Experiment - 2(A)

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Subject Name: Advanced Programming Lab-2 **Subject Code:** 22CSH-359

1. <u>Title:</u> Linked Lists (Remove duplicates from a sorted list) https://leetcode.com/problems/remove-duplicates-from-sorted-list

2. <u>Objective:</u> Given the head of a sorted linked list, the task is to remove all duplicates such that each element appears only once. Return the modified linked list, which is still sorted.

3. Algorithm:

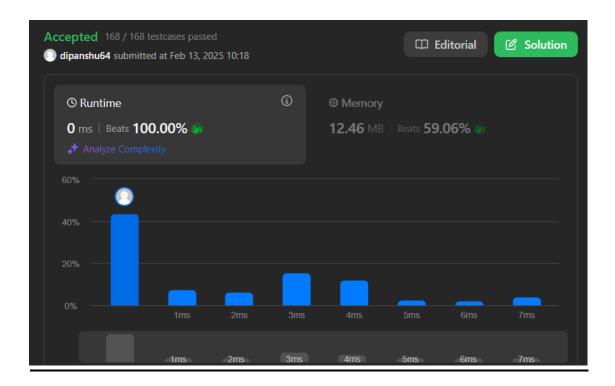
- Iterate Through the Linked List:
- Start with the head of the list.
- Traverse the list using a pointer current starting from the head.
 - Check for Duplicates:
- For each node, compare the value of the current node with the value of the next node.
- If the values are equal (i.e., a duplicate), update the current node's next pointer to skip the next node.
- If the values are not equal, simply move the current pointer to the next node.
 - End of List:
- Continue this process until you reach the end of the list (i.e., current.next is None).
 - Return Modified List:
- The linked list will be modified in-place with the duplicates removed.
- Return the modified head of the linked list.

4. Implementation/Code:

```
class Solution:
def deleteDuplicates(self, head):
    current = head
    while current and current.next:
        # If current node's value is equal to the next node's value
```

```
if current.val == current.next.val:
        current.next = current.next next # Skip the duplicate node
else:
        current = current.next # Move to the next node
return head
```

5. Output:



6. **Time Complexity:** O(n)

7. **Space Complexity:** O(1)

7. Learning Outcomes:

• In-place Modifications in Linked Lists:

• Understand how to perform operations on linked lists without using additional memory (i.e., modifying the list in-place).

• Traversing Linked Lists:

• Gain experience with traversing a linked list and manipulating the next pointers.

• Handling Edge Cases:

• The solution works even for edge cases where the linked list is empty (i.e., head is None) or has only one node.

Experiment 2(B)

- 1. **Title:** Reverse a linked list (https://leetcode.com/problems/reverse-linked-list/)
- **2.** Objective: Given the head of a singly linked list, the task is to reverse the list and return the reversed list.

3. Algorithm:

- Initialization:
 - Start with three pointers:
 - o prev initially set to None.
 - o current initialized to head.
 - o next_node will help in keeping track of the next node in the list.

• Traverse the List:

- For each node in the linked list:
 - Save the next node: next node = current.next.
 - Reverse the direction of the current node's next pointer: current.next = prev.
 - Move prev to the current node: prev = current.
 - Move current to the next node: current = next_node.

• End of List:

• When current becomes None, the list is reversed, and prev will be the new head of the reversed list.

• Return the Reversed List:

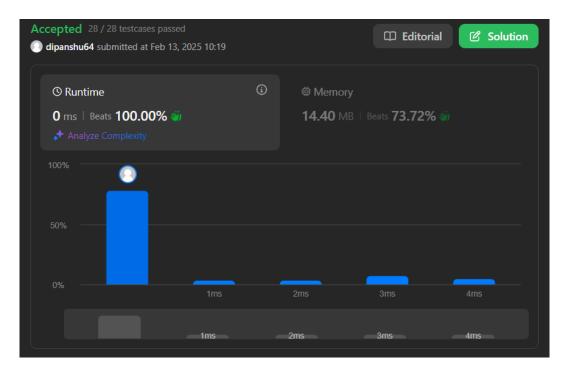
• The new head of the reversed linked list is prev.

4. Implementation/Code:

```
class ListNode:
    def __init__(self, val=0, next=None):
        self.val = val
        self.next = next

class Solution:
    def reverseList(self, head):
        prev = None
        current = head
        while current:
            next_node = current.next # Save next node
            current.next = prev # Reverse the current node's pointer
            prev = current # Move prev to current node
            current = next_node # Move current to next node
            return prev # Return the new head
```

5. Output:



6. Time Complexity: O(n)

7. **Space Complexity:** O(1)

8. <u>Learning Outcomes:</u>

• Reversing a Linked List:

• Learn how to reverse the direction of pointers in a singly linked list, both iteratively and recursively.

• Recursive vs Iterative Approaches:

• Understand the trade-offs between iterative and recursive approaches for solving linked list problems.

• Linked List Manipulation:

• Gain experience in manipulating linked list nodes and pointers to achieve desired outcomes (such as reversal).

Experiment 2(C)

- 1. <u>Title:</u> Delete middle node of a list (https://leetcode.com/problems/delete-the-middle-node-of-a-linked-list)
- **2.** Objective: Given the head of a linked list, the task is to delete the middle node and return the head of the modified list. The middle node is the $\lfloor n/2 \rfloor$ th node, where n is the length of the list, and $\lfloor x \rfloor$ represents the largest integer less than or equal to x.

3. Algorithm:

- Find the Length of the List:
 - Traverse the linked list and calculate the length n.
- Determine the Middle Node:
 - The middle node is at the index n // 2.
- Handle Edge Case for Small Lists:
 - If the list contains only one node (n == 1), return None (the list becomes empty).
- Traverse to the Node Before the Middle Node:
 - Use a pointer to traverse to the node just before the middle node.
- Delete the Middle Node:
 - Modify the next pointer of the node before the middle node to point to the node after the middle node.
- Return the Modified List:
 - Return the modified head of the linked list after the middle node is removed.
- 4. <u>Implementation/Code:</u>

```
class ListNode:
    def __init__(self, val=0, next=None):
        self.val = val
        self.next = next

class Solution:
    def deleteMiddle(self, head):
        # If the list contains only one node, return None (empty list after removal)
        if not head or not head.next:
            return None

        # Initialize two pointers: slow (to find the middle) and fast (to find the end)
        slow = head
        fast = head
        prev = None
```

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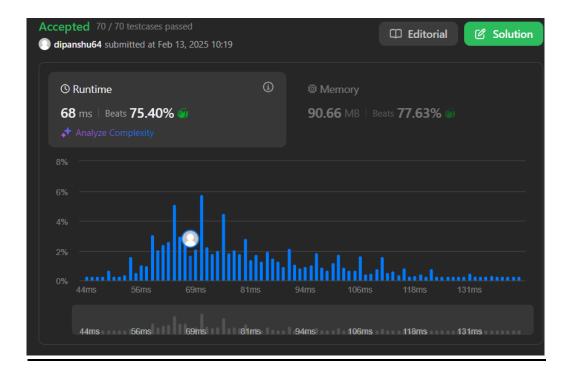
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```

```
# Traverse the list with fast moving two steps at a time and slow
moving one step
  while fast and fast.next:
     fast = fast.next.next
     prev = slow
     slow = slow.next

# Remove the middle node by connecting prev to slow.next
if prev:
     prev.next = slow.next

return head
```

5. Output:



9. Time Complexity: O(n)

7. Space Complexity: O(1)

10. Learning Outcomes:

- In-place Modifications in Linked Lists:
 - Learn how to modify linked lists in place, without creating new nodes or arrays.
- Two Pointer Technique:
 - Understand the usage of the slow and fast pointers to find specific nodes in a linked list (in this case, the middle node).