CD: COMPILER DESIGN SPM: UNIT-1 18/09/2022

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(CD)
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#### PREPARED FOR

Engineering Students
All Engineering College

# CD: COMPILER DESIGN

TOPIC On: UNIT-1 LEXICAL ANALYZER | Thompson method | RE TO NFA | NFA TO DFA

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Under On: INTRODUCTION TO COMPILER

# TOPIC On: UNIT-1 LEXICAL ANALYZER | Thompson method | RE TO NFA | NFA TO DFA

# Regular expression to NFA using Thompson's rule

- ☐ Regular expression to NFA (Thompson's construction)
- 1. For ∈, construct the NFA

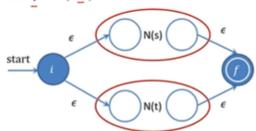


#### 2. For a in $\Sigma$ , construct the NFA

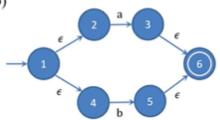


### ☐ Regular expression to NFA (Thompson's construction)

3. For regular expression s|t or (s+t)



For Example: (a|b) or (a+b)



### ☐ Regular expression to NFA (Thompson's construction)

4. For regular expression st

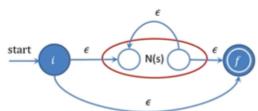


For Example: ab

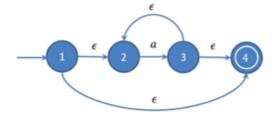


#### ☐ Regular expression to NFA (Thompson's construction)

5. For regular expression s\*

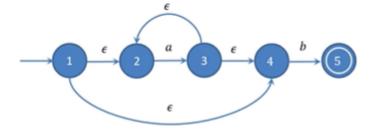


For Example: a\*

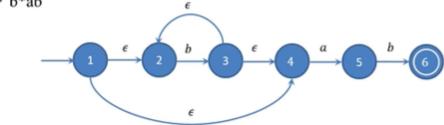


## ☐ Regular expression to NFA (Examples)

➤ a\*b

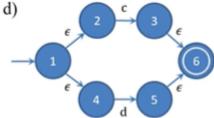


➤ b\*ab

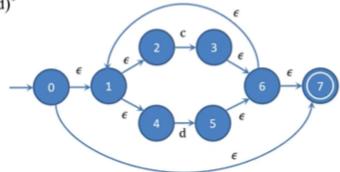


# ☐ Regular expression to NFA (Examples)

 $\triangleright$  (c|d) or (c+d)



 $\triangleright$  (c|d)\*or (c+d)\*



#### **□** Exercise

#### Convert following regular expression to NFA:

- 1. abba
- 2. bb(a)\*
- 3. (a|b)\*
- 4. a\* | b\*
- 5. a(a)\*ab
- 6. aa\*+ bb\*
- 7. (a+b)\*abb
- 8. 10(0+1)\*1
- 9. (a+b)\*a(a+b)
- 10. (0+1)\*010(0+1)\*
- 11. (010+00)\*(10)\*
- 12. 100(1)\*00(0+1)\*

# Conversion from NFA to DFA using subset construction method

#### ☐ Subset construction algorithm

Input: An NFA N.

Output: A DFA D accepting the same language.

*Method*: Algorithm construct a transition table *Dtran* for D. We use the following operation:

OPERATION	DESCRIPTION
$\in -closure(s)$	Set of NFA states reachable from NFA state $s$ on $\in$ - transition alone.
∈ − closure	Set of NFA states reachable from some NFA state $s$ in $T$ on $\in$ - transition alone.
$\delta(T,a)$	Set of NFA states to which there is a transition on input symbol $a$ from some NFA state $s$ in $T$ .

#### ■ Exercise

❖ Convert following regular expression to DFA using subset construction method.

#### OR

- Construct a NFA for following regular expression using Thompson's notation and then convert it into DFA.
- 1. (010+00)\*(10)\*
- 2. (a+b)\*a(a+b)
- 3. (a|b)\*a(a|b)#

# **DFA Optimization**

#### ☐ DFA Optimization Algorithm

- 1. Construct an initial partition  $\Pi$  of the set of states with two groups: the accepting states F and the non-accepting states S F.
- 2. Apply the repartition procedure to  $\Pi$  to construct a new partition  $\Pi new$ .
- 3. If  $\Pi new = \Pi$ , let  $\Pi final = \Pi$  and continue with step (4). Otherwise, repeat step (2) with  $\Pi = \Pi new$ .

#### for each group G of $\Pi$ do begin

partition G into subgroups such that two states s and t of G are in the same subgroup if and only if for all input symbols a, states s and t have transitions on a to states in the same group of  $\Pi$ .

replace G in  $\Pi new$  by the set of all subgroups formed.

end

#### ☐ DFA Optimization Algorithm

- 4. Choose one state in each group of the partition Πfinal as the representative for that group. The representatives will be the states of M'. Let s be a representative state, and suppose on input a there is a transition of M from s to t. Let r be the representative of t's group. Then M' has a transition from s to r on a. Let the start state of M' be the representative of the group containing start state s<sub>0</sub> of M, and let the accepting states of M' be the representatives that are in F.
- 5. If M' has a dead state d, then remove d from M'. Also remove any state not reachable from the start state.