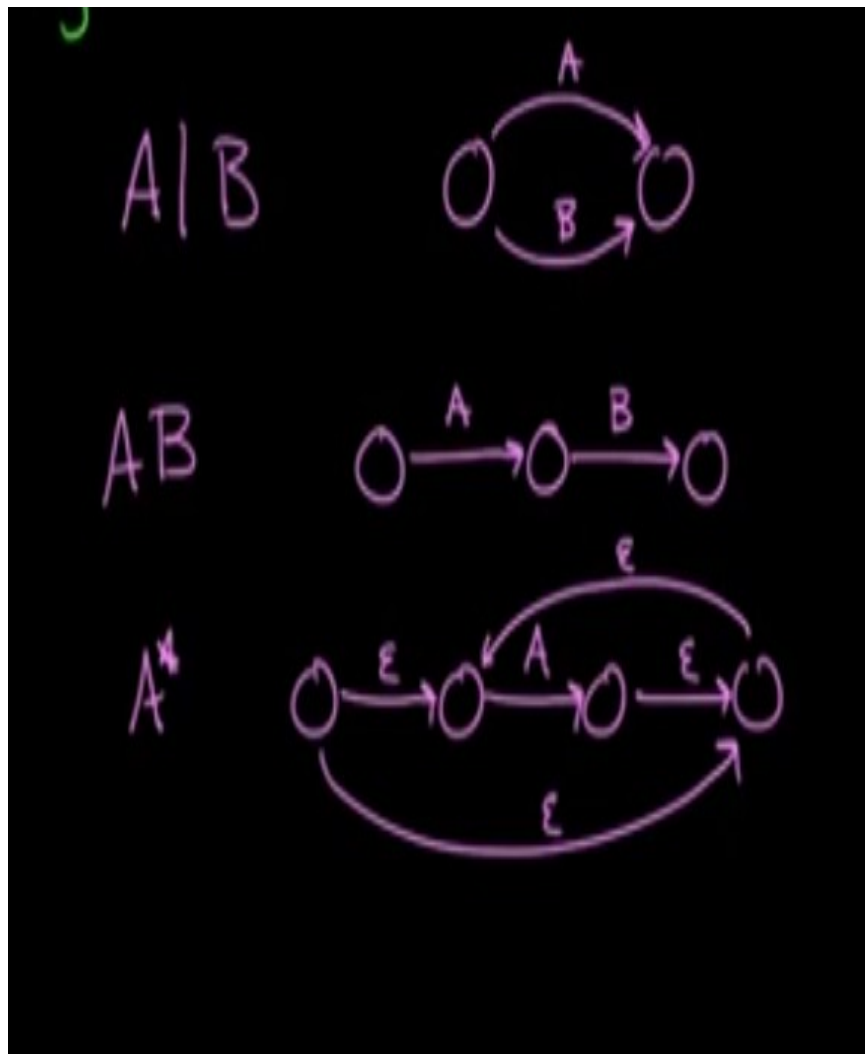
Conversion from RE to NFA



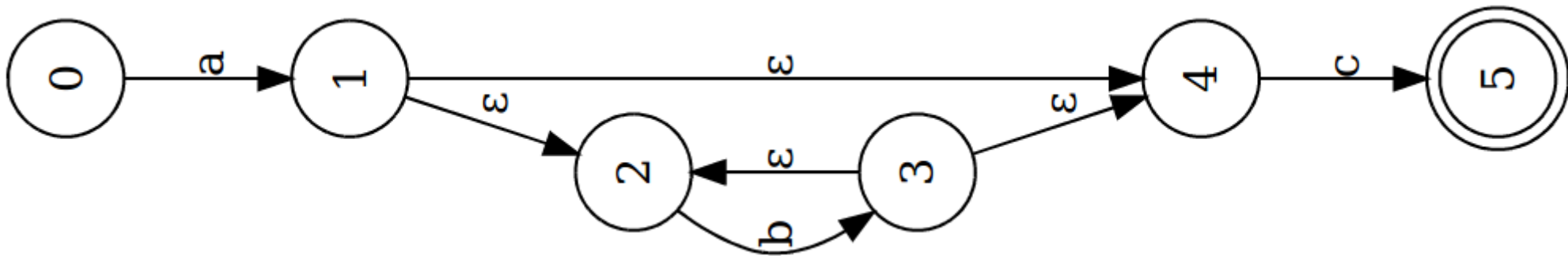
Conversion from RE to NFA



Conversion from RE to NFA



For ab^*c



NFA

END