Experiment 3 A

Student Name: Sachin UID: 22BCS15411

Branch: CSE Section/Group: Ntpp 602-A

Semester: 6TH Date of Performance:03/02/25

Subject Name: AP Lab-2 Subject Code: 22CSH-352

1. TITLE:

Maximum Depth of Binary Tree

2. AIM:

Given the root of a binary tree, return its maximum depth.

A binary tree's maximum depth is the number of nodes along the longest path from the root node down to the farthest leaf node.

3. Algorithm

- Start DFS with the root node at depth 0.
- o If the node is null, return the current depth.
- o Recursively explore left and right children, increasing depth by 1.
- o Return the maximum depth from left or right subtree.

Implemetation/Code

```
class Solution {
public:
    int maxDepth(TreeNode* root) {
        if (root == nullptr) {
            return 0;
        } else {
            int leftDepth = maxDepth(root->left);
            int rightDepth = maxDepth(root->right);
            return 1 + std::max(leftDepth, rightDepth);
```

```
}
};
```

Output

```
Testcase >_ Test Result

• Case 1
• Case 2

Input

root =
[3,9,20,null,null,15,7]

Output

3

Expected

3
```

 $\textbf{Time Complexity}: O(\ n)$

Space Complexity : O(h)

Learning Outcomes:-

- Understand how to use depth-first search for tree traversal.
- o Gain skills in calculating the depth or height of binary trees



Experiment 3 B

Student Name: Sachin UID: 22BCS15411

Branch: CSE Section/Group: Ntpp 602-A

Semester: 6TH Date of Performance:03/02/25

Subject Name: AP Lab-2 Subject Code: 22CSH-352

1. TITLE:

KTH Smallest Element in a BST

2. AIM:

Given the root of a binary search tree, and an integer k, return the kth smallest value (1-indexed) of all the values of the nodes in the tree.

3. Algorithm

- Perform an in-order traversal of the binary tree starting from the root.
- Use a generator to yield nodes' values one by one in their in-order sequence.
- Iterate up to the kth element of the generator.
- Return the kth smallest element from the traversal.

Implementation/Code:

```
class Solution {
public:
    std::vector<std::vector<int>> zigzagLevelOrder(TreeNode* root) {
        std::vector<std::vector<int>> result;
        if (root == nullptr) {
            return result;
        }

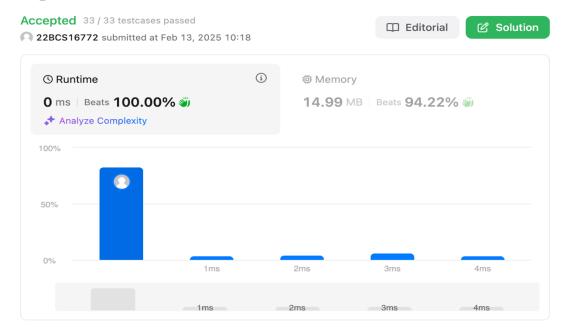
        std::queue<TreeNode*> q;
        q.push(root);
        bool leftToRight = true;

    while (!q.empty()) {
        int levelSize = q.size();
        std::vector<int> currentLevel(levelSize);

        for (int i = 0; i < levelSize; ++i) {</pre>
```

```
TreeNode* node = q.front();
                q.pop();
                // Store the value in the correct order
                if (leftToRight) {
                    currentLevel[i] = node->val;
                } else {
                    currentLevel[levelSize - 1 - i] = node->val; // Reverse
order
                }
                if (node->left) {
                    q.push(node->left);
                }
                if (node->right) {
                    q.push(node->right);
                }
            }
            result.push_back(currentLevel);
            leftToRight = !leftToRight; // Toggle direction for the next level
        }
        return result;
    }
};
```

Output





Time Complexity : O(k)

Space Complexity : O(h)

Learning Outcomes:-

- Learn how to perform and apply in-order traversal in binary trees to solve problems.
- Python generators to manage state and produce results on demand during tree traversal.

Experiment 3 C

Student Name: Sachin UID: 22BCS15411

Branch: CSE Section/Group:NTPP-602-A

Semester: 5 Date of Performance: 11/10/24

Subject Name: AP Lab Subject Code: 22CSH-311

1. TITLE: Fibonacci Numbers

2. AIM: The Fibonacci sequence appears in nature all around us, in the arrangement of seeds in a sunflower and the spiral of a nautilus for example. The Fibonacci sequence begins with Fibonacci (0)=0 and Fibonacci(1)=1 and as its first and second terms. After these first two elements, each subsequent element is equal to the sum of the previous two elements

3. Algorithm

- o Base Case Check: If n is 1, return 1 as the first Fibonacci number.
- o Initialize Variables: Set a and b to 1, representing the first two Fibonacci numbers.
- o Iterate Through Positions: Loop from position 2 to n-1.
- O Update Fibonacci Values: In each iteration, update a to b and b to a + b.
- o Return Result: After completing the loop, return a as the n-th Fibonacci number.

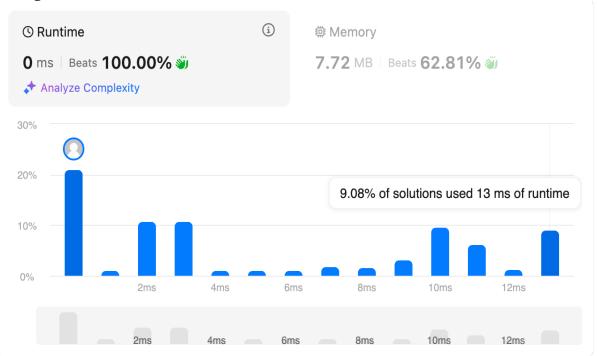
4. Implementation/Code

```
class Solution {
public:
    int fib(int n) {
        if (n <= 1) {
            return n;
        }

    // Efficient iterative approach (using only two variables)
    int a = 0; // F(0)
    int b = 1; // F(1)
    for (int i = 2; i <= n; ++i) {
        int temp = a + b; // Calculate F(i)
        a = b; // Update F(i-2)
        b = temp; // Update F(i-1)</pre>
```

% return b; // F(n) is now stored in b } }; Time Complexity : O(N) Space Complexity : O(1)

5. Output



6. Learning Outcomes:-

- 1. Understand how to compute Fibonacci numbers using an iterative approach.
- 2. Learn to optimize space by maintaining only two variables instead of using an entire array.
- 3. Gain proficiency in implementing efficient loop constructs for sequential computations.