



DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

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Experiment 1.4

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Subject Name: Advanced programming

Subject Code: 22ITP-351

Lab II

PROBLEM 1:

1. **Aim:** Longest Nice Substring (**Easy**)
2. **Objective:** Given a string *s*, return the longest substring of *s* that is nice. If there are multiple, return the substring of the earliest occurrence. If there are none, return an empty string.

3. **Code:**

```
class Solution {
    public String longestNiceSubstring(String s) {
        if (s.length() < 2) return "";
        for (int i = 0; i < s.length(); i++) {
            char c = s.charAt(i);
            if (s.contains(Character.toString(Character.toUpperCase(c))) && s.contains (Character.toString
(Character.toLowerCase(c)))) {
                continue;
            }
            String left = longestNiceSubstring(s.substring(0, i));
            String right = longestNiceSubstring(s.substring(i + 1));
            return left.length() >= right.length() ? left : right;
        }
        return s;
    }
}
```

4. **Time Complexity:**

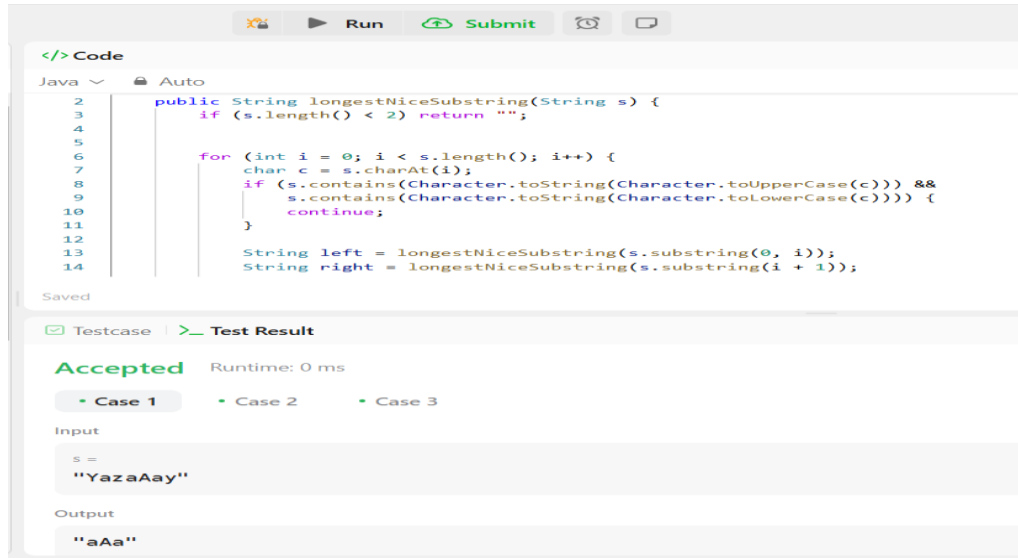
Best Case (String is already "nice") = $O(n)$

Average Case (Some splits, but balanced) = $O(n \log n)$

Worst Case (Unbalanced splits at every character) = $O(n^2)$

Space complexity is $O(n)$

5. **Output:**



```
</> Code
Java v Auto
2 public String longestNiceSubstring(String s) {
3     if (s.length() < 2) return "";
4
5
6     for (int i = 0; i < s.length(); i++) {
7         char c = s.charAt(i);
8         if (s.contains(Character.toString(Character.toUpperCase(c))) &&
9             s.contains(Character.toString(Character.toLowerCase(c)))) {
10            continue;
11        }
12        String left = longestNiceSubstring(s.substring(0, i));
13        String right = longestNiceSubstring(s.substring(i + 1));
14    }
15
16    return "";
17 }
Saved
Testcase Test Result
Accepted Runtime: 0 ms
Case 1 Case 2 Case 3
Input
s =
"YazaAay"
Output
"aAa"
```

PROBLEM 2:

1. **Aim:** Maximum Subarray (**Medium**).
2. **Objective:** Given an integer array `nums`, find the subarray with the largest sum, and return its sum.

3. Code:

```
class Solution {
    public int maxSubArray(int[] nums) {
        int maxSum = nums[0];
        int currentSum = nums[0];

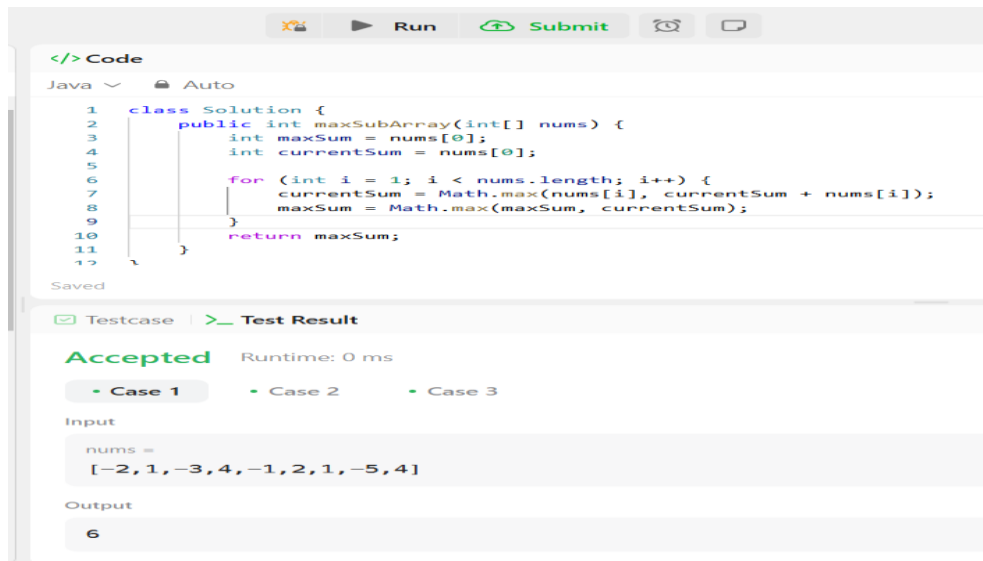
        for (int i = 1; i < nums.length; i++) {
            currentSum = Math.max(nums[i], currentSum + nums[i]);
            maxSum = Math.max(maxSum, currentSum);
        }
        return maxSum;
    }
}
```

4. Time Complexity:

Time Complexity: $O(n)$ (linear time)

Space Complexity: $O(1)$ (constant space)

5. Output:



```
</> Code
Java Auto
1 class Solution {
2     public int maxSubArray(int[] nums) {
3         int maxSum = nums[0];
4         int currentSum = nums[0];
5
6         for (int i = 1; i < nums.length; i++) {
7             currentSum = Math.max(nums[i], currentSum + nums[i]);
8             maxSum = Math.max(maxSum, currentSum);
9         }
10        return maxSum;
11    }
12}
Saved
Testcase Test Result
Accepted Runtime: 0 ms
Case 1 Case 2 Case 3
Input
nums =
[-2, 1, -3, 4, -1, 2, 1, -5, 4]
Output
6
```

PROBLEM 3:

1. **Aim:** Reverse Pairs (Hard).

2. **Objective:** Given an integer array nums, return the number of reverse pairs in the array.
A reverse pair is a pair (i, j) where:

- $0 \leq i < j < \text{nums.length}$ and
- $\text{nums}[i] > 2 * \text{nums}[j]$.

3. **Code:**

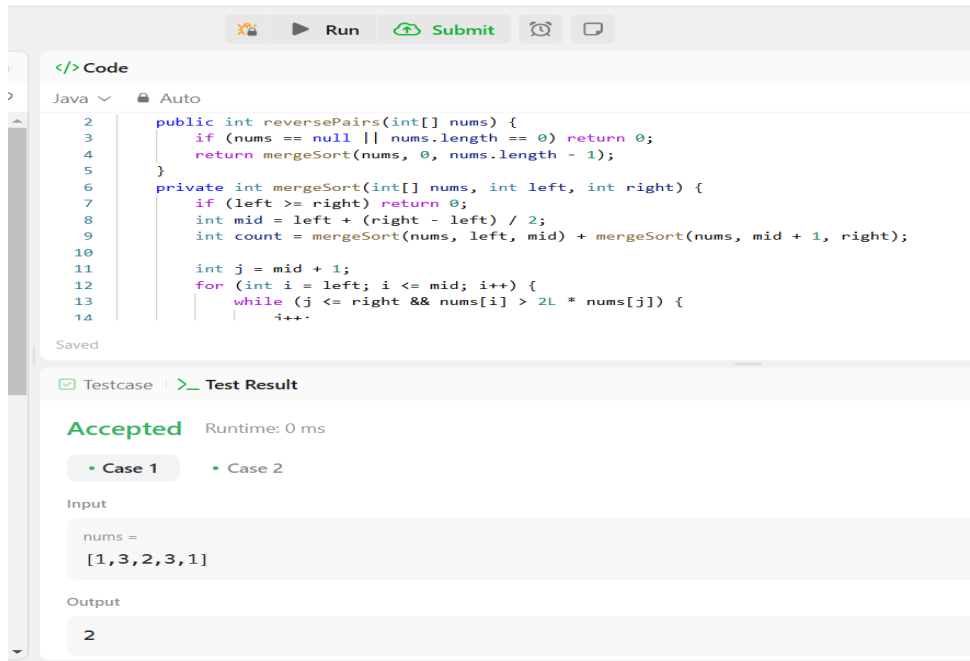
```
class Solution {
    public int reversePairs(int[] nums) {
        if (nums == null || nums.length == 0) return 0;
        return mergeSort(nums, 0, nums.length - 1);
    }
    private int mergeSort(int[] nums, int left, int right) {
        if (left >= right) return 0;
        int mid = left + (right - left) / 2;
        int count = mergeSort(nums, left, mid) + mergeSort(nums, mid + 1, right);
        int j = mid + 1;
        for (int i = left; i <= mid; i++) {
            while (j <= right && nums[i] > 2L * nums[j]) {
                j++;
            }
            count += j - (mid + 1);
        }
    }
}
```

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```
        merge(nums, left, mid, right);
        return count;
    }

    private void merge(int[] nums, int left, int mid, int right) {
        int[] temp = new int[right - left + 1];
        int i = left, j = mid + 1, k = 0;
        while (i <= mid && j <= right) {
            if (nums[i] <= nums[j]) {
                temp[k++] = nums[i++];
            } else {
                temp[k++] = nums[j++];
            }
        }
        while (i <= mid) temp[k++] = nums[i++];
        while (j <= right) temp[k++] = nums[j++];
        System.arraycopy(temp, 0, nums, left, temp.length);
    }
}
```

4. Output:



The screenshot shows a Java IDE with the following code:

```
2 public int reversePairs(int[] nums) {
3     if (nums == null || nums.length == 0) return 0;
4     return mergeSort(nums, 0, nums.length - 1);
5 }
6 private int mergeSort(int[] nums, int left, int right) {
7     if (left >= right) return 0;
8     int mid = left + (right - left) / 2;
9     int count = mergeSort(nums, left, mid) + mergeSort(nums, mid + 1, right);
10
11     int j = mid + 1;
12     for (int i = left; i <= mid; i++) {
13         while (j <= right && nums[i] > 2L * nums[j]) {
14             j++;
15         }
16     }
17     return count;
18 }
```

The test result shows:

Accepted Runtime: 0 ms

Case 1

Input: nums = [1,3,2,3,1]

Output: 2

5. Time Complexity:

Time Complexity = $O(n \log n)$

Space Complexity = $O(n)$

PROBLEM 4:

1. **Aim:** Longest Increasing Subsequence II (**Hard**).
2. **Objective:** You are given an integer array `nums` and an integer `k`. Find the longest subsequence of `nums` that meets the following requirements:
 - The subsequence is strictly increasing and
 - The difference between adjacent elements in the subsequence is at most `k`.Return the length of the longest subsequence that meets the requirements.

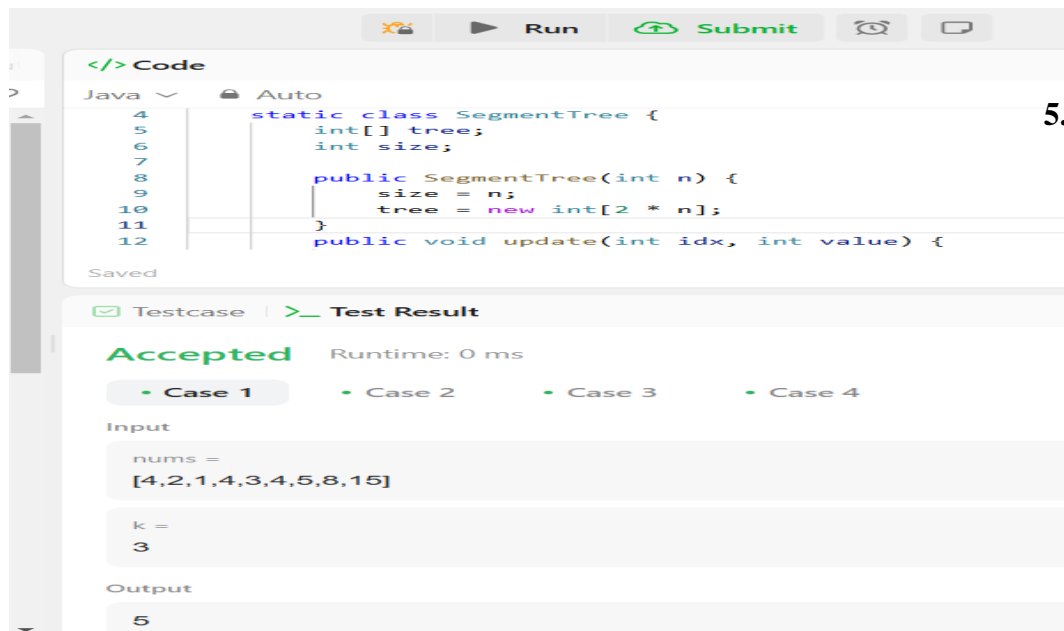
3. **Code:**

```
import java.util.*;
class Solution {
    static class SegmentTree {
        int[] tree;
        int size;
        public SegmentTree(int n) {
            size = n;
            tree = new int[2 * n];
        }
        public void update(int idx, int value) {
            idx += size;
            tree[idx] = value;
            while (idx > 1) {
                idx /= 2;
                tree[idx] = Math.max(tree[2 * idx], tree[2 * idx + 1]);
            }
        }
        public int query(int left, int right) {
            int res = 0;
            left += size;
            right += size;
            while (left <= right) {
                if ((left & 1) == 1) res = Math.max(res, tree[left++]);
                if ((right & 1) == 0) res = Math.max(res, tree[right--]);
                left /= 2;
                right /= 2;
            }
            return res;
        }
    }
    public int lengthOfLIS(int[] nums, int k) {
```

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```
int maxVal = Arrays.stream(nums).max().getAsInt();
SegmentTree segTree = new SegmentTree(maxVal + 1);
int maxLen = 1;
for (int num : nums) {
    int bestPrev = segTree.query(Math.max(1, num - k), num - 1);
    int newLength = bestPrev + 1;
    segTree.update(num, newLength);
    maxLen = Math.max(maxLen, newLength);
}
return maxLen;
}
```

4. Output:



The screenshot shows a code editor with the following Java code for a Segment Tree:

```
5.
4 static class SegmentTree {
5     int[] tree;
6     int size;
7
8     public SegmentTree(int n) {
9         size = n;
10        tree = new int[2 * n];
11    }
12    public void update(int idx, int value) {
```

Below the code editor, the test results are shown:

Accepted Runtime: 0 ms

• Case 1 • Case 2 • Case 3 • Case 4

Input

nums =
[4,2,1,4,3,4,5,8,15]

k =
3

Output

5

5. Time Complexity:

Time Complexity = $O(n \log n)$

Space Complexity = $O(n)$

6. Learning Outcome:

- Learned how different problems have varying complexities, ranging from $O(n)$ (Kadane's Algorithm) to $O(n \log n)$ (Merge Sort & Segment Tree).
- Explored recursion-based solutions (Longest Nice Substring) and Divide & Conquer techniques (Reverse Pairs using Merge Sort).
- Implemented Segment Tree for optimized Longest Increasing Subsequence II, reducing time complexity to $O(n \log n)$.
- Applied Kadane's Algorithm for Maximum Subarray, Merge Sort-based counting for Reverse Pairs, and Segment Tree-based LIS for efficient computations.