

Functional Programming	WS 2018/2019	LVA 703024+703025
Exercises 5		November 30, 2018

All Solutions

2 We take boolToInt = λb ite b 1 0 and compute the normal form of boolToInt True as follows:

$$\begin{array}{l} \mathsf{boolToInt} \; \mathsf{True} = \underline{(\lambda b. \, \mathsf{ite} \; b \; 1 \; 0) \; \mathsf{True}} \\ \to_{\beta} \; \mathsf{ite} \; \mathsf{True} \; 1 \; 0 = \underline{(\lambda x y z. \, x \; y \; z) \; \mathsf{True}} \; 1 \; 0 \\ \to_{\beta} \underline{(\lambda y z. \, \mathsf{True} \; y \; z) \; 1} \; 0 \\ \to_{\beta} \underline{(\lambda z. \, \mathsf{True} \; 1 \; z) \; 0} \\ \to_{\beta} \mathsf{True} \; 1 \; 0 = \underline{(\lambda x y. \, x) \; 1 \; 0} \\ \to_{\beta} (\lambda y. \, 1) \; 0 \\ \to_{\beta} 1 \end{array}$$

3 Since E = A A, we want the following property to hold

$$A A t \rightarrow_{\beta}^{*} A A.$$

This tells us that A is a "function" that takes two arguments (A and t), duplicates the first one (A) and ignores the second one (t). Thus a possible solution is $A = \lambda ax$. a, which is shown by the following β -reduction:

$$\begin{split} E \ t &= A \ A \ t \\ &= \underbrace{(\lambda ax. \ a \ a) \ A}_{\beta} \ t \\ &\to_{\beta} \underbrace{(\lambda x. \ A \ A) \ t}_{A \ A = E} \end{split}$$