

# DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

Discover. Learn. Empower.

## Experiment-7(A)

**Student Name:** Dipanshu Sharma

**Branch:** CSE

**Semester:** 6

**Subject Name:** Advanced Programming Lab-2

**UID:** 22BCS16030

**Section/Group:** NTPP\_602-A

**Date of Performance:** 10-03-25

**Subject Code:** 22CSH-359

1. **Title:** Greedy ( Maximum Units on a Truck )
2. **Objective:** The objective is to maximize the total number of units loaded onto a truck given an array of box types and a truck size constraint.
3. **Algorithm:**
  - **Input:** An array boxTypes where each element represents [numberOfBoxes, unitsPerBox] and an integer truckSize.
  - **Sorting Step:** Sort boxTypes in **descending order** based on the number of units per box.
  - **Initialization:**
    - maxUnits = 0 □ Stores the total units collected.
  - **Iteration:**
    - For each box in the sorted list:
      - If truckSize == 0, stop the iteration.
      - Add min(box[0], truckSize) \* box[1] to maxUnits.
      - Subtract min(box[0], truckSize) from truckSize.
  - **Return maxUnits** as the maximum units loaded onto the truck.

## 4. **Implementation/Code:**

```
class Solution:
    def maximumUnits(self, boxTypes, truckSize):
        # Sort box types in descending order of units per box
```



# DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

Discover. Learn. Empower.

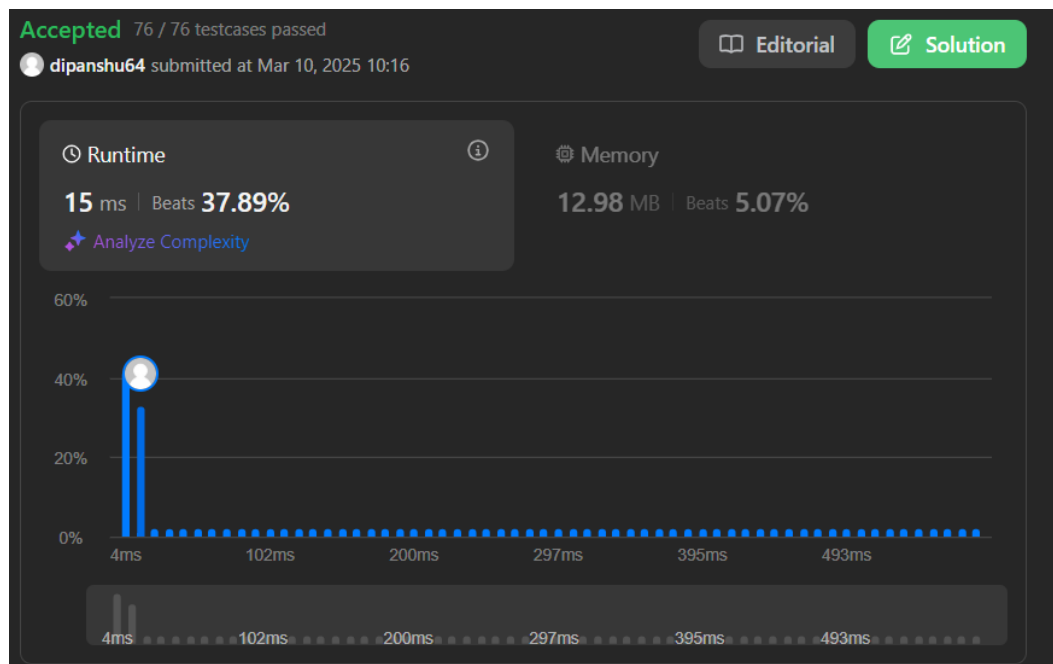
```
boxTypes.sort(key=lambda x: x[1], reverse=True)

maxUnits = 0
for boxes, units in boxTypes:
    if truckSize == 0:
        break
    count = min(boxes, truckSize)
    maxUnits += count * units
    truckSize -= count

return maxUnits

# Example Usage
print(Solution().maximumUnits([[1,3],[2,2],[3,1]], 4)) # Output: 8
print(Solution().maximumUnits([[5,10],[2,5],[4,7],[3,9]], 10)) # Output:
91
```

## 5. Output:



## 6. Time Complexity: $O(n \log n)$

## 7. Space Complexity: $O(1)$



# DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

Discover. Learn. Empower.

## Experiment 7(B)

1. **Title:** Maximum Subarray
2. **Objective:** To find the contiguous subarray with the largest sum in a given integer array.
3. **Algorithm:**

1. **Input:** An array `nums` containing integers.

2. **Initialization:**

- o `currentSum = 0` (Tracks the sum of the current subarray)
- o `maxSum = -∞` (Tracks the maximum subarray sum found so far)

3. **Iteration:**

- o For each element `num` in `nums`:

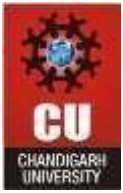
1. Add `num` to `currentSum`.
2. Update `maxSum = max(maxSum, currentSum)`.
3. If `currentSum` becomes negative, reset it to zero.
4. **Return** `maxSum` as the maximum subarray sum.

4. **Implementation/Code:**

```
class Solution:
    def maxSubArray(self, nums):
        currentSum = 0
        maxSum = float('-inf')

        for num in nums:
            currentSum += num
            maxSum = max(maxSum, currentSum)
            if currentSum < 0:
                currentSum = 0
        return maxSum

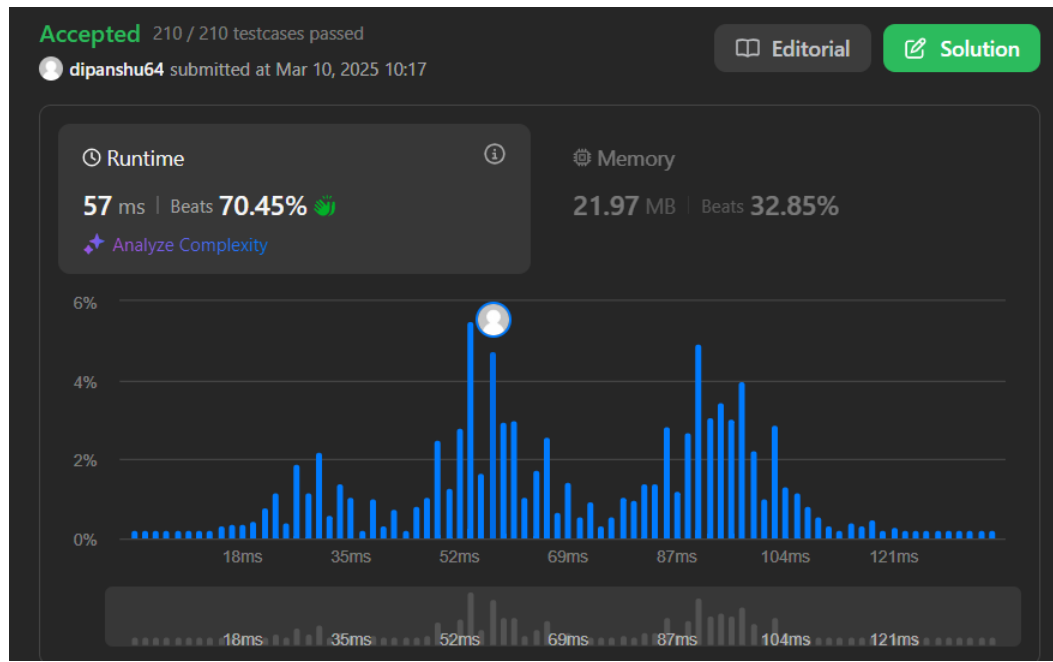
# Example Usage
print(Solution().maxSubArray([-2,1,-3,4,-1,2,1,-5,4]))
print(Solution().maxSubArray([5,4,-1,7,8]))
```



# DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

Discover. Learn. Empower.

## 6. Output:

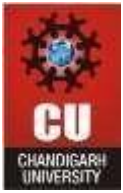


8. Time Complexity:  $O(n)$

9. Space Complexity:  $O(1)$

## 10. Learning Outcome:

- Mastered Kadane's Algorithm for maximum subarray sum.
- Learned efficient handling of negative values in subarrays.
- Improved understanding of optimal subarray identification in  $O(n)$  time.



# DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

Discover. Learn. Empower.

## Experiment 7(C)

1. **Title:** Remove Stones to Minimize the Total

2. **Objective:** To minimize the total number of stones remaining in the piles by repeatedly removing  $\text{floor}(\text{piles}[i] / 2)$  stones from the pile with the maximum stones, exactly  $k$  times.

### **3. Algorithm:**

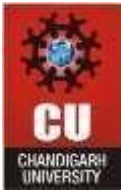
- **Input:** An array `piles` representing the number of stones in each pile, and an integer  $k$ .
- **Max Heap Conversion:**
  - Convert all elements of `piles` into their **negative values** and insert them into a **heap** (since Python's `heapq` only supports min-heaps, negating values makes it behave like a max-heap).
- **Perform  $k$  Operations:**
  - Repeat  $k$  times:
    - Extract the largest element from the heap (using `heapq.heappop`).
    - Remove  $\text{floor}(\text{largest} / 2)$  stones from it.
    - Insert the updated pile value back into the heap.
- **Calculate Remaining Stones:**
  - Sum all the remaining elements in the heap (converting them back to positive values).
- **Return** the total number of remaining stones.

### **5. Implementation/Code:**

```
import heapq
import math

class Solution:
    def minStoneSum(self, piles, k):
        # Convert to negative values to simulate max heap
        max_heap = [-pile for pile in piles]
        heapq.heapify(max_heap)

        # Perform k operations
        for _ in range(k):
            largest = -heapq.heappop(max_heap) # Get the max element
            reduced_stones = largest - largest // 2 # Correct floor
            division
```



# DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

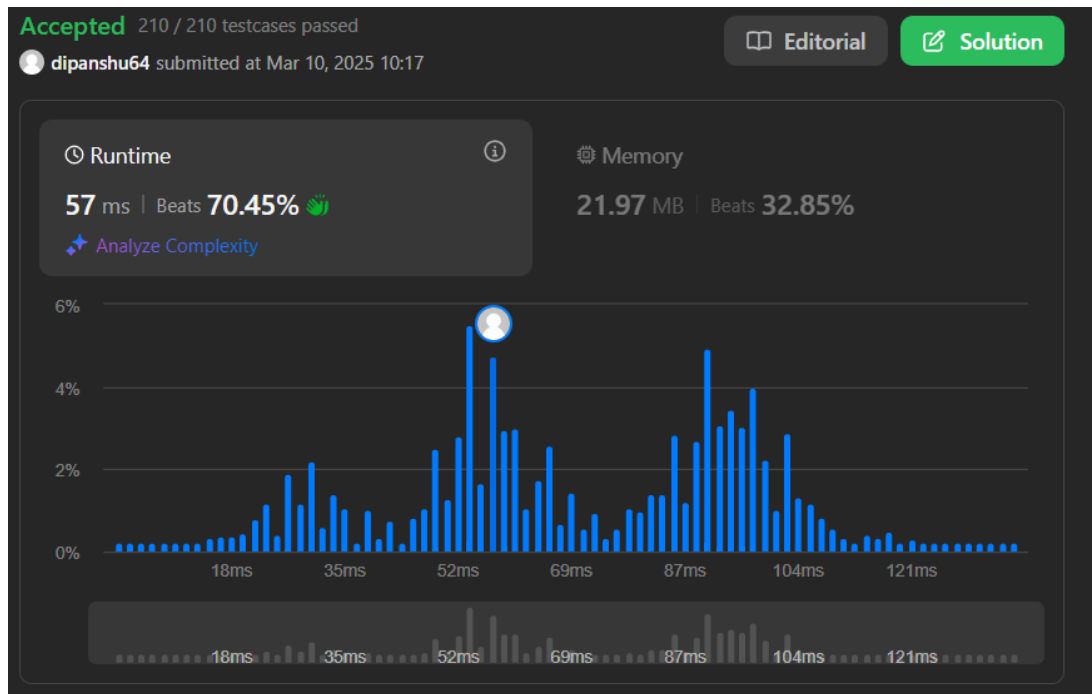
Discover. Learn. Empower.

```
        heapq.heappush(max_heap, -reduced_stones)    # Push back the
reduced pile

        # Calculate remaining stones (Convert back to positive)
        return int(-sum(max_heap))

# Example Usage
print(Solution().minStoneSum([5, 4, 9], 2)) # Output: 12
print(Solution().minStoneSum([4, 3, 6, 7], 3)) # Output: 12
```

## 6. Output:



8. Time Complexity:  $O(k \log n)$

9. Space Complexity:  $O(n)$

## 10. Learning Outcomes:

- Mastered Kadane's Algorithm for maximum subarray sum.
- Implemented DP solutions for optimal substructure problems.
- Applied Greedy Approach for maximizing outcomes.
- Utilized Max Heap for efficient element extraction and modification.
- Improved understanding of optimization techniques.