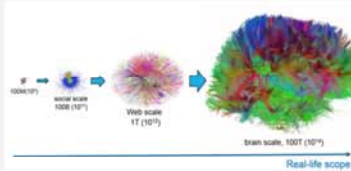


OVERVIEW

Motivation

Querying Big graphs

Given query Q , find answers $Q(G)$ for Q from data graph G



Challenge

- Real graphs are **Huge**.
- Tractable methods could be in feasible!
- Real-world applications require searching with limited resource**
- How to evaluate query using limited resource?

Contributions

Resource-bounded query answering

- Accessing small amount of data for accurate answers
- A tunable graph querying framework
- Balance computing resource and answer quality

Efficient resource bounded algorithms for

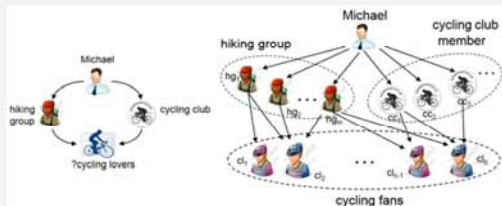
- Localized queries: strong simulation & subgraph queries
- Non-localized queries: reachability

Related Work: BlinkDB, budgeted search, graph indexing & compression, MapReduce, GraphLab...

Graph Queries

Localized graph pattern queries

- A match can be determined by exploring its d_Q (diameter of Q) hops
- Simulation queries: strong simulation relation with a personalized node
- Subgraph isomorphism: injective mapping



Matching relation: (Michael, Michael), (hiking group, hg_m), (cycling club, cc_1), (cycling club, cc_2), (cycling lover, cl_{n-1}), (cycling lover, cl_n)

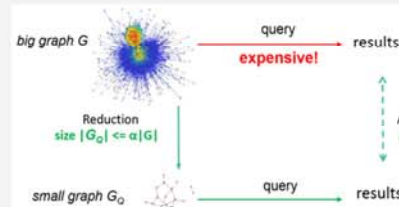
Non-localized queries

- We consider reachability queries: visit the entire G in the worst case

Question

How to effectively evaluate graph pattern queries and reachability with lim-

Resource-bounded Graph Querying



Resource-bounded algorithm for query class with resource bound α and accuracy guarantee η

- Access small (bounded) amount of G
- Guaranteed result quality
- Balance resource and answer quality

Resource-bounded graph query answering
Given: a query class L , α in $(0, 1]$ and η in $(0, 1]$, find algorithm with resource bound α and accuracy guarantee η

NP-hard for simulation and subgraph queries
Impossible for reachability queries

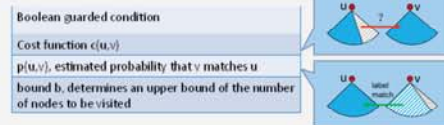
Resource-bounded Graph Simulation

Auxiliary Information (Preprocessing)

Local information that benefits dynamic reduction

degree |neighbor| <label, frequency> ...

Dynamically updated during processing



Dynamic Reduction

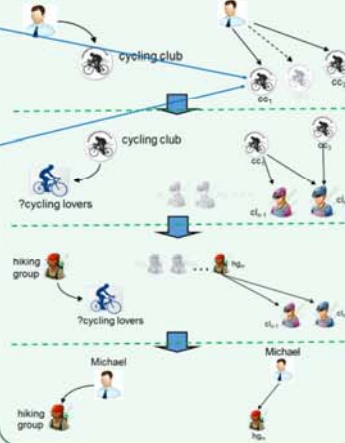
- Iteratively process each query edge
- Dynamically update cost and probability
- Select promising nodes
- Update G_0 until resource bound reached/no unvisited nodes

Approximate Querying

- apply graph simulation algorithms over G_0 to compute matches

Resource bounded graph simulation

Dynamic Reduction

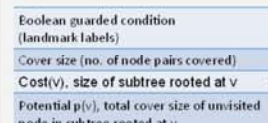


Resource-bounded Reachability

Hierarchical Landmark Index

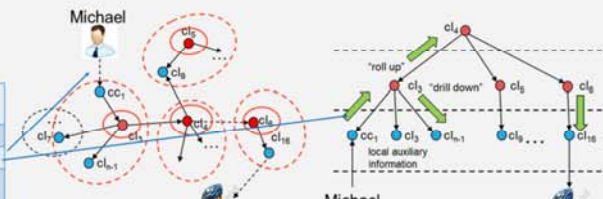
Landmarks are used to encode reachability.
Idea: select bounded number of landmarks for approximate reachability.

Dynamically updated auxiliary information during processing



Guided reachability search

- Bi-directed search with guided "roll-up"/"drill-down"
- Terminates if "yes" is determined or no unvisited nodes in the index



Resource-bounded Querying: Summary

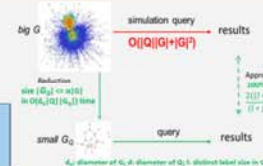
Resource-bounded

Simulation querying/

Subgraph isomorphism

$$\text{accuracy measure} \quad \text{precis} = \frac{|V \cap Q(G)|}{|V|} \quad \text{recall} = \frac{|V \cap Q(G)|}{|Q(G)|}$$

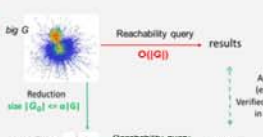
$$\text{acc} = 2 \text{precis} * \text{recall} / (\text{precis} + \text{recall})$$



Resource-bounded reachability

$$\text{accuracy measure} \quad \text{precis} = \frac{tp}{tp + fp} \quad \text{recall} = \frac{tp}{tp + fn}$$

Other types of resource bounds, accuracy measures, graph query classes, applications...



RESULTS

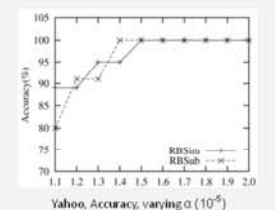
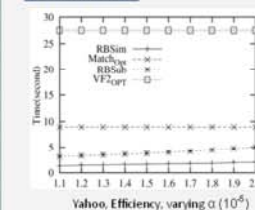
Dataset

- Yahoo Web graph (<http://webscope.sandbox.yahoo.com/catalog.php?datatype=g>)
- Youtube (<http://netsa.cs.sfu.ca/youtubedata>)

Baseline

- Resource bounded reachability **RBReach** VS **LM**: applying landmark vectors; **BFS** and **optimized BFS** over compressed graphs
- Resource bounded simulation algorithm **RBSim** VS **Optimized strong simulation MatchOpt**
- Resource bounded subgraph isomorphism **RBSub** VS **Optimized VF2**

Evaluation



Big Picture

