

DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

Discover. Learn. Empower.

Experiment-4(A)

Student Name: Pranav Saini

Branch: CSE

Semester: 6

Subject Name: Advanced Programming Lab-2

UID: 22BCS15994

Section/Group: 602-A

Date of Performance: 13-02-25

Subject Code: 22CSH-359

1. **Title:** Sorting and Searching (Merge Sorted Array)
<https://leetcode.com/problems/merge-sorted-array/>

2. **Objective:** To merge two sorted arrays `nums1` and `nums2` into one sorted array, while ensuring that the result is stored in `nums1`. The algorithm should be efficient and merge the arrays in-place.

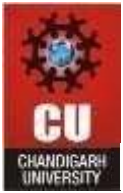
3. **Algorithm:**

- Start with three pointers: `i` for the last element of `nums1` (index `m-1`), `j` for the last element of `nums2` (index `n-1`), and `k` for the last position in `nums1` (index `m+n-1`).
- Compare the elements at `nums1[i]` and `nums2[j]`. Place the larger of the two at `nums1[k]`, and decrement `k`.
- If `nums1[i]` is greater, move `i` left. If `nums2[j]` is greater, move `j` left.
- If any elements are left in `nums2`, copy them to `nums1` (since `nums1` already contains the first `m` elements in sorted order).
- Continue the process until all elements from both arrays are merged.

4. **Implementation/Code:**

```
class Solution:
    def merge(self, nums1, m, nums2, n):
        # Start from the end of both arrays i, j, k
        i = m - 1, j = n - 1, k = m + n - 1

        # Merge in reverse order
        while i >= 0 and j >= 0:
            if nums1[i] > nums2[j]:
                nums1[k] = nums1[i]
                i -= 1
            else:
                nums1[k] = nums2[j]
                j -= 1
            k -= 1
```

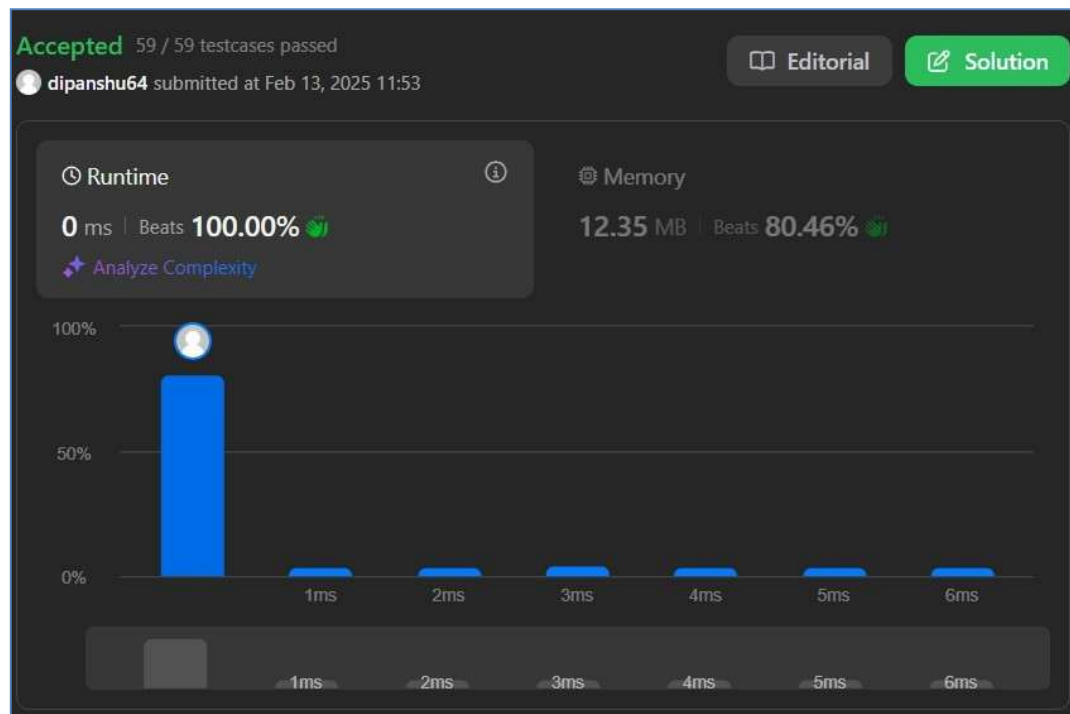


DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

Discover. Learn. Empower.

```
# If nums2 is not fully traversed, copy the remaining elements while j >= 0:  
    nums1[k] = nums2[j]  
    j -= 1  
    k -= 1
```

5. Output:

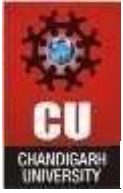


6. Time Complexity: $O(m+n)$

7. Space Complexity: $O(1)$

8. Learning Outcomes:

- **Array Manipulation:** Understanding how to manipulate arrays in-place.
- **Two Pointer Technique:** Using multiple pointers to traverse an array efficiently and merge them.
- **In-place Sorting:** Learning how to merge sorted arrays without extra space.
- **Optimization:** Achieving linear time complexity by processing the arrays from the end, avoiding unnecessary moves.



Experiment 4(B)

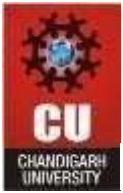
1. **Title:** FirstBadColors(<https://leetcode.com/problems/first-bad-version/>)
2. **Objective:** The task is to find the first bad version in a series of product versions, where all versions after the first bad version are also bad. You should minimize the number of calls to the `isBadVersion(version)` API.
3. **Algorithm:**
 - Initialize two pointers `left` as 1 and `right` as `n`.
 - Perform a binary search:
 - While `left` is less than `right`, calculate the mid-point `mid = (left + right) // 2`.
 - If `isBadVersion(mid)` is `True`, then the first bad version is at `mid` or to the left of it, so update `right = mid`.
 - If `isBadVersion(mid)` is `False`, then the first bad version must be to the right, so update `left = mid + 1`.
 - When the loop terminates, `left` will point to the first bad version.

4. **Implementation/Code:**

```
class Solution:
    def firstBadVersion(self, n):
        left, right = 1, n

        while left < right:
            mid = left + (right - left) // 2
            if isBadVersion(mid):
                right = mid # Potential first bad version
            else:
                left = mid + 1 # Bad version must be on the right

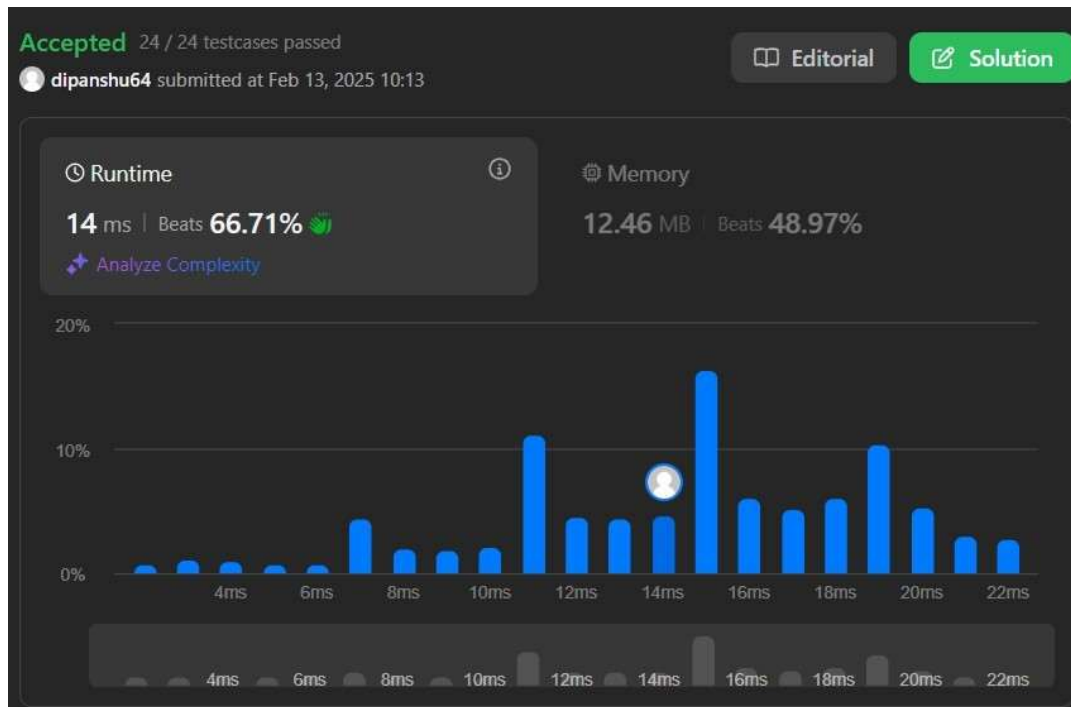
        return left
```



DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

Discover. Learn. Empower.

6. Output:

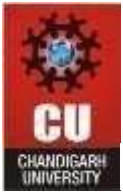


8. Time Complexity: $O(\log n)$

9. Space Complexity: $O(1)$

9. Learning Outcomes:

1. **Binary Search:** Understanding and applying the binary search algorithm to reduce the number of checks (calls to `isBadVersion`).
2. **Optimization:** Minimizing the number of calls to an external API, which is a key performance concern in real-world applications.
3. **Problem Solving:** Effectively narrowing down the range for the first bad version using a divide-and-conquer strategy.



Experiment 4(C)

1. **Title:** SortColors(<https://leetcode.com/problems/sort-colors/>)
2. **Objective:** To sort an array containing three distinct values (0, 1, 2) in place, where 0 represents red, 1 represents white, and 2 represents blue. The array should be sorted in such a way that all 0s come first, followed by all 1s, and then all 2s.

4. Algorithm:

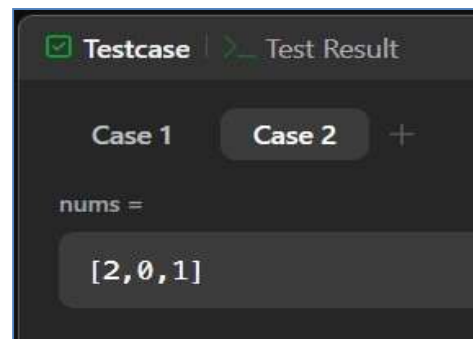
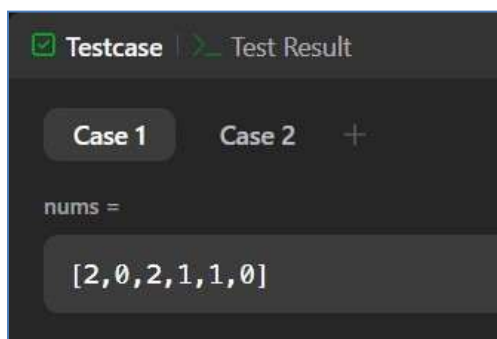
- Initialize three pointers: `low` (starting at the beginning of the array), `mid` (also starting at the beginning), and `high` (starting at the end of the array).
- Traverse the array using the `mid` pointer:
 - If `nums[mid]` is 0, swap it with `nums[low]` and increment both `low` and `mid`.
 - If `nums[mid]` is 1, just increment `mid`.
 - If `nums[mid]` is 2, swap it with `nums[high]` and decrement `high`.
- Continue this process until `mid` exceeds `high`.

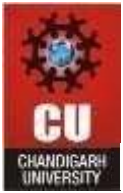
5. Implementation/Code:

```
class Solution:
    def sortColors(self, nums):
        low, mid, high = 0, 0, len(nums) - 1

        while mid <= high:
            if nums[mid] == 0:
                nums[low], nums[mid] = nums[mid], nums[low]
                low += 1
                mid += 1
            elif nums[mid] == 1:
                mid += 1
            else:
                nums[mid], nums[high] = nums[high], nums[mid]
                high -= 1
```

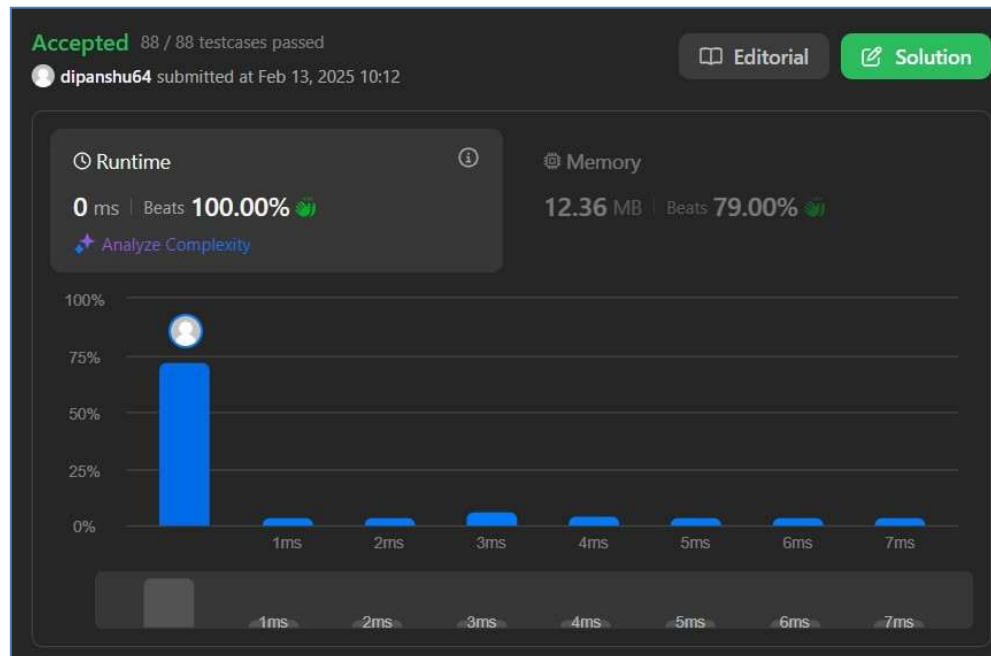
6. Output:





DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

Discover. Learn. Empower.



8. Time Complexity: $O(N)$

9. Space Complexity: $O(1)$

10. Learning Outcomes:

1. **Three-way Partitioning**: Understanding how to partition an array into three groups based on a specific condition (here, 0, 1, 2).
2. **In-place Sorting**: Sorting without using extra space for another array.
3. **Two Pointer Technique**: Efficient use of multiple pointers (low, mid, high) to traverse the array in a single pass.
4. **Optimization**: Solving the problem in linear time ($O(n)$) and constant space ($O(1)$).