

DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

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Experiment-6(A)

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Semester: 6

Subject Name: Advanced Programming Lab-2

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Section/Group: NTPP_602-A

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Subject Code: 22CSH-359

1. **Title:** Dynamic Programming (Climbing Stairs)

2. **Objective:** The problem is to find the total number of distinct ways to reach the top of a staircase with n steps, where at each step, you can either climb 1 or 2 steps.

3. **Algorithm:**

- **Understanding the Problem:**

- For $n = 1$: Only 1 way $\rightarrow [1]$
- For $n = 2$: Two ways $\rightarrow [1, 1], [2]$
- For $n = 3$: Three ways $\rightarrow [1, 1, 1], [1, 2], [2, 1]$
- For $n = 4$: Five ways $\rightarrow [1, 1, 1, 1], [1, 1, 2], [1, 2, 1], [2, 1, 1], [2, 2]$

- **Pattern Identification:**

- The problem resembles the Fibonacci sequence: • **$\text{Ways}(n) = \text{Ways}(n-1) + \text{Ways}(n-2)$** •

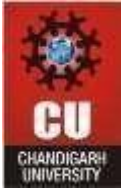
Base cases:

- $\text{Ways}(1) = 1$ ◦
- $\text{Ways}(2) = 2$

- **Approach:**

- Initialize two variables $\text{first} = 1$ (Ways to reach step 1) and $\text{second} = 2$ (Ways to reach step 2).
- For steps from 3 to n , calculate the number of ways using:

```
python CopyEdit
current = first + second
first = second
second = current
```



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- Return `second` as it will hold the answer for n .

4. Implementation/Code:

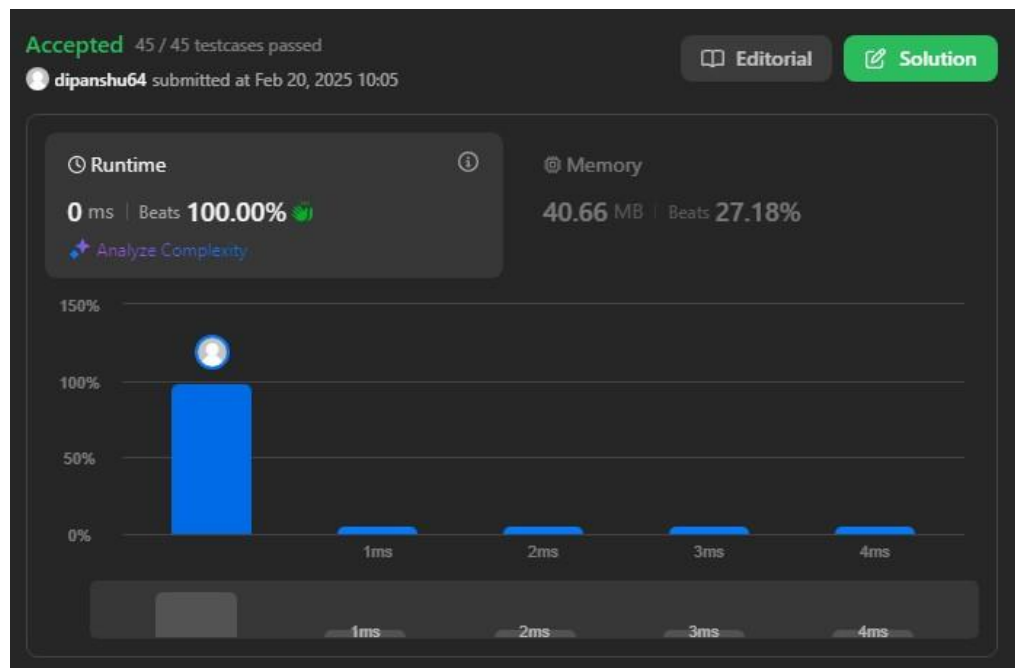
```
class Solution:    def climbStairs(self,
n: int) -> int:
    # Base cases
    if n <= 2:
        return n

    # Initialize the first two steps
    first, second = 1, 2

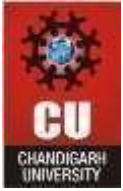
    # Compute the ways to reach each step from 3 to n
    for i in range(3, n + 1):
        current = first + second
        first = second
        second = current

    return second
```

5. Output:



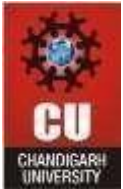
6. Time Complexity: $O(N)$



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7. Space Complexity: $O(1)$



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Experiment 6(B)

1. Title: Jump Game

2. Objective: Determine if you can reach the last index of an array where each element represents the maximum jump length from that position.

3. Algorithm:

1. **Initialize:** A variable `maxReach` to 0, representing the furthest index we can reach.

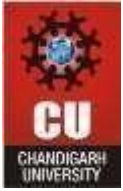
2. **Iterate:** Through each index `i` in the array:

- If `i > maxReach`, return `false` (i.e., current index is unreachable).
- Update `maxReach` as `max(maxReach, i + nums[i])`.
- If `maxReach` is greater than or equal to the last index, return `true`.

Return: If the loop completes without returning, it means the last index is reachable, so return `true`.

4. Implementation/Code:

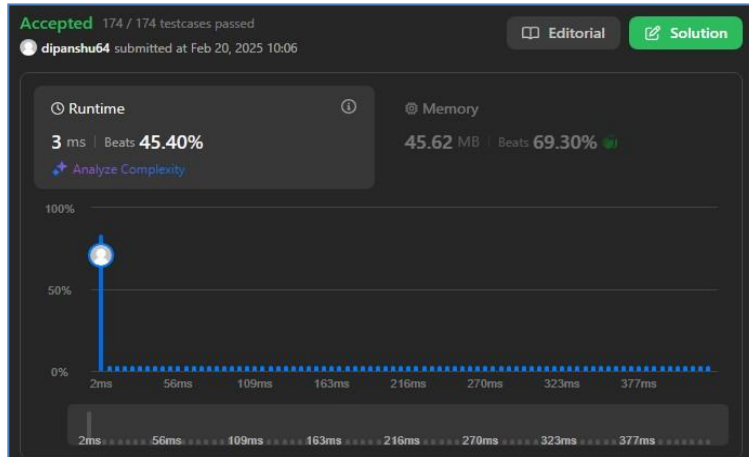
```
class Solution:
    def canJump(self, nums:
list[int]) -> bool:
        maxReach = 0
        for i in
range(len(nums)):
            if i > maxReach:
                return False
            maxReach = max(maxReach,
i + nums[i])
            if maxReach >= len(nums)
- 1:
                return True
        return True
```



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6. Output:



8. Time Complexity: $O(N)$

9. Space Complexity: $O(1)$

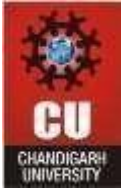
Experiment 6(C)

1. Title: Maximum Subarray

2. Objective: To find the contiguous subarray with the largest sum in a given integer array nums.

3. Algorithm:

- **Initialization:**
 - `currentSum = 0` (stores sum of the current subarray)
 - `maxSum = -infinity` (stores the maximum sum found so far)
- **Iteration through the array:**
 - For each element `num` in `nums`:
 - Add `num` to `currentSum`.
 - Update `maxSum` to the maximum of `maxSum` and `currentSum`.
 - If `currentSum` becomes negative, reset it to 0 (discard the current subarray).
- **Result:**



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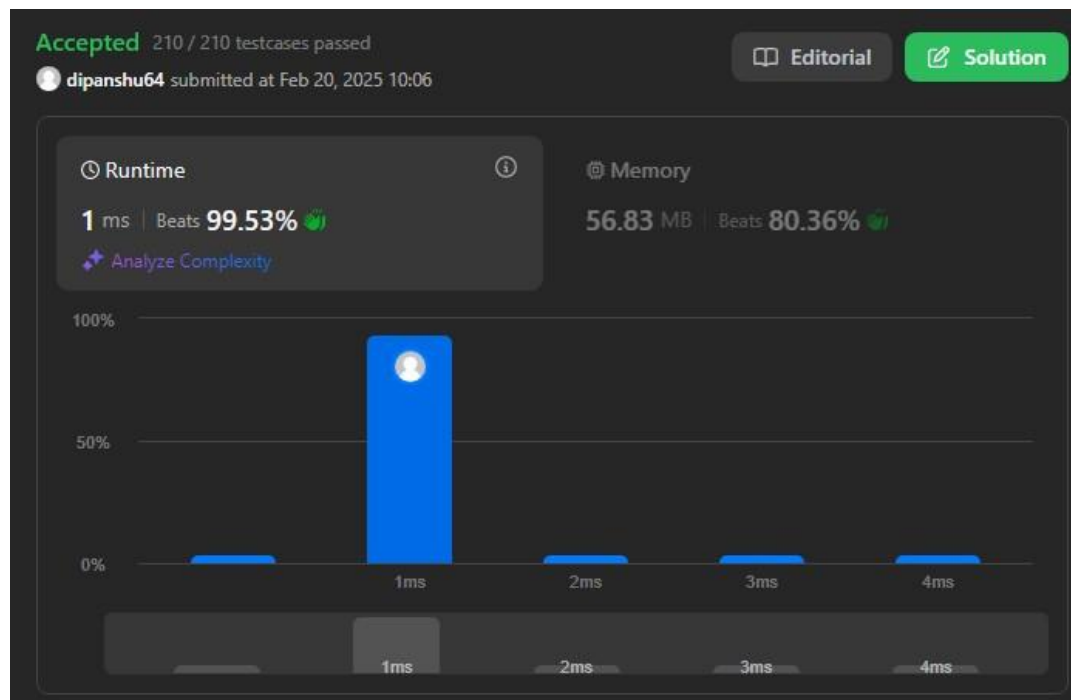
- Return `maxSum` as the maximum sum of the subarray.

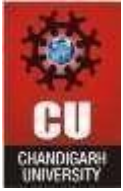
5. Implementation/Code:

```
class Solution:    def maxSubArray(self, nums:
list[int]) -> int:
        currentSum = 0
maxSum = float('-inf')
        for num
in nums:
            currentSum += num
            maxSum = max(maxSum, currentSum)
if currentSum < 0:
    currentSum = 0

return maxSum
```

6. Output :





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8. Time Complexity: $O(N)$

9. Space Complexity: $O(1)$

10. Learning Outcomes:

- **Kadane's Algorithm:** A powerful technique to solve maximum subarray problems in linear time.
- **Handling Negatives:** Resetting `currentSum` when it goes negative is key to maintaining the optimal subarray.
- **Optimized Approach:** Avoids nested loops, ensuring efficiency even for large input sizes.