

Value Iteration Convergence

Review

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- How do we reason about the **future consequences** of actions in an MDP?

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- How do we reason about the **future consequences** of actions in an MDP?
- What are the basic **algorithms for solving MDPs**?

Guiding Questions

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- Does value iteration always converge?
- Is the value function unique?
- Can there be multiple optimal policies?
- Is there always a deterministic optimal policy?

Value Iteration: The Bellman Operator

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Algorithm: Value Iteration

while $\|V - V'\|_\infty > \epsilon$

$V \leftarrow V'$

$V' \leftarrow B[V]$

return V'

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$$B[V](s) = \max_{a \in A} (R(s, a) + \gamma E [V(s')])$$

Value Iteration Convergence

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Theorem 1: Let $\{V_1, \dots, V_\infty\}$ be a sequence of value functions for a discrete MDP generated by the recurrence $V_{k+1} = B[V_k]$. If $\gamma < 1$, then $\lim_{k \rightarrow \infty} V_k = V^*$.

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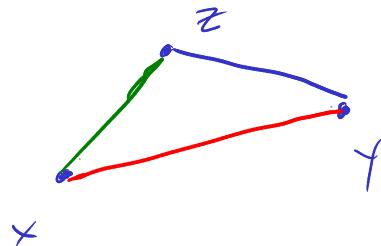
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Example : $M = \mathbb{R}^2$ $d(x, y) = \|x - y\|_2$

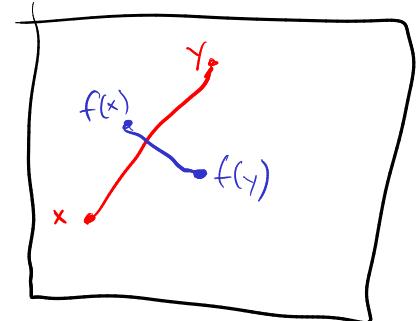
Contraction Mappings

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Definition: A *contraction mapping* on metric space (M, d) is a function $f : M \rightarrow M$ satisfying

$$d(f(x), f(y)) \leq \alpha d(x, y)$$

for some α , $0 \leq \alpha \leq 1$ and all x and y in M .



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$$f(x) = \begin{bmatrix} \frac{x[1]}{2} + 1 \\ \frac{x[2]}{2} + \frac{1}{2} \end{bmatrix}$$

$$\begin{aligned} M &= \mathbb{R}^2 \\ d(x, y) &= \|x - y\|_2 \end{aligned}$$

Script: contraction_mapping.jl

$$\begin{aligned} d(f(x), f(y)) &= \sqrt{\left(\frac{x[1]}{2} + 1 - \frac{y[1]}{2} - 1\right)^2 + \left(\frac{x[2]}{2} + \frac{1}{2} - \frac{y[2]}{2} - \frac{1}{2}\right)^2} \\ &= \frac{1}{2} \sqrt{(x[1] - y[1])^2 + (x[2] - y[2])^2} \\ &= \frac{1}{2} d(x, y) \end{aligned}$$

f is a contraction mapping with $\alpha = \frac{1}{2}$

Banach's Theorem

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By Lemma 2 and Banach's theorem (part 2), repeated application of the Bellman operator always has a fixed point limit, \hat{V} .

By Banach's theorem (part 1), $\hat{V} = B[\hat{V}]$. Since \hat{V} satisfies Bellman's equation, it is optimal and $\hat{V} = V^*$. *optimality*

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3. There are a finite number of possible policies

Does Policy Iteration Converge?

Theorem: Policy iteration converges to an optimal policy for a finite MDP in finite time.

Proof (sketch):

1. The policy will either improve or stay the same at each iteration
2. The policy will stay the same if and only if $V^\pi = V^*$
3. There are a finite number of possible policies
4. By (1), (2), and (3), the policy will improve until it finds the optimal policy, and it will always find the optimal policy.

Properties of optimal MDP solutions

- Does every MDP have a unique optimal value function, V^* ? Yes
- Does every MDP have a unique optimal policy, π^* ? No
- Does every MDP have a *deterministic* optimal policy? Yes $\pi(a|s)$

Because
of
Banach's Theorem

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Justification

- Suppose that $\tilde{\pi}$ is optimal and, for some s , $\tilde{\pi}(a^1 | s) > 0$, $\tilde{\pi}(a^2 | s) > 0$, and $\tilde{\pi}(a^1 | s) + \tilde{\pi}(a^2 | s) = 1$.
- Then $\underline{Q^*(s, a^1)} = \underline{Q^*(s, a^2)} = \underline{V^*(s)}$. If this were not true, then $\tilde{\pi}$ would not be optimal.
- As a consequence, a deterministic policy $\tilde{\pi}'$ with $\tilde{\pi}'(s) = a^1$ is also optimal!

Guiding Questions

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- Does value iteration always converge?
- Is the value function unique?
- Can there be multiple optimal policies?
- Is there always a deterministic optimal policy?

Break

Conservation MDP

