

Grometrie mangin W -> 10W 6-106 doss not change Greanetric margin Maximum classifier: VIVI b y (i) (w x x (i) + b) = ) hw11=1 Lecture 7 - Optimal Margin Classifier - Primal / Dual optimization problem (KKT) - SVM dual - Kernels hw,6 (x)= g(wTx+6) 0 g(z) = { 1 if 2 20 otherwise 0 Je {-1, 1} & WTX tb=0 Eunc. Margin: D'= y (1) (wTx(1)+b) double with V= many (i) Gres. Margin: NU)=40 (WX X(1)+11 WI) will not D = min ()(1) Change the prosition of 1 w11 = 1 eguivalent edition of The hyperplane [w, = 1 the optimal problem: W,2+ W, 1=17

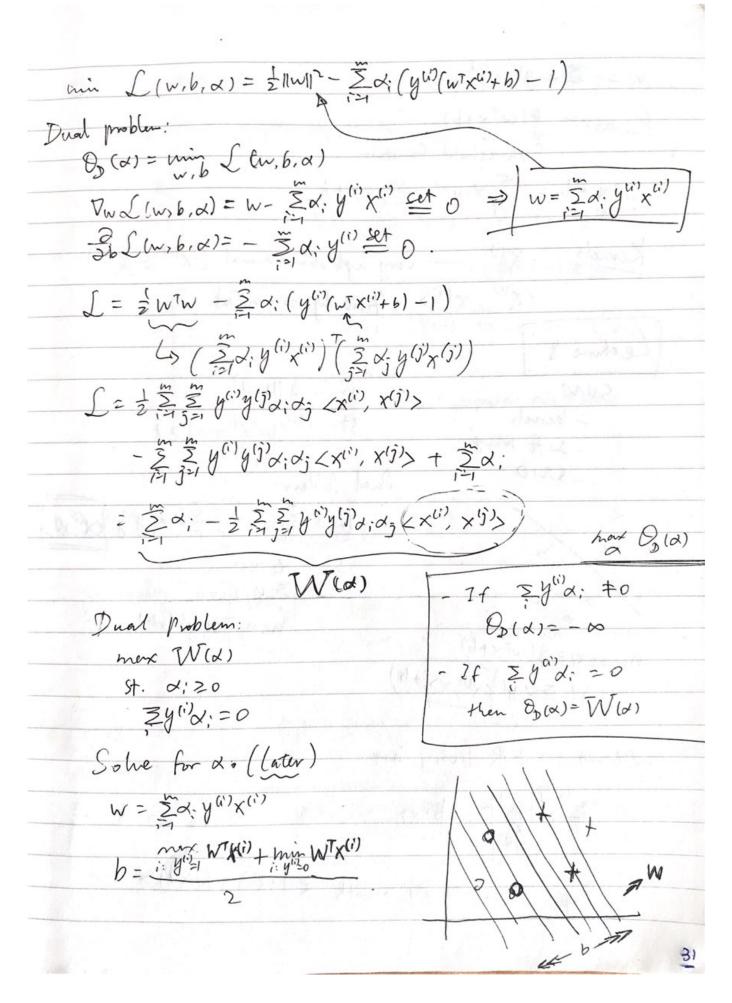
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#1. St. y" (wTx")+b) > V (i=1,...,m) 4t. y" (wx(i)+b) 2) B=1 } ← impose this constraint then min y (1) (wtx (1)+b)=1. #3: more Tiwil St. y (1) (wrx(1)4b) 21 Convex optimization: him few) St. hicw)=0, i=1,..., l. OP: quadratic programming software Lograngian: gradient descent algo.  $\mathcal{L}(w,\beta) = fw_1 + \frac{2}{2} f_i h_i(w)$   $\mathcal{L}(w,\beta) = fw_1 + \frac{2}{2} f_i h_i(w)$   $h_i(w) = \begin{pmatrix} h_i(w) \\ h_2(w) \end{pmatrix} = 0$   $\frac{\partial \mathcal{L}}{\partial w} \xrightarrow{\text{cot}} 0$   $\frac{\partial \mathcal{L}}{\partial w} \xrightarrow{\text{cot}} 0$ 

for wo to be a solution, it is necessary that 3 pt st. of (w\*, p\*) =0. of (w\*,p\*) =0. Primal Problem min few) St. g;(w) ≤0, i=1...k ("g(w) ≤0") hy(w) ±0. j=1.-1 ("how)=0) Lagrangian: L(ω,σ,β)=fw)+\(\int\_{\infty}\alpha;g:(w)+\(\infty\) (\(\beta;\frac{\partial}{\partial};\frac{\p Define: Op(w) = max & (w, a, B) Consider: px = min mans S (m, d, p) = min Op(w) Op(w): - If g:(w)>0 then Op(w) = 00 - If hi(w) +0 then Op(w) =00 - Otherwise, Op (w) = faw) Then op(w) = { f(w), if constraints satisfied (g,h) or wherenice um Op(w) = original problem

Dual Problem Op(d, B)= win L(w, d, B) d = max min L(w,d,p) = max OD(d,p) max min (...) < min max (...) dx & px A general result for any function nex ( min 1 {x=y} ) < min ( mars 1 {x=y} ) yeso,13 ( xeso,13 ( xeso,13 ( yeso,13 ) Some conditions: d'=p\* for optimal margin classifier. dual problem how better character. let f be conses (Hessian H20) suppose hi is affine (hick)=a; wtb) and suppose g; is (strictly) feasible. (In st. H; g:(w) <0) Then I w\*, xx, B\* st. w\* solve primal d. px solve dual, and px=dx= L(wx, xx, px, DW L(w\*, x\*, B\*)=0 KKT complementary condition x; + g;(w\*) =0 /€ 3B S (W\*, 2\*, B\*) =0 KKT: short for g: (w+) =0 Karush-Kuhn-Tucker X: >20

KKT condition: If x: >0 => g: (w\*)=0 (g: (w) is an active "constraint). Lagrange multiplier: di si in svr problem di : W is SVM problem w.b (notation change Parameters min = 11/12 St. y(1) (wTx(1)+b)≥1, i=1...m. g, (w, b) = -y(i) (wx(i)+b) + 1 ≤0  $d: >0 \Rightarrow d: (w, b) =0$  (active constraint) (=) (x", y") has functional margin ! X: =0 for most non-support Fune. Maying (usually 2 \$0) "Support Vectors"



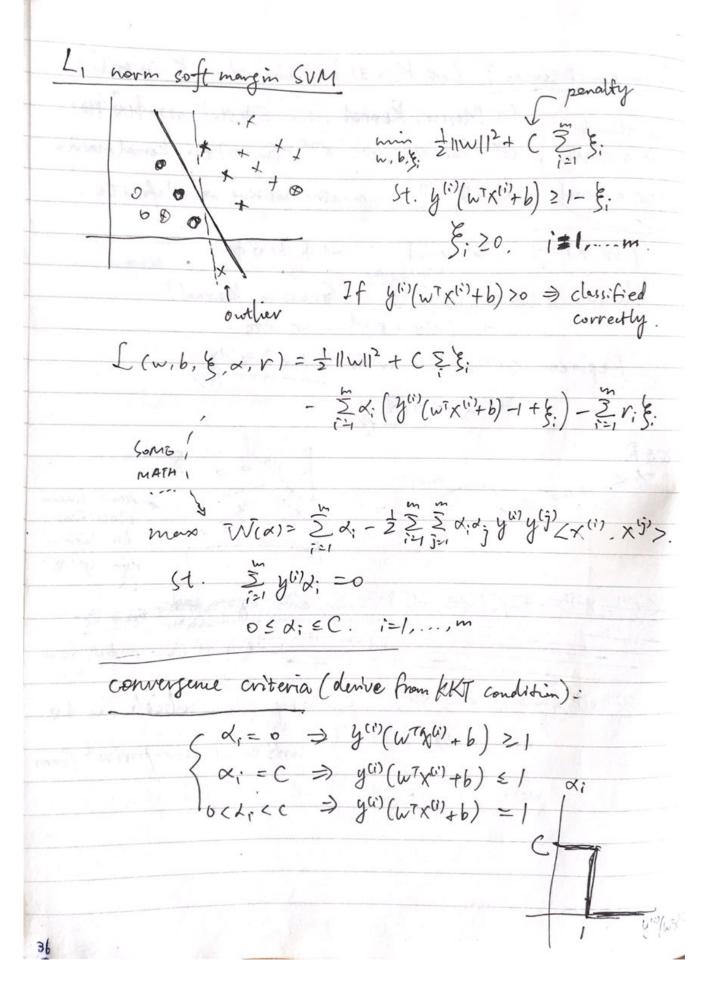
W= Zz; y"x") hw,b(x)= &(WTX+b) 1 threshold Function wix+b= = xxy"(x",x)+b. Kernels X(i) - very ligh dimentional (X(i) C/K ) (X(i), X(J)) efficiently computed Lecture 8 min \$ 11W112 - benels st. yci)(wTx1)+b) >1 - Loft Margin - SMO Dual Problem: man Za: - & TEX: of your st. d; 20 = y;d; =0 w = Zd, y(1)x(1) hw,6(x) = g(wTx+b) = g( \(\frac{2}{2}\)\(\frac{2}{3 Heme X+R living area  $\times \frac{P}{X} \begin{bmatrix} x \\ x^2 \\ x^4 \end{bmatrix} = \phi(x)$ Replace <x(1), x(j)> with <\p(x(1)), \ph(x(1))>

\$ (x) - very high din. cannot compute efficiently K(xi'), x'j') = ( \$\phi(x'i'), \$\phi(x'j') >: sometimes compute efficiently x, & ERn  $K(x_{i}z) = (x_{i}z)^{2} = (\sum_{j=1}^{n} X_{i}z_{j})(\sum_{j=1}^{n} X_{j}z_{j}) = \sum_{j=1}^{n} \sum_{j=1}^{n} (X_{i}X_{j})(z_{i}z_{j})$ = ( \phi(x)) \phi(\partial) X, X, Need O(n2) to compute E(x). D(N)= But need O(n) time to XL XZ Compute K(X, Z) X2 X3 XI XI x3 X2 X3 X3 52 X1 JIX \((x,z)=(x^Tz+6)) Scherally: K(x, z) = (x = +c) -) (n+d) features of all monomials up to degree d. 5 m 0(5) X moder < p(x), 0(2)> X, & similar: < o(x), o(x)> maybe large; - - - dismilar: small.

K(x, 2) - large if x, 2 similar small if x12 dissimilar 3 \$ st. K(x, 2)= < \$(x), \$(2)>? Suppose Kis a kernel let {x1, ..., x(m)} he given let K & Rmxn Kij = K (xw, x'j) Then for any vector 2 & Rm. (21 K & = = = = = (x") (x")) =; = 3 22 2; ( ( ( ( (x ()))) ( ( ( (x ()))) 2; = = = ( = = : p(x13) 20 K ≥ 0. (positive semi-definite) And the inverse Judgement is true

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Theonem (Mercer): Let K(x, 2) be given, then K is walled (is Mercer) Konnel (:. e. 3 p st. K(x,2) = \$(x) \$(2)) Use this to iff. for all {x(1)... x(m)} (mcoo) the Kernel matrix test a function KERMXM is symmetric positive semidefinite. is/not a kernel eg: K(x,x)=-1 not valid . '.' -1 \$ \$ (x) \$ (x) (pare K(x, 5)= 6- 11x-8115 (Gaussian Kernel) or (x72+6)d < x(1), x(1)> with K(x(1), x(1)). xER map dim. space written as inner-product form  $\langle x^{(i)}, x^{(i)} \rangle \rightarrow \langle (x^{(i)}, x^{(i)}) \rangle$ for any also which can be un'tten as inner-product form



Digression:	
coordinate ascent (a	nother opt. question)
	(no constraints on d; s)
Repeat ?	Gold everything except X;
Topan	1 fixed.
for 1=1 to m.	Fold everything except X:
2 di = argman	W(d,di, d: , diffdm)
regard for the forman	source 12th china tout
	W(d, dm)
1/1/0//	W(d, dm)
	Charles Strain Strain Cotton
	30/30/31/36 3 AB
X THE STATE OF THE	and the second of the second of
1 1 1 1 1 1 1	1. ~ Alexander de la
high dum: not pixed	order d, du
henvistic val	he function decide which is chosen to vary (change)
	to converge (compared to Nepoton's agethod)
* I'nner loops	executes very quick.
0	
SMO	Lossella grant pattoriane el vinos x
Coordinate ascent can	not work directly on SVM dual opt. problem
: constraints :	Sy (i) 2: =0.   Smo Algo. is due to
: Change 2 d:	cat a time.
Je z vi	John Platt @ Microsofo
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Select d; d; (heuristic) (\*) key step hold all di's fixed except di di Optimize W(d) wirit di di st. constraints < So do extremely efficiently: Although many iterations (large number) But each iter runs very fast (cheap) Know = 3"d: =0. y (1) of 1 + y (2) die = - 3 y (1) die def Ø € d; € C. W (d. . dz , dz ... dn) = W(5-402, x, x, ... xm) = a 22+ box + C ( ghadratic fune.) optimice quadratic fune. => x = 5-ywd2 in inner loop operates very efficiently. Compute b is not hard. Do it after class.

Applications of SIM @ Handler's Digit Recognition K(x, y) = (x, h) d as 6 \_ 50, 0000 X E Kloo Sway is comparable with best neural networks @ Classify protein seg.'s BAJTSIAJBAJTAU anino acid seg. (A. - 2) \$(x) ≥ ? ( hard problem. ~ \$(x) € \$(204) = R160000 MABA use Dynamic Programming (DP): BATT occur 2 times compute \$ (m) \$ (2). TSIA 3538 Basic (Applicative) part of this course is over. Understanding these algorithms

Lecture 9 | learning theory - Bios Variance - Engineed Risk Minimization - Union Bound/ bloeffding Inequality - Uniform convergence. 80+01X O. + O, x + O, x + O, x + Oq x + "underfit" "overfit" " bias" "Variance" Oo+ Orx+Ozx2 "balanced" horx1=g(00+0,x+".+0,x4)

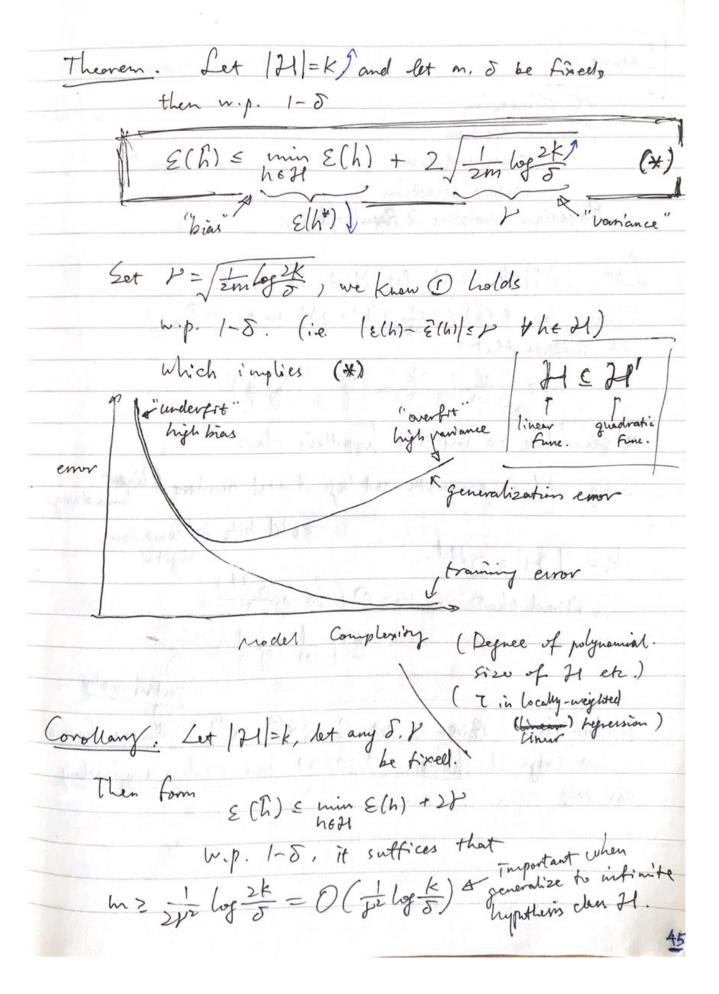
Whear Classification: ho(x)=g(oix) g(x)=1{220} (note yefo,13) S = {(x(i), y(i))} = (training set) (x(i), y(i)) ~ D training ever of ho: E(ho) = Es(ho) = = = 1 Sho(x") + y")} € SVM. logistic negression. ERM: 6 = ang min Es(ho) is approx. of GRM. (GRM is more general) Hypothesis class I = {ho: O ∈ Rn+1}. ho: X -> fo, 1}. GRM: h= argmin & (h) Creneralization error: E(h) = P(x,y)&p(hox) + y). let A. Az ... Ax be k event: Union Bound: ( not necessarily independent) Then P(A,UA,U...UA) = P(A,) + P(A,) + ... + P(A) P(A,UA,UA) < P(A,)+P(A,)+P(Az)

Hoeffding Inequality: Let Z... 2m be ~ I.ID Bernoulli ( $\phi$ ) random variable ( $P(2; =1) = \phi$ ). Let \$ : 1 2 8; and let any 1 > 0 be finel. Then p ( -0 > × ) = 2 e -2 xm distribution of p Central Limit Theorem: works only in is large. The case of finite H H = {h, h2,...hk} K pypothein L= arguin Eg(hi) Strategy: (1) } ≈ E

(2) Show bound on E(h).

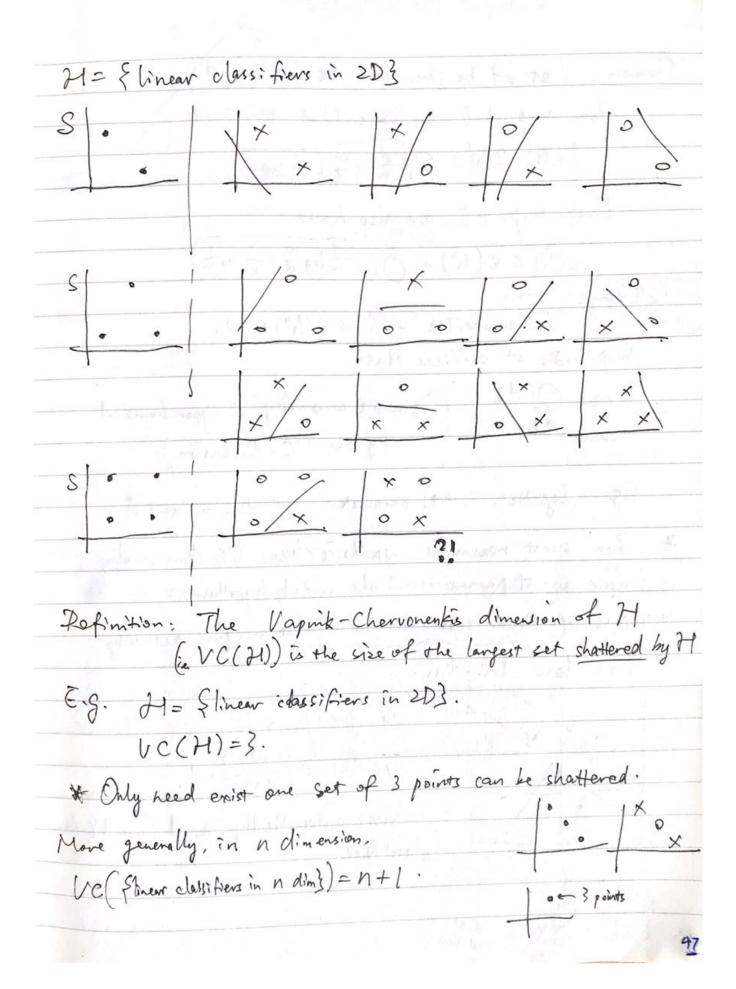
(U Fix any by & H. Define Z:= Ishi(xi) \$ yii) = {0,1} ~ Bernoulli  $P(2:=1) = \epsilon(h_j).$  (All 2: ~ 7.7.D) É(hj) = m = = = = = [ ] { hj(x") + y"} Thean E(hj) By tweffding Zneg: P(( E(hj) - E(hj) = 2 e - 2 m A; = event that \E(h\_1) - E(h\_1) > 2 P(Aj) = 2e 2 m. P( ] h; 6 H st. | E(hj) - ê(hj) /> ). = P(A, UA, U. - · UAx) = = P(A;) = \$ 2e^21/m = 2ke-21/m (1 - both sides). P(\$ h, & H st. | E(h) - E(h) /> L) = P(Vhj & H st. | E(hj)- E(hj) | = 1-2ke-21m So w.p. 1-2Ke-22m, E(h) will be within 2 probability. I of (16) C. Pali 1. of ECh) for all h & H. Uni form Convergence"

(winer I and &, what is m	? He is then my
S = 2Ke-222m, so he for	m.
So long as m ? 2 2 by 2k	, 171 = (12,535)
then w.p. 1-8, we h	one (E(h)-E(h))=> for all het
"Sample complexity" bound.	in practical
Solve for I for fixed in	
w.p. 1-0, we have	that the H.
$ \tilde{\epsilon}(h) - \epsilon(h)  \in \int$	Zim log 2k
(2) Let assume & ho H. 18Ch	1- E(h)/5 r. 0
Can we prove Something a	
L' = ang min ÉCh)	(2) (3)
h* = ang min Elh) he H	3
€ (2) € € (2) + 2	- by D
€ Ê (h*) + ×	- by @
≤ E(h*)+Y+Y	E(h) - traing enor of h
$= \varepsilon(h^*) + 2V.$	Elhi - generalization error
	of h.
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Leture 10  VC dimension  Model Selection  - Cross validation  - Feature selection  Bayesian statistics & Regularisation  Let [H]=k, and let Y, 5 be fixed.  Then for E(h) < min E(h) + 2 Y m.p. 1-5,  it suffices that  mz zir by = O( je log = )  ** generalize to infinite hypothesis class:  Sony H is parameterized by d real numbers ( linear -> 640 bits in computer.  V=  H  = 260d.  Sufficed that mz O( je log = )  Cufficient that mz O( je log = )  Pefinition: Given a set S = fx(i). x(d) I more form we say H shatters S if H can realize any labeling on it.	(1	
Model Solection  - Cross validation  - Feature selection  Bayesian statistics & Regularitation  Let $ \mathcal{H} =k$ , and let $\mathcal{H}$ , $\mathcal{J}$ be fixed.  Then for $\mathcal{E}(h) \leq \min_{k \in \mathcal{H}} \mathcal{E}(h) + 2\mathcal{V}$ w.p. $\mathcal{H}\mathcal{J}$ , it suffices that $m \geq \frac{1}{2} \operatorname{log} \frac{2k}{\mathcal{J}} = \mathcal{O}\left(\frac{1}{2^2} \log \frac{k}{\mathcal{J}}\right)$ * generalize to infinite hypothesis class:  Sony $\mathcal{H}$ is parameterized by direal numbers (linear boundary) $= \mathcal{J}(h) = \mathcal{J}($	Len	me 10
Model Selection  - Cross validation  - feature selection  Bayesian statistics & Regularisation  Let $ \mathcal{H} =k$ , and let $\mathcal{H}$ , $\mathcal{S}$ be fixed.  Then for $\mathcal{E}(h) \leq \min_{\mathcal{E}(h)+2\mathcal{Y}} \sup_{\mathcal{H}} \mathcal{E}(h)$ , it suffices that $m^2 = \lim_{\mathcal{H}} \log \frac{2k}{\mathcal{S}} = \mathcal{O}\left(\frac{1}{\mathcal{H}}\log \frac{k}{\mathcal{S}}\right)$ * generalize to infinite hypothesis class:  Say $\mathcal{H}$ is parameterized by d real numbers ( linear boundar) $= \mathcal{G}(d)$ bits in computer. $k =  \mathcal{H}  = 260 d$ .  Sufficed that: $m \geq \mathcal{O}\left(\frac{1}{\mathcal{H}}\log \frac{1}{\mathcal{S}}\right)$ .  Intuitively:  Definition: Given a set $S = f_{\mathcal{H}}(l)$ . $\mathcal{H}^{(d)}$ ! I more form we say $\mathcal{H}$ shatters $S$ if $\mathcal{H}$ can realize any labeling on it.		
- Cross validation - feature selection  Bayesian statistics & Regularisation  Let $ \mathcal{H} =k$ , and let $\mathcal{H}$ , $\mathcal{S}$ be fixed.  Then for $\mathcal{E}(h) \leq \min_{k \in \mathcal{H}} \mathcal{E}(h) + 2\mathcal{H}$ w.p. $I - \mathcal{S}$ ,  it suffices that $m \geq \frac{1}{2} \operatorname{log} \frac{2k}{\mathcal{S}} = \mathcal{O}\left(\frac{1}{2^2} \log \frac{k}{\mathcal{S}}\right)$ * generalize to infinite hypothesis class:  Say $\mathcal{H}$ is parameterized by direal numbers (linear boundar) $= 640 \text{ bits in Computer}$ . $k =  \mathcal{H}  = 2600 \text{ deg}$ .  Sufficed that: $m \geq \mathcal{O}\left(\frac{1}{2^2} \log \frac{1}{2^2}\right)$ .  Intuitively:  Definition: Given a set $S = f_X(I)$ . $X^{(0)}$ ! I more form we say $\mathcal{H}$ shatters $S$ if $\mathcal{H}$ can realize any labeling on it.		
Soyesian statistics & Regularisation  Let $ \mathcal{H} =k$ , and let $\mathcal{H}$ , $\delta$ be fixed.  Then for $\mathcal{E}(h) \leq \min_{k \in \mathcal{H}} \mathcal{E}(h) + 2\mathcal{V}$ w.p. $I - \delta$ ,  It suffices that $m \geq \frac{1}{2\nu^2} \log \frac{1}{\delta} = O\left(\frac{1}{2^2} \log \frac{1}{\delta}\right)$ * generalize to infinite hypothesis class:  Sony $\mathcal{H}$ is parameterized by direal numbers (linear boundar) $= \delta(4d) \text{ bits in Computer.}$ Oligital $ \mathcal{H}  = 2\delta(d)$ .  Cuffied that: $m \geq O\left(\frac{1}{2^2} \log \frac{1}{\delta}\right)$ .  Intuitively:  Definition: Given a set $S = f_X(l)$ . $\mathcal{H}^{(d)}$ ! In one form we say $\mathcal{H}$ shatters $S$ if $\mathcal{H}$ can realize any labelity on it.	¥)_	
Let $ \mathcal{H} =k$ , and let $\mathcal{H}$ , $\delta$ be fixed.  Then for $\mathcal{E}(h) \leq \min_{h \in \mathcal{H}} \mathcal{E}(h)+2\mathcal{Y} \text{ w.p. } I-\delta$ ,  it suffices that $m \geq \frac{2}{2} \operatorname{pr} \log \frac{2k}{\delta} = O\left(\frac{1}{2^2} \log \frac{k}{\delta}\right)$ * generalize to infinite hypothesis class:  Say $\mathcal{H}$ is parameterized by $d$ real numbers $\begin{pmatrix} \lim_{h \to \infty} \int \frac{k}{k} dh \end{pmatrix}$ $k =  \mathcal{H}  = 2^{6kd}$ .  Sufficed that: $m \geq O\left(\frac{1}{2^2} \log \frac{k}{\delta}\right)$ .  Thurstooly  Definition: Given a set $S = f_X(i)$ . $f_X^{(i)}$ ! I have form we say $\mathcal{H}$ shatters $S$ if $\mathcal{H}$ can realize any labeling on it.	THE STATE OF THE S	- Feature selection
Let $ \mathcal{H} =k$ , and let $\mathcal{H}$ , $\delta$ be fixed.  Then for $\mathcal{E}(h) \leq \min_{h \in \mathcal{H}} \mathcal{E}(h)+2\mathcal{Y} \text{ w.p. } I-\delta$ ,  it suffices that $m \geq \frac{2}{2} \operatorname{pr} \log \frac{2k}{\delta} = O\left(\frac{1}{2^2} \log \frac{k}{\delta}\right)$ * generalize to infinite hypothesis class:  Say $\mathcal{H}$ is parameterized by $d$ real numbers $\begin{pmatrix} \lim_{h \to \infty} \int \frac{k}{k} dh \end{pmatrix}$ $k =  \mathcal{H}  = 2^{6kd}$ .  Sufficed that: $m \geq O\left(\frac{1}{2^2} \log \frac{k}{\delta}\right)$ .  Thurstooly  Definition: Given a set $S = f_X(i)$ . $f_X^{(i)}$ ! I have form we say $\mathcal{H}$ shatters $S$ if $\mathcal{H}$ can realize any labeling on it.		Bayesian Statistics & Regularisation
m 2 2pr log \$\frac{2k}{3} = O(\frac{1}{3^2}\log \frac{k}{3})\$  ** generalize to infinite hypothesis class:  Say H is parameterized by d real numbers (linear boundar) \$\frac{-}{64d}\$ bits in computer.  \$\frac{-}{564d}\$ bits in computer.  \$\frac{-}{504d}\$	ſ	
m 2 2pr log \$\frac{2k}{3} = O(\frac{1}{3^2}\log \frac{k}{3})\$  ** generalize to infinite hypothesis class:  Say H is parameterized by d real numbers (linear boundar) \$\frac{-}{64d}\$ bits in computer.  \$\frac{-}{564d}\$ bits in computer.  \$\frac{-}{504d}\$	2	t   H = k, and let P, J be fixed.
m 2 2pr log \$\frac{2k}{3} = O(\frac{1}{3^2}\log \frac{k}{3})\$  ** generalize to infinite hypothesis class:  Say H is parameterized by d real numbers (linear boundar) \$\frac{-}{64d}\$ bits in computer.  \$\frac{-}{564d}\$ bits in computer.  \$\frac{-}{504d}\$	(	hen for E(h) < min E(h) + 22 w.p. 1-5,
The special state of the special states any labeling on it.  **Seneralize to infinite hypothesis class:  Song H is parameterized by d real numbers ( linear boundar) $\rightarrow 6400$ bits in computer. $\Rightarrow 6400$ bits in computer. $\Rightarrow 6400$ bits in computer.  Olighted  Sufficed that: $m \ge O(\frac{1}{2^{2}}\log \frac{1}{2^{2}})$ .  Intrinitively  Petinitian: Given a set $S = \{x^{(j)}, x^{(d)}\}$ there form  we say H shatters $S = \{x^{(j)}, x^{(d)}\}$ there form  on it.	ì	t suffices that next
* generalize to infinite hypothesis class:  Say It is parameterized by d real numbers ( linear boundar)  ->64d bits in computar.    K =   H   = 264d.   Oligitar    Sufficient that:   m = 0 ( for log 264d)		
Say H is parameterized by d real numbers ( linear boundary $\Rightarrow$ 64d bits in computer. $\Rightarrow$ 64d bits in computer. Objected Sufficient that: $m \ge O\left(\frac{1}{p^2}\log\frac{56kd}{5}\right)$ $= O\left(\frac{d}{p^2}\log\frac{1}{5}\right).$ Intuitively Definition: Given a set $S = \{x^{(i)}, x^{(d)}\}$ knone form we say H shatters $S$ if $H$ can realize any labeling on it.	1266	The state of the s
Say H is parameterized by d real numbers ( linear boundary $\Rightarrow$ 64d bits in computer. $\Rightarrow$ 64d bits in computer. Objected Sufficient that: $m \ge O\left(\frac{1}{p^2}\log\frac{1}{p^2}\right)$ $= O\left(\frac{d}{p^2}\log\frac{1}{p^2}\right).$ Intuitively Definition: Given a set $S = \{x^{(i)}, x^{(d)}\}$ knone form we say H shatters $S$ if $H$ can realize any labeling on it.	×	generalize vo infinite hypothesis class:
$K =  \mathcal{H}  = 260d$ .		and It is personaterized by a seal number ( linear
Sufficed that: $m \ge O(\sqrt{y^2 \log 5})$ = $O(\sqrt{y^2 \log 5})$ .  Intuitively  Definition: Given a set $S = \{x^{(i)}, x^{(d)}\}$ . Imone form  we say $\mathcal{H}$ shatters $S$ if $\mathcal{H}$ can realize any labeling on it.	٥	boundar
Sufficed that: $m \ge O(\sqrt{p^2 \log 5})$ = $O(\sqrt{p^2 \log 5})$ .  Intuitively  Definition: Given a set $S = \{x^{(i)}, x^{(d)}\}$ . Imone form  we say $\mathcal{H}$ shatters $S$ if $\mathcal{H}$ can realize any labeling on it.		-> 64d bits in computer.
Sufficient that: m=0(pr log 5)  = 0(pr log 5).  Intuitively  Definition: Given a set S = fx(1). x(d). I more form  we say H shatters S if H can realize any labeling on it.		$\langle =  \mathcal{H}  = 2^{040}$ .
$= \mathcal{O}\left(\frac{d}{pr}\log 5\right).$ Intuitively  Definition: Given a set $S = \{x^{(i)}, x^{(d)}\}$ knone form  we say $\mathcal{H}$ shatters $S$ if $\mathcal{H}$ can realize any labeling on it.		Sufficed that: mzO(jr log)
Definition: Given a set $S = \{x^{(i)}, x^{(d)}\}$ more form we say $\mathcal{H}$ shatters $S$ if $\mathcal{H}$ can realize any labeling on it.		$= O(\frac{d}{d} \log 1)$
Definition: Given a set $S = \{x^{(i)}, x^{(d)}\}$ more form we say $\mathcal{H}$ shatters $S$ if $\mathcal{H}$ can realize any labeling on it.		( Dr 102 2).
we say It shatters S if It can realize any labeling on it.	_	
we say It shatters S if It can realize any labeling on it.	P	efinition: Given a set S= {x(1). x(1)} more form
on it.		we say I shatters S if I can realize any labeling
	01	
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Theorem. Let I be given and lot VC(II)=d. then w-p. ho, we have that the H: 1 2(h) - 2(h) 5 0 alog m + in log 1). Thus, w.p. 1-5, are also have E(h) = E(h) + O( Imlog m + inlog o) Collollary: To gharantee E(h) s E(h\*)+2), W.p. L-S, it suffices that m= O(d) ("sample complexity" is upper bounded by the VC-dimension) Eg. logistic. het parameter Ve din = n+1. It For most reasonable hypothesis classes. VC. ohim usually linear in It parameters of the model/hypotheris \* Crass of linear separaters with large margin actually has low Ve-dim. 24 11x(i)1/2 5 R Ve (21) 5 [ +2]+1. SUM automatically find low Verdin classifier. 11 x112 = \$ K:3 1 X 1/2 = X; (converge condition)