Experiment - 4(A)

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Branch: CSE **Section/Group:** NTPP_602-A **Date of Performance:** 17-02-25

Subject Name: Advanced Programming Lab-2 **Subject Code:** 22CSH-359

1. <u>Title:</u> Sorting and Searching (Merge Sorted Array) https://leetcode.com/problems/merge-sorted-array/

2. <u>Objective:</u> To merge two sorted arrays nums1 and nums2 into one sorted array, while ensuring that the result is stored in nums1. The algorithm should be efficient and merge the arrays in-place.

3. Algorithm:

- Start with three pointers: i for the last element of nums1 (index m-1), j for the last element of nums2 (index n-1), and k for the last position in nums1 (index m+n-1).
- Compare the elements at nums1[i] and nums2[j]. Place the larger of the two at nums1[k], and decrement k.
- If nums1[i] is greater, move i left. If nums2[j] is greater, move j left.
- If any elements are left in nums2, copy them to nums1 (since nums1 already contains the first m elements in sorted order).
- Continue the process until all elements from both arrays are merged.

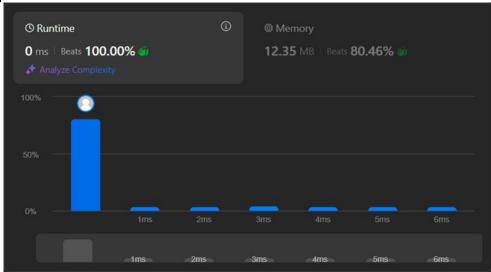
4. Implementation/Code:

```
class Solution:
def merge(self, nums1, m, nums2, n):
    # Start from the end of both arrays
i, j, k = m - 1, n - 1, m + n - 1

# Merge in reverse order
while i >= 0 and j >= 0:
    if nums1[i] > nums2[j]:
        nums1[k] = nums1[i]
        i -= 1

else:
        nums1[k] = nums2[j]
        j -= 1
        k -= 1
```

5. Output:



- **6.** Time Complexity: O(m+n)
- 7. **Space Complexity:** O(1)

8. <u>Learning Outcomes:</u>

- Array Manipulation: Understanding how to manipulate arrays in-place.
- **Two Pointer Technique**: Using multiple pointers to traverse arrays efficiently and merge them.
- **In-place Sorting**: Learning how to merge sorted arrays without extra space.
- **Optimization**: Achieving linear time complexity by processing the arrays from the end, avoiding unnecessary moves.

Experiment 4(B)

- **1. <u>Title:</u>** First Bad Colors (https://leetcode.com/problems/first-bad-version/)
- **2.** <u>Objective:</u> The task is to find the first bad version in a series of product versions, where all versions after the first bad version are also bad. You should minimize the number of calls to the isBadVersion(version) API.

3. Algorithm:

- Initialize two pointers left as 1 and right as n.
- Perform a binary search:
 - While left is less than right, calculate the mid-point mid = (left + right) // 2.
 - If isBadVersion (mid) is True, then the first bad version is at mid or to the left of it, so update right = mid.
 - If isBadVersion (mid) is False, then the first bad version must be to the right, so update left = mid + 1.
- When the loop terminates, left will point to the first bad version.

4. Implementation/Code:

return left

```
class Solution:
    def firstBadVersion(self, n):
        left, right = 1, n

    while left < right:
        mid = left + (right - left) // 2
        if isBadVersion(mid):
            right = mid # Potential first bad version
        else:
        left = mid + 1 # Bad version must be on the right</pre>
```

6. Output:



- **8.** <u>Time Complexity:</u> O(log n)
- **9.** Space Complexity: O(1)

9. Learning Outcomes:

- 1. **Binary Search**: Understanding and applying the binary search algorithm to reduce the number of checks (calls to isBadVersion).
- 2. **Optimization**: Minimizing the number of calls to an external API, which is a key performance concern in real-world applications.
- 3. **Problem Solving**: Effectively narrowing down the range for the first bad version using a divide-and-conquer strategy.

Experiment 4(C)

- **1.** <u>Title:</u> Sort Colors (https://leetcode.com/problems/sort-colors/)
- **2. Objective:** To sort an array containing three distinct values (0, 1, 2) in place, where 0 represents red, 1 represents white, and 2 represents blue. The array should be sorted in such a way that all 0s come first, followed by all 1s, and then all 2s.

4. Algorithm:

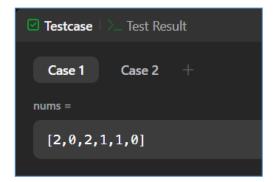
- Initialize three pointers: low (starting at the beginning of the array), mid (also starting at the beginning), and high (starting at the end of the array).
- Traverse the array using the mid pointer:
 - If nums [mid] is 0, swap it with nums [low] and increment both low and mid.
 - If nums [mid] is 1, just increment mid.
 - If nums [mid] is 2, swap it with nums [high] and decrement high.
- Continue this process until mid exceeds high.

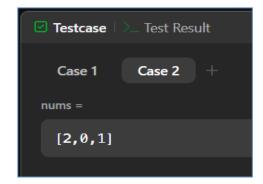
5. Implementation/Code:

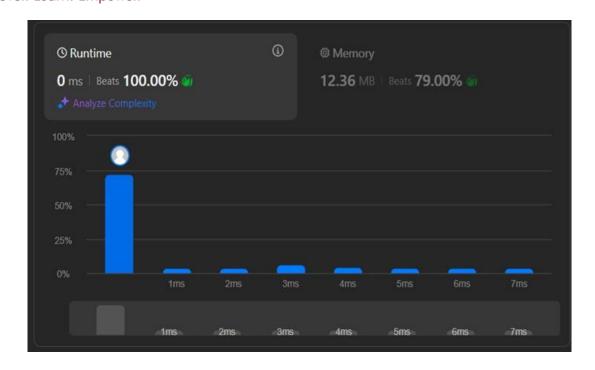
```
class Solution:
    def sortColors(self, nums):
        low, mid, high = 0, 0, len(nums) - 1

    while mid <= high:
        if nums[mid] == 0:
            nums[low], nums[mid] = nums[mid], nums[low]
            low += 1
            mid += 1
        elif nums[mid] == 1:
            mid += 1
        else:
            nums[mid], nums[high] = nums[high], nums[mid]
            high -= 1</pre>
```

6. Output:







8. Time Complexity: O(N)

9. Space Complexity: O(1)

10. Learning Outcomes:

- 1. **Three-way Partitioning**: Understanding how to partition an array into three groups based on a specific condition (here, 0, 1, 2).
- 2. **In-place Sorting**: Sorting without using extra space for another array.
- 3. **Two Pointer Technique**: Efficient use of multiple pointers (low, mid, high) to traverse the array in a single pass.
- 4. **Optimization**: Solving the problem in linear time (O(n)) and constant space (O(1)).