

Research on Strong Subadditivity of Quantum Information

量子信息的强次可加性的研究

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- 1 Introduction
 - Information theory
- 2 operator extension of SSA
 - Proof
- 3 SSA and Quantum markov chain
 - SSA equality
 - Operator SSA equality
- 4 Future work

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concepts

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- $H(X|Y) \equiv H(X, Y) - H(Y)$

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- $H(X)$
- $H(X|Y) \equiv H(X, Y) - H(Y)$
- $H(X : Y) \equiv H(X) + H(Y) - H(X, Y)$

inequalities

independency

- $H(X : Y) \geq 0$ Subadditivity
 $H(X) + H(Y) - H(X, Y) \geq 0$

inequalities

independency

- $H(X : Y) \geq 0$ Subadditivity
 $H(X) + H(Y) - H(X, Y) \geq 0$
- $H(X : Z|Y) \geq 0$ Strong Subadditivity(SSA)
 $H(X, Y) + H(Y, Z) - H(X, Y, Z) - H(Y) \geq 0$

Quantum Information

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Quantum Information

- $S(\rho) = -\text{Tr}(\rho \log \rho)$
- Does SSA holds ?
- Lieb et al.(1973) proof SSA
- Patrick Hayden et al.(2004) proof $\text{SSA}=0$ implies quantum markov state
- Ting-Chun Lin et al.(2023) proof a operator extension of SSA

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operator form

$$\hat{A} = \log \rho_{AB} + \log \rho_{BC} - \log \rho_A - \log \rho_C \leq 0$$

recover by considering $\langle \hat{A} \rangle$

operator norm

consider density matrix in AB,BC.

$$\rho_{AB} \otimes \sigma_C^{-1} \leq \rho_A \otimes \sigma_{BC}^{-1}$$

(f(t)=logt operator monotone)

$$(\rho_A^{-\frac{1}{2}} \otimes \sigma_{BC}^{\frac{1}{2}})(\rho_{AB} \otimes \sigma_C^{-1})(\rho_A^{-\frac{1}{2}} \otimes \sigma_{BC}^{\frac{1}{2}}) \leq I_{ABC}$$

$$\|(\rho_{AB}^{\frac{1}{2}} \otimes \sigma_C^{-\frac{1}{2}})(\rho_A^{-\frac{1}{2}} \otimes \sigma_{BC}^{\frac{1}{2}})\| \leq 1$$

diagram

proof the below operator has norm ≤ 1 .

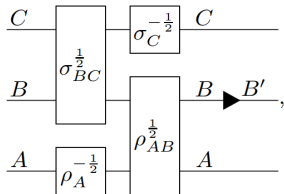
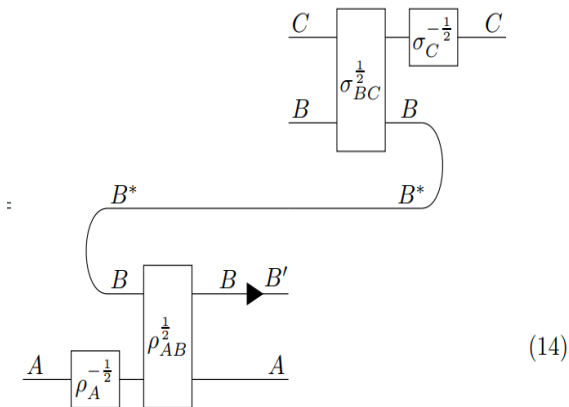


Figure: Caption

Isometry

proof the left and right part has norm 1.



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Quantum markov chain

ABC is a quantum markov chain $\iff H(A:C|B)$

Two Characterization

ABC is a quantum markov chain, then

$$\rho_{ABC} = (\mathcal{I}_A \otimes \mathcal{R}_{B \rightarrow BC})(\rho_{AB}),$$

Figure: Caption

$$\mathcal{T}_{B \rightarrow BC} : X_B \mapsto \rho_{BC}^{\frac{1}{2}} (\rho_B^{-\frac{1}{2}} X_B \rho_B^{-\frac{1}{2}} \otimes \text{id}_C) \rho_{BC}^{\frac{1}{2}}$$

Figure: Caption

Two Characterization

Theorem 6 *A state ρ_{ABC} on $\mathcal{H}_A \otimes \mathcal{H}_B \otimes \mathcal{H}_C$ satisfies strong subadditivity (eq. (1)) with equality if and only if there is a decomposition of system B as*

$$\mathcal{H}_B = \bigoplus_j \mathcal{H}_{b_j^L} \otimes \mathcal{H}_{b_j^R}$$

into a direct sum of tensor products, such that

$$\rho_{ABC} = \bigoplus_j q_j \rho_{Ab_j^L} \otimes \rho_{b_j^R C},$$

with states $\rho_{Ab_j^L}$ on $\mathcal{H}_A \otimes \mathcal{H}_{b_j^L}$ and $\rho_{b_j^R C}$ on $\mathcal{H}_{b_j^R} \otimes \mathcal{H}_C$, and a probability distribution $\{q_j\}$.

equivalent condition

$$\blacksquare \hat{A} = 0 \implies \langle \hat{A} \rangle = 0 \iff \text{SSA} = 0$$

equivalent condition

- $\hat{A} = 0 \implies \langle \hat{A} \rangle = 0 \iff SSA = 0$
- $\hat{A} = 0 \iff (\rho_{AB}^{\frac{1}{2}} \otimes \sigma_B^{-\frac{1}{2}})(\rho_A^{-\frac{1}{2}} \otimes \sigma_{BC}^{\frac{1}{2}})$ is an isometry



Condition of isometry

recall another form of $(\rho_{AB}^{\frac{1}{2}} \otimes \sigma_B^{-\frac{1}{2}})(\rho_A^{-\frac{1}{2}} \otimes \sigma_{BC}^{\frac{1}{2}})$

$$\left\| (I_A \otimes V_{B \rightarrow B'}^\dagger \otimes I_C)(I_A \otimes I_{B'} \otimes V_{C \rightarrow BB^*C}^\sigma)^\dagger (V_{A \rightarrow AB'B^*}^\rho \otimes I_B \otimes I_C) \right\|$$

Figure: Caption

denote the three term as A, B^*, C , the above can be written as AB^*C

Condition of isometry

$A : v_1 \rightarrow W, B : v_2 \rightarrow W$ A, B is isometry, use Direct Sum Decomposition, we can get

theorem

$$w \notin \text{Im}_A \iff \|A^*w\| < \|w\| \text{ [Nyc, Wyf]}$$

Condition of isometry

$A : v_1 \rightarrow W, B : v_2 \rightarrow W$ A, B is isometry, use Direct Sum Decomposition, we can get

theorem

$$w \notin \text{Im}_A \iff \|A^*w\| < \|w\| \text{ [Nyc, Wyf]}$$

corollary

$$A^*B \text{ is isometry} \iff \text{Im}_B \subseteq \text{Im}_A$$

equality condition

consider dimension of Im_B and Im_C , we get

operator SSA equality condition

$$\hat{A} = 0 \iff \dim(B) = 0$$

The condition is too strong.

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Future work

- new SSA extension between SSA and operator SSA.



Future work

- new SSA extension between SSA and operator SSA.
- use operator SSA inequality to characterize quantum state when $SSA=0$.

Thank you

Thank you for listening!

Q&A

Questions?