



DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

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Experiment 4

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Subject Name: Advanced programming

Lab II

UID: 22BET10188

Section/Group: 22BET_IOT-703/A

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Subject Code: 22ITP-351

PROBLEM 1:

1. **Aim:** Longest Nice Substring (**Easy**)
2. **Objective:** Given a string s , return the longest substring of s that is nice. If there are multiple, return the substring of the earliest occurrence. If there are none, return an empty string.

3. **Code:**

```
class Solution {
    public String longestNiceSubstring(String s)
    { if (s.length() < 2) return "";
      for (int i = 0; i < s.length(); i++)
      { char c = s.charAt(i);
        if (s.contains(Character.toString(Character.toUpperCase(c))) && s.contains (Character.toString
(Character.toLowerCase(c)))) {
            continue;
        }
        String left = longestNiceSubstring(s.substring(0, i));
        String right = longestNiceSubstring(s.substring(i + 1));
        return left.length() >= right.length() ? left : right;
      }
      return s;
    }
}
```

4. **Time Complexity:**

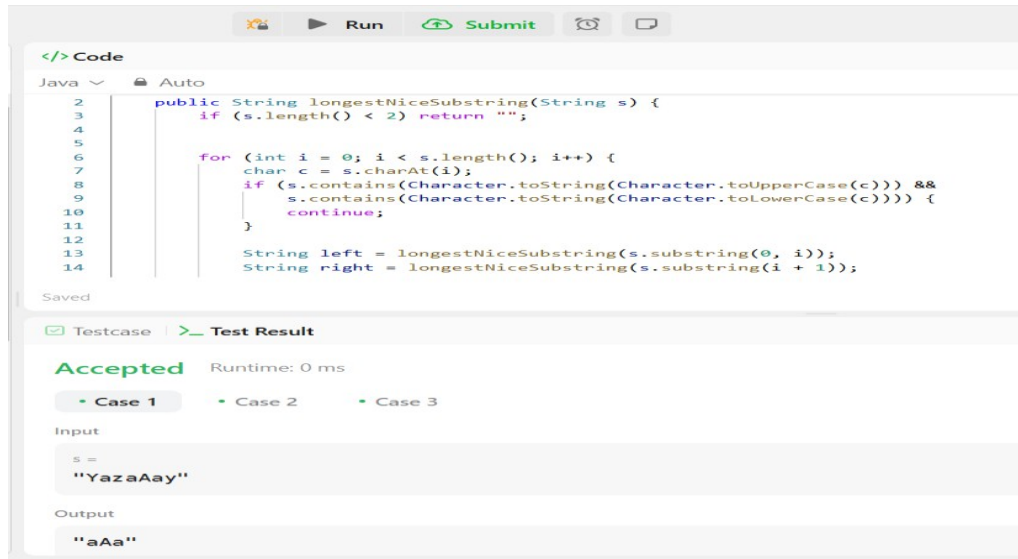
Best Case (String is already "nice") = $O(n)$

Average Case (Some splits, but balanced) = $O(n \log n)$

Worst Case (Unbalanced splits at every character) = $O(n^2)$

Space complexity is $O(n)$

5. **Output:**



```
</> Code
Java ▾ Auto
2 public String longestNiceSubstring(String s) {
3     if (s.length() < 2) return "";
4
5
6     for (int i = 0; i < s.length(); i++) {
7         char c = s.charAt(i);
8         if (s.contains(Character.toString(Character.toUpperCase(c))) &&
9             s.contains(Character.toString(Character.toLowerCase(c)))) {
10             continue;
11         }
12
13         String left = longestNiceSubstring(s.substring(0, i));
14         String right = longestNiceSubstring(s.substring(i + 1));
15     }
16 }
Saved
Testcase | Test Result
Accepted Runtime: 0 ms
• Case 1 • Case 2 • Case 3
Input
s =
"YazaAay"
Output
"aAa"
```

PROBLEM 2:

1. **Aim:** Maximum Subarray (**Medium**).
2. **Objective:** Given an integer array nums, find the subarray with the largest sum, and return its sum.
3. **Code:**

```
class Solution {
    public int maxSubArray(int[] nums)
    { int maxSum = nums[0];
      int currentSum = nums[0];

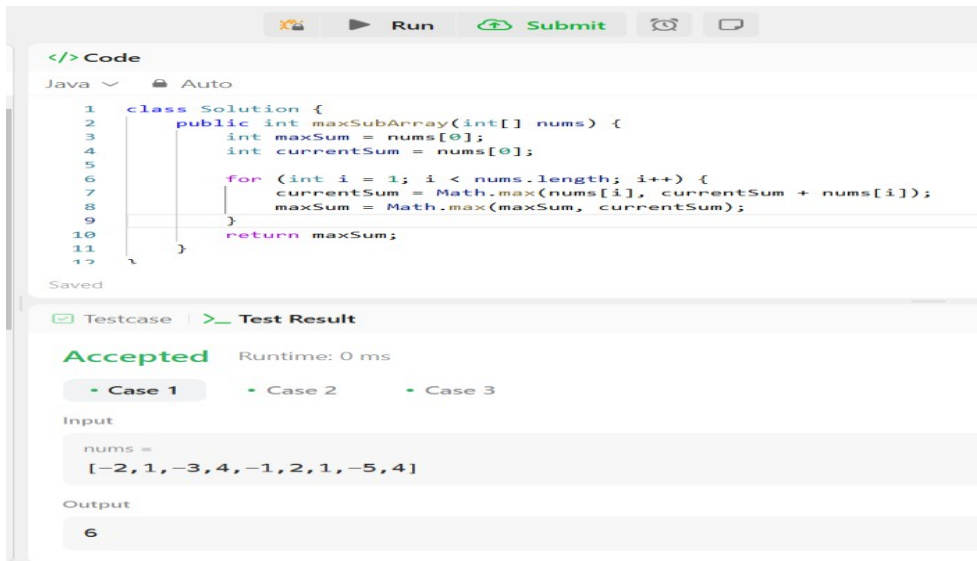
      for (int i = 1; i < nums.length; i++) {
          currentSum = Math.max(nums[i], currentSum + nums[i]);
          maxSum = Math.max(maxSum, currentSum);
      }
      return maxSum;
    }
}
```

4. **Time Complexity:**

Time Complexity: $O(n)$ (linear time)

Space Complexity: $O(1)$ (constant space)

5. **Output:**



```

1  class Solution {
2      public int maxSubArray(int[] nums) {
3          int maxSum = nums[0];
4          int currentSum = nums[0];
5
6          for (int i = 1; i < nums.length; i++) {
7              currentSum = Math.max(nums[i], currentSum + nums[i]);
8              maxSum = Math.max(maxSum, currentSum);
9          }
10         return maxSum;
11     }
12 }

```

Testcase | Test Result

Accepted Runtime: 0 ms

Case 1 Case 2 Case 3

Input

nums =
[-2, 1, -3, 4, -1, 2, 1, -5, 4]

Output

6

PROBLEM 3:

1. Aim: Reverse Pairs (Hard).

2. Objective: Given an integer array nums, return the number of reverse pairs in the array. A reverse pair is a pair (i, j) where:

- $0 \leq i < j < \text{nums.length}$ and
- $\text{nums}[i] > 2 * \text{nums}[j]$.

3. Code:

```

class Solution {
    public int reversePairs(int[] nums) {
        if (nums == null || nums.length == 0) return 0;
        return mergeSort(nums, 0, nums.length - 1);
    }
    private int mergeSort(int[] nums, int left, int right)
    { if (left >= right) return 0;
      int mid = left + (right - left) / 2;
      int count = mergeSort(nums, left, mid) + mergeSort(nums, mid + 1, right);
      int j = mid + 1;
      for (int i = left; i <= mid; i++) {
          while (j <= right && nums[i] > 2L * nums[j]) { j++;
              +;
          }
          count += j - (mid + 1);
      }
    }
}

```

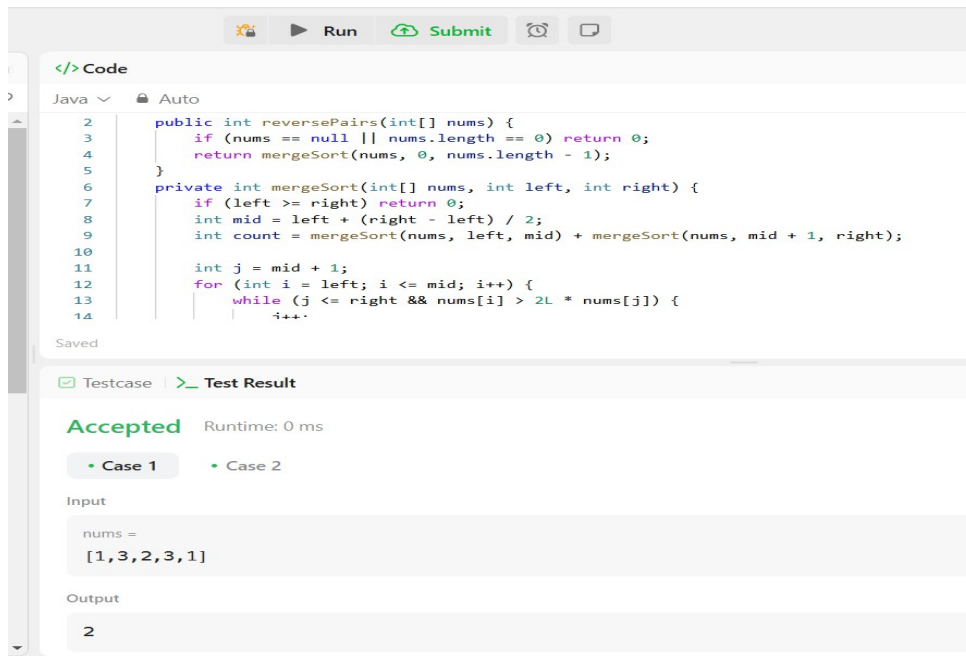
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```

merge(nums, left, mid, right);
return count;
}
private void merge(int[] nums, int left, int mid, int right)
{
    int[] temp = new int[right - left + 1];
    int i = left, j = mid + 1, k = 0;
    while (i <= mid && j <= right) {
        if (nums[i] <= nums[j]) { temp[k++] = nums[i++]; }
        else {
            temp[k++] = nums[j++];
        }
    }
    while (i <= mid) temp[k++] = nums[i++];
    while (j <= right) temp[k++] = nums[j++];
    System.arraycopy(temp, 0, nums, left, temp.length);
}
}

```

4. Output:



```

</> Code
Java Auto
2 public int reversePairs(int[] nums) {
3     if (nums == null || nums.length == 0) return 0;
4     return mergeSort(nums, 0, nums.length - 1);
5 }
6 private int mergeSort(int[] nums, int left, int right) {
7     if (left >= right) return 0;
8     int mid = left + (right - left) / 2;
9     int count = mergeSort(nums, left, mid) + mergeSort(nums, mid + 1, right);
10
11     int j = mid + 1;
12     for (int i = left; i <= mid; i++) {
13         while (j <= right && nums[i] > 2L * nums[j]) {
14             j++;
15         }
16     }
17 }

```

Saved

Testcase Test Result

Accepted Runtime: 0 ms

Case 1 Case 2

Input

nums =
[1, 3, 2, 3, 1]

Output

2

5. Time Complexity:

Time Complexity = $O(n \log n)$

Space Complexity = $O(n)$



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PROBLEM 4:

1. **Aim:** Longest Increasing Subsequence II (**Hard**).
2. **Objective:** You are given an integer array `nums` and an integer `k`. Find the longest subsequence of `nums` that meets the following requirements:
 - The subsequence is strictly increasing and
 - The difference between adjacent elements in the subsequence is at most `k`.Return the length of the longest subsequence that meets the requirements.

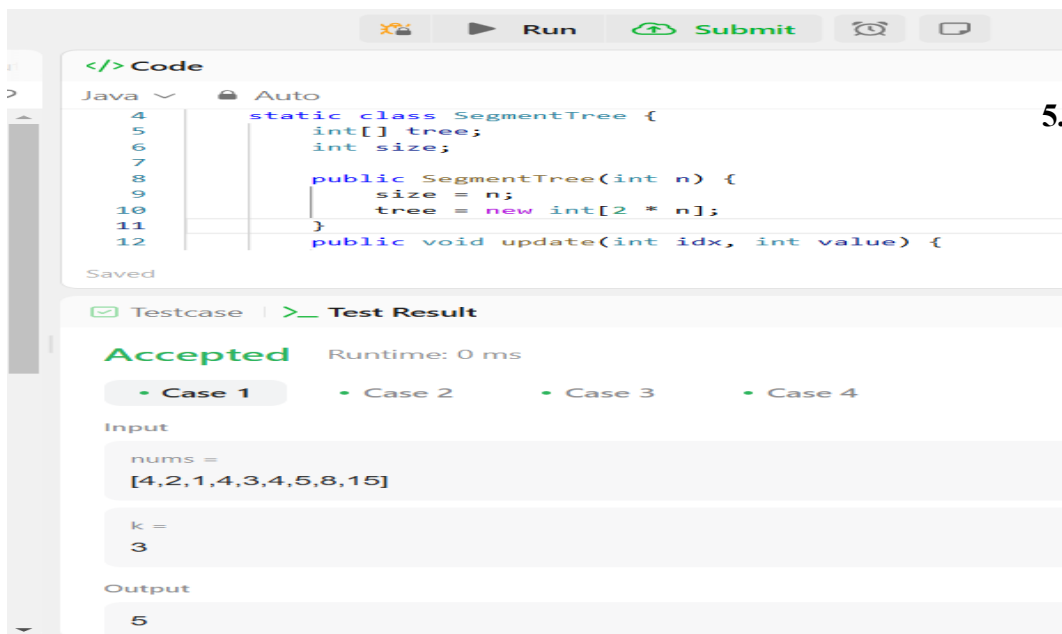
3. **Code:**

```
import java.util.*;
class Solution {
    static class SegmentTree
    {
        int[] tree;
        int size;
        public SegmentTree(int n)
        {
            size = n;
            tree = new int[2 * n];
        }
        public void update(int idx, int value)
        {
            idx += size;
            tree[idx] = value;
            while (idx > 1) {
                idx /= 2;
                tree[idx] = Math.max(tree[2 * idx], tree[2 * idx + 1]);
            }
        }
        public int query(int left, int right)
        {
            int res = 0;
            left += size;
            right += size;
            while (left <= right) {
                if ((left & 1) == 1) res = Math.max(res, tree[left++]);
                if ((right & 1) == 0) res = Math.max(res, tree[right--]);
                left /= 2;
                right /= 2;
            }
            return res;
        }
    }
}
```

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```
public int lengthOfLIS(int[] nums, int k) {
    int maxVal = Arrays.stream(nums).max().getAsInt();
    SegmentTree segTree = new SegmentTree(maxVal + 1);
    int maxLen = 1;
    for (int num : nums) {
        int bestPrev = segTree.query(Math.max(1, num - k), num - 1);
        int newLength = bestPrev + 1;
        segTree.update(num, newLength);
        maxLen = Math.max(maxLen, newLength);
    }
    return maxLen;
}
```

4. Output:



The screenshot shows a code editor with a Java class named `SegmentTree`. The class has a static array `tree` and an integer `size`. It includes a constructor `SegmentTree(int n)` that initializes `size` to `n` and `tree` to a new array of size `2 * n`. It also has a method `update(int idx, int value)`. The test results section shows that the code was accepted with a runtime of 0 ms. The input for the test case is `nums = [4,2,1,4,3,4,5,8,15]` and `k = 3`. The output is `5`.

5. Time Complexity:

Time Complexity = $O(n \log n)$

Space Complexity = $O(n)$

6. Learning Outcome:

- Learned how different problems have varying complexities, ranging from $O(n)$ (Kadane's Algorithm) to $O(n \log n)$ (Merge Sort & Segment Tree).
- Explored recursion-based solutions (Longest Nice Substring) and Divide & Conquer techniques (Reverse Pairs using Merge Sort).
- Implemented Segment Tree for optimized Longest Increasing Subsequence II, reducing time complexity to $O(n \log n)$.
- Applied Kadane's Algorithm for Maximum Subarray, Merge Sort-based counting for Reverse Pairs, and Segment Tree-based LIS for efficient computations.