



Experiment 5

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Subject Name: Ap

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1. Aim : Sort Colors

2. **Objective :** Given an array `nums` with `n` objects colored red, white, or blue, sort them in-place so that objects of the same color are adjacent, with the colors in the order red, white, and blue.

We will use the integers 0, 1, and 2 to represent the color red, white, and blue, respectively.

You must solve this problem without using the library's sort function.

3. Code:

```
void sortColors(vector<int>& nums) {  
  
    int low = 0, mid = 0, high = nums.size() - 1;  
  
    while (mid <= high) {  
  
        if (nums[mid] == 0) {  
  
            swap(nums[low++], nums[mid++]);  
  
        } else if (nums[mid] == 1) {  
  
            mid++;  
  
        } else {  
  
            swap(nums[mid], nums[high--]);  
  
        }  
    }  
}
```



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}

}

4. Output:

Accepted Runtime: 0 ms

- Case 1
- Case 2

Input

```
nums =  
[2,0,2,1,1,0]
```

Output

```
[0,0,1,1,2,2]
```

Expected

```
[0,0,1,1,2,2]
```

Accepted Runtime: 0 ms

- Case 1
- Case 2

Input

```
nums =  
[2,0,1]
```

Output

```
[0,1,2]
```

Expected

```
[0,1,2]
```

5. Learning Outcomes

- Use of Binary Search Approach in it.
- Use of Swap Function in it.
- Use of Pointer Approach in it.

1. **Aim :- Median of Two Sorted Arrays**
2. **Objective :-** Given two sorted arrays nums1 and nums2 of size m and n respectively, return **the median** of the two sorted arrays. The overall run time complexity should be $O(\log(m+n))$.
3. **Code:-**

```
void merge(vector<int>& nums1, vector<int>& nums2, vector<long long >&v){

    int k = 0;
    int idx1 = 0;
    int idx2 = 0;

    while(idx1 < nums1.size() and idx2 < nums2.size()){

        if(nums1[idx1] <= nums2[idx2]){
            v.push_back(nums1[idx1++]);
        }else{
            v.push_back(nums2[idx2++]);
        }

    }

    while(idx1 < nums1.size()){
        v.push_back(nums1[idx1++]);
    }
    while(idx2 < nums2.size()){
        v.push_back(nums2[idx2++]);
    }

}

public:
    double findMedianSortedArrays(vector<int>& nums1, vector<int>& nums2) {

        vector<long long > v;

        merge(nums1,nums2,v);

        int index = v.size()/ 2;

        if(v.size() % 2 != 0){
            return v[index];
        }
```



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```
}else{  
  
    return (double)(v[index] + v[index - 1])/2;  
  
}  
}
```

4. OutPut:-

Case 1

Input

nums1 =[1,2]

nums2 =[3,4]

Output

2.50000

Expected

2.50000

Input

nums1 =[1,2]

nums2 =[3,4]

Output

2.50000

Expected

2.50000

5. Learning Outcomes:-

- 1.Using MergeSort Approach.
- 2.we can make it as an ideal approach in it.
- 3.using of function call.

1. **Aim :-** Kth Smallest Element in a Sorted Matrix
2. **Objective :-** Given an $n \times n$ matrix where each of the rows and columns is sorted in ascending order, return *the k^{th} smallest element in the matrix.*
Note that it is the k^{th} smallest element in the sorted order, not the k^{th} distinct element.
You must find a solution with a memory complexity better than $O(n^2)$.

3. Code:-

```
int kthSmallest(vector<vector<int>>& matrix, int k) {
```

```
    priority_queue<int>ans;
    for(int i=0;i<matrix.size();i++){
        for(int j=0;j<matrix.size();j++){

            if(ans.size() < k){
                ans.push(matrix[i][j]);
            }else{

                if(matrix[i][j]<ans.top()){
                    ans.pop();
                    ans.push(matrix[i][j]);
                }
            }
        }
    }
    return ans.top();}
```

4. Output :-

Case 1

Input

matrix =[[1,5,9],[10,11,13],[12,13,15]]

k = 8

Output

13

Expected

13



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Input

matrix = [[-5]]

k = 1

Output

-5

Expected

-5

5. Learning Outcomes:-

1. Use of Priority Queue.
2. Use of neasted loop.
3. Using the approach of queue.