

# Funktionale Programmierung Mitschrift

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„Avoid success at all cost “

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## Vorlesung 1

```
-- Hello World Haskell
main :: IO ()
main = putStrLn "Chewie, we're home"
```

Codebeispiel 1: Hello World

## Functional Programming (FP)

A programming language is a medium for expressive ideas (not to get a computer to perform operations ). Thus programs must be written for people to read, and only incidentally for machines.

## Computational Model in FP : *Reduction*

Replace expressions by their value.

IN FP, expressions are formed by applying functions to values.

1. Function as in maths:  $x = y \rightarrow f(x) = f(y)$
2. Functions are values like numbers or text

	FP	Imperative
construction	function application and composition	statement sequencing
execution	reduction (expression evaluation)	state changes
semantics	$\lambda$ -calculus	denotational

$n \in \mathbb{N}, n \geq 2$  is a prime number  $\Leftrightarrow$  the set of non-trivial factors of  $n$  is empty.  
 $n$  is prime  $\Leftrightarrow \{m \mid m \in \{2, \dots, n-1\}, n \bmod m = 0\} = \{\}$

```
int IsPrime(int n)
{
    int m;
    int found_factor;
    found_factor
    for (m = 2; m <= n - 1; m++)
    {
        if (n % m == 0)
        {
            found_factor = 1 ;
            break;
        }
    }
    return !found_factor;
}
```

Codebeispiel 2: isPrime in C

```
isPrime :: Integer -> Bool
isPrime n = factors n == []
  where
    factors :: Integer -> [Integer]
    factors n = [ m | m <- [2..n-1], mod n m == 0]

main :: IO ()
main = do
  let n = 42
  print (isPrime n)
```

Codebeispiel 3: isPrime in Haskell

```
let xs = [ x+1 | x <- [0..9] ]
:sprint xs = _
length xs
:sprint xs = [_,_,_,_,_,_,_,_,_,_]
```

Codebeispiel 4: Lazy Evaluation in der ghci REPL

## Haskell Ramp Up

Read  $\equiv$  as "denotes the same value as"

Apply  $f$  to value  $e$ :  $f \sqcup e$

(juxtaposition, "apply", binary operator  $\sqcup$ , Haskell speak:  $\text{infixL } 10 \sqcup$ ) =  $\sqcup$  has

max precedence (10):  $f e_1 + e_2 \equiv (f e_1) + e_2$   $\sqcup$  associates to the left  $g \sqcup f \sqcup e \equiv (g \sqcup f) e$

Function composition:

-  $g (f e)$

- Operator "." ("after") :  $(g.f) e \quad (.\ = \circ) = g(f e)$

- Alternative "apply" operator  $\$$  (lowest precedence, associates to the right),  
 $\text{infix } 0 \$$ :  $f \$ e_1 + e_2 = f (e_1 + e_2)$

## Vorlesung 2

Prefix application of binary infix operator  $\oplus$

$(\oplus) e_1 e_2 \equiv e_1 \oplus e_2$

$(\&\&) \text{ True False} \equiv \text{False}$

Infix application of binary function  $f$ :

$e_1 \text{ `f` } e_2 \equiv f e_1 e_2$

$x \text{ `elem` } xs \equiv x \in xs$

User defined operators with characters : !#%&\*+ / <=> ? @ \ ^ |

```
cos 2 * pi
cos (2 * pi)
cos $ 2 * pi
isLetter (head (reverse ("It's a " ++ "Trap")))
(isLetter . head . reverse) ("It's a " ++ "Trap")
isLetter $ head $ reverse $ "It's a " ++ "Trap"
```

Codebeispiel 5: Verschiedene Schreibweise einer Applikation

```
epsilon :: Double
epsilon = 0.00001
(==~) :: Double -> Double -> Bool
x ==~ y = abs (x - y) < epsilon
infix 4 ==~
```

Codebeispiel 6: Eigener  $\approx$  Operator

## Values and Types

Read `::` as "has type"

Any Haskell value `e` has a type `t` (`e :: t`) that is determined at compile time.

The `::` type assignment is either given explicitly or inferred by the computer

## Types

Type	Description	Value
Int	fixed precision integers ( $-2^{63} \dots 2^{63} - 1$ )	0, 1, 42
Integer	arbitrary Precision integers	0, $10^{100}$
Float, Double	Single/Double precision floating points	0.1, 1e03
Char	Unicode Character	'x', '\t', ' ', '\8710'
Bool	Booleans	True, False
()	Unit (single-value type)	()

```
2
it :: Integer
42 :: Int
it :: Int
'a'
it :: Char
True
it :: Bool
10^100
it :: Integer
10^100 :: Double
it :: Double
```

## Type Constructors

- Build new types from existing Types
- Let a,b denote arbitrary Types (type variables)

Type Constructor	Description	Values
(a,b)	pairs of values of types a and b	(1,True) :: (Int, Bool) 2,False :: (Int, Bool)
(a <sub>1</sub> , a <sub>2</sub> , ..., a <sub>n</sub> )	n-Types	[] :: [a]
[a]	list of values of type a	Just 42 :: Maybe Integer Nothing :: Maybe a
Maybe a	optional value of type a	Left 'x' :: Either Char b Right pi :: Either a Double
Either a b	Choice between values of Type a and b	print 42 :: IO ()
IO a	I/O action that returns a value of type a (can have side effects )	getChar :: IO Char isLetter :: Char -> Bool
a -> b	function from type a to b	

```

(1, '1', 1.0)
it :: (Integer, Char, Double)
[1, '1', 1.0]
it :: Fehler
[0.1, 1.0, 0.01]
it :: [Double]
[]
it :: [t]
"Yoda"
it :: [Char]
['Y', 'o', 'd', 'a']
"Yoda"
[Just 0, Nothing, Just 2]
it :: [Maybe Integer]
[Left True, Right 'a']
it :: [Either Bool Char]
print 'x'
it :: ()
getChar
*
it :: Char
:t getChar
getChar :: IO Char
:t fst
fst :: (a,b) -> a
:t snd
snd :: (a,b) -> b
:t head
head :: [a] -> a
:t (++)
(++) :: [a] -> [a] -> [a]

```

## Currying

- Recall:
  - $e_1 ++ e_2 \equiv (++) e_1 e_2$
  - $++ e_1 e_2 \equiv ((++) e_1) e_2$
- Function application happens one argument at a time (currying, Haskell B. Curry)
- Type of n-ary function:  $: a_1 -> a_2 \dots -> a_n -> b$
- Type constructor  $->$  associates to the right thus read the type as:  
 $a_1 -> (a_2 -> a_3 (\dots -> (a_n -> b) \dots))$

- Enables partial application: "Give me a value of type  $a_1$ , I'll give you a (n-1)-ary function of type  $a_2 \rightarrow a_3 \rightarrow \dots \rightarrow a_n \rightarrow b$ "

```
"Chew" ++ "bacca"
"Chewbacca"
(++) "Chew" "bacca"
"Chewbacca"
((++) "Chew") "bacca"
"Chewbacca"
:t (++) "Chew"
"Chew" :: [Char] -> [Char]
let chew = (++) "Chew"
chew "bacca"
"Chewbacca"
let double (*) 2
double 21
42
```

## Vorlesung 3

### Defining Values (and thus: Functions)

- = binds names to values, names must not start with A-Z (Haskell style: camelCase)
- Define constant (0-ary) c, value of c is that of expression:  
 $c = e$
- Define n-ary function, arguments  $x_i$  and f may occur in e (no "letrec" needed)  
 $f\ x_1\ x_2 \dots x_n = e$
- Hskell programm = set of top-level bindings (order immaterial, no rebinding)
- Good style: give type assignment for top-level bindings:  
 $f :: a1 \rightarrow a2 \rightarrow b$   
 $f\ x_1\ x_2 = e$
- Guards (introduced by |).

```
f x1 x2 ... xn
  | q1 = e1
  | q2 = en
```

- $q_i$  (expressions of type Bool) evaluated top to bottom, first True guards "wins"
- $$\text{fac } n = \begin{cases} 1 & \text{if } n \geq 1 \\ n \cdot \text{fac}(n-1) & \text{else} \end{cases}$$

```
fac :: Integer -> Integer
fac n = if n <= 1 then 1 else n * fac (n - 1)

fac2 n | n <= 1    = 1
      | otherwise = n * fac2 (n - 1)

main :: IO ()
main = print $ fac 10
```

Codebeispiel 7: fac in Haskell

```
power :: Double -> Integer -> Double
power x k | k == 1 = x
          | even k = power (x * x) (halve k)
          | otherwise = x * power (x * x) (halve k)
where
  even :: Integer -> Bool -- Nicht typisch
  even n = n `mod` 2 == 0
  halve n = n `div` 2

main :: IO ()
main = print $ power 2 16
```

Codebeispiel 8: Power in Haskell

## Lokale Definitionen

1. **where** - binding : Local definitions visible in the entire right-hand-side (rhs) of a definition

```
f x1 x2 ... xn
  | q1 = e1
  | q2 = en
where
  g1 ... = b1
  gi ... = bi
```

2. **let** - expression Local definitions visible inside an expression:

```
let g1 ... = b1
    g2 ... = b1
in e
```

## Lists([a])

- Recursive definition:

1. **[]** ist a list (nil), type **[] :: [a]**



2.  $x : xs$  (head, tail) is a list, if  $x :: a$ , and  $xs :: [a]$ .

`cons: (:) :: a -> [a] -> [a] -> [a]` with `infixr : 5`

- Notation:  $3:(2:1:[]) \equiv 3:2:1:[] \equiv [3,2,1]$

```
[ ]
it :: [t]
[1]
it :: [Integer]
[1,2,3]
it :: [Integer]
['z']
"z"
it :: [Char]
['z','x']
"zx"
it :: [Char]
[] == ""
True
it :: Bool
[[1],[2,3]]
it :: [[Integer]]
[[1],[2,3],[]]
[[1],[2,3]]
it :: [[Integer]]
False:[]
[False]
it :: [Bool]
(False:[]):[]
it :: [[Bool]]
:t [(<),(=<),(>)]
[(<),(=<),(>)] :: Ord a => [a -> a-> Bool]
[(1,"one"),(2,"two"),(3,"three")]
it :: [(Integer,[Char])]
:t head
head :: [a] -> a
:t tail :: [a] -> [a]
head "It's a trap"
'I'
it :: Char
tail "It's a trap"
"t's a trap"
it :: [Char]
reverse "Never odd or even"
"neve ro ddo reveN"
it :: [Char]
```

- Law  $\forall xs \neq []: \text{head } xs : \text{tail } xs$

```
:i String
type String = [Char]
```

## Type Synonyms

- Introduce your own type synonyms. (type names : *Uppercase*) `type  $t_1 = t_2$`

```
type Bits = [Integer]
```

```
type Predicate a = a -> Bool
```

```
bits :: Integer -> Bits
bits n | n == 0      = [0]
       | otherwise = (n `mod` 2) : bits (n `div` 2)
```

```
isEven :: Predicate Integer
isEven n = head (bits n) == 0
```

```
main :: IO ()
main = print $ isEven 35
```

Sequence (lists of enumerable elements)

- $[x..y] \equiv [x, x+1, x+2, \dots, y]$

```
['a'..'z']
"abcdefghijklmnopqrstuvwxyz"
```

- $x, s..y \equiv [x, x+i, x+(2*i), \dots, y]$  where  $i = x-s$

```
[1,3..20]
[1,3,5,7,9,11,13,15,17,19]
[2,4..20]
[2,4,6,8,10,12,14,16,18,20]
```

- Infinite List `[1..]`