

# CMPS 431 W - Homework 2.

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1. Declarative query language = describe what you want rather than how you compute it. Relational calculus query is declarative query language.

Procedural query language = used to represent execution plans  
Relational algebra is procedural query language.

2. Assume eid are unique.

(1.) relational algebra =

$$\pi_{fname, lname, branch\_name, num\_employee} \left( \left( \text{Branch} \bowtie_{\text{Branch.eid} = \text{Personal Banker.eid}} \text{Personal Banker} \right) \bowtie \text{Employee} \right)$$

TRC =

$$\{R \mid \exists e \in \text{Employee}, \exists p \in \text{Personal Banker}, \exists b \in \text{Branch} ( e.eid = p.eid \wedge \\ b.manager\_eid = e.eid \wedge R.eid = p.eid \wedge R.fname = e.fname \wedge \\ R.lname = e.lname \wedge R.num\_employee = b.num\_employee \wedge \\ R.branch\_name = b.branch\_name ) \}$$

(2.) relational algebra =

$$\pi_{fname, lname} ( \sigma_{\text{birthday} > "06/02/2001" \wedge \text{income} > 10000} (\text{Customer}) )$$

TRC =

$$\{R \mid \exists p \in \text{Customer} ( p.birthday > "06/02/2001" \wedge p.fname = R.fname \wedge \\ R.lname = p.lname \wedge p.birthday = R.birthday \wedge p.income > 10000 ) \}$$

(3.) relational algebra:

$$\pi_{cid, fname, lname, birthday} (\forall num\_employee > 25 (Customer \bowtie Duns \bowtie Account \bowtie Branch))$$

TRC =

$$\{ R \mid \exists D \in Duns, \exists A \in Account, \exists B \in Branch, \exists C \in Customer (B.num\_employee > 25 \wedge A.accno = D.accno \wedge A.branch\_no = B.branch\_no \wedge D.cid = C.cid \wedge R.cid = D.cid \wedge R.fname = C.fname \wedge R.lname = C.lname \wedge R.birthday = C.birthday) \}$$

(4.) relational algebra:

$$\pi_{cid}(Duns) - \pi_{cid}(\forall Duns.cid \neq d.cid \wedge Duns.accno = d.accno (P_d(Duns) \times Duns))$$

TRC =

$$\{ R \mid \forall A \in Duns, \nexists B \in Duns (A.accno = B.accno \wedge A.cid \neq B.cid \wedge R.cid = B.cid) \}$$

(For partial) Owners that DO have joint account:

$$A.accno = B.accno \wedge A.cid \neq B.cid$$

same account number, but not same cid (different person).

(5.) relational algebra: # I considered "Pittsburgh West" as branch-name. But I think it can also be street

$$\pi_{cid}((\forall branch\_name = "Pittsburgh West" (Branch \bowtie Employee \bowtie \overset{\text{Banker}}{\text{Personal}}))) \cap (\forall branch\_name = "Harrisburgh South" (Branch \bowtie Employee \bowtie \overset{\text{Banker}}{\text{Personal}}))$$

TRC = same cid (unique).  $P_1, P_2$  should be same person  $\begin{cases} \textcircled{1} P_1 \text{ in Pitts.} \\ \textcircled{2} P_2 \text{ in Harris.} \end{cases}$

$$\{ R \mid \exists C \in Customer, \exists P_1 \in \text{Personal Banker}, \exists P_2 \in \text{Personal Banker}, \exists E_1 \in Employee, \exists E_2 \in Employee, \exists B_1 \in Branch, \exists B_2 \in Branch (C.cid = P_1.cid \wedge P_1.cid = P_2.cid \wedge P_1.eid = E_1.eid \wedge E_1.branch\_no = B_1.branch\_no \wedge B_1.branch\_name = "Pittsburgh West" \wedge P_2.eid = E_2.eid \wedge E_2.branch\_no = B_2.branch\_no \wedge B_2.branch\_name = "Harrisburgh South" \wedge R.cid = C.cid) \}$$

get the cids as required.

(6.) relational algebra:

$$\pi_{\text{branch\_name}} (\sigma_{\text{name} = \text{"Jackson"}} (\text{Employee} \bowtie \text{Branch}) \cap \sigma_{\text{salary} < 20000} (\text{Employee} \bowtie \text{Branch}))$$

TRC =

$$\{ R \mid \exists E_1 \in \text{Employee}, \exists E_2 \in \text{Employee}, \exists B \in \text{Branch} (E_1.\text{branch\_no} = E_2.\text{branch\_no} \wedge E_1.\text{name} = \text{"Jackson"} \wedge E_2.\text{salary} < 20000 \wedge E_1.\text{branch\_no} = B.\text{branch\_no} \wedge R.\text{branch\_name} = B.\text{branch\_name}) \}$$

work in the same branch.      no matter if  $E_1 = E_2$ .

(7.) relational algebra:

$$\pi_{\text{cid}} (\sigma_{\text{txn\_type} = \text{"check"} \wedge \text{amount} > 10000} (\text{Transaction} \bowtie \text{Ouns}))$$

absolute value.

TRC =

$$\{ R \mid \exists T \in \text{Transaction}, \exists O \in \text{Ouns} (T.\text{txn\_type} = \text{"check"} \wedge T.\text{accno} = O.\text{accno} \wedge |T.\text{amount}| > 10000 \wedge R.\text{cid} = O.\text{cid}) \}$$

absolute value

(8.) relational algebra:

$$\pi_{\text{fname}, \text{lname}, \text{birthday}} (\sigma_{\text{birthday} < \text{"08/21/1950"}} (\text{Customer}))$$

select all customer meet the birthday requirement.

$$\sigma_{A.\text{income} > B.\text{income}} (\rho_A (\sigma_{\text{birthday} < \text{"08/21/1950"}} (\text{Customer})) \times \rho_B (\sigma_{\text{birthday} < \text{"08/21/1950"}} (\text{Customer})))$$

all possible combination of customer

(8.) TRC =

$$\{R \mid \exists C_1 \in \text{Customer}, \forall C_2 \in \text{Customer} (C_1.\text{birthday} < "08/21/1950" \wedge C_2.\text{birthday} < "08/21/1950" \wedge C_1.\text{income} < C_2.\text{income} \wedge R.\text{fname} = C_1.\text{fname} \wedge R.\text{lname} = C_1.\text{lname} \wedge R.\text{birthday} = C_1.\text{birthday})\}$$

select all (by t)  
customer with  
required birthday.

(9.) relational algebra =

$$\pi_{\text{cid}, \text{salary}} (\sigma_{\text{Employee.salary} > M.\text{salary}} (\text{Employee} \bowtie (\rho_M (\text{Employee} \bowtie \text{Branch}))))$$

$\bowtie_{\text{eid} = \text{manager\_eid}} \text{Employee}$

TRC = regular employee

manager

$$\{R \mid \exists E \in \text{Employee}, \exists M \in \text{Employee}, \exists B \in \text{Branch} (E.\text{branch\_no} = M.\text{branch\_no} \wedge M.\text{eid} = B.\text{manager\_eid} \wedge E.\text{salary} > M.\text{salary} \wedge R.\text{eid} = E.\text{eid} \wedge R.\text{salary} = E.\text{salary})\}$$

10. relational algebra =

get full info

get "real" id.

get id of "fake info".

all possible combination (with wrong info)

$$\pi_{\text{fname}, \text{lname}, \text{income}} ((\pi_{\text{cid}} (\text{Customer}) - \pi_{\text{cid}} (\pi_{\text{cid}, \text{branch\_no}} (\text{Customer} \times \text{Branch}))) \bowtie \pi_{\text{cid}, \text{branch\_no}} (\text{Ouns} \bowtie \text{Account})) \bowtie \text{Customer}$$

that actual have the account

"fake info".

TRC =

$$\{R \mid \exists C \in \text{Customer}, \exists D \in \text{Ouns}, \exists A \in \text{Account}, \forall B \in \text{Branch} (C.\text{cid} = D.\text{cid} \wedge D.\text{accno} = A.\text{aceno} \wedge A.\text{branch\_no} = B.\text{branch\_no} \wedge R.\text{fname} = C.\text{fname} \wedge R.\text{lname} = C.\text{lname} \wedge R.\text{income} = C.\text{income})\}$$

11. relational algebra:

$$\pi_{cid}(Customer) - \pi_{cid} \left( \pi_{cid, branch\_no} \left( \frac{Customer \times (\sigma_{20000 < budget < 200000}(Branch))}{\downarrow \text{all possible combination}} \right) - \left( \pi_{cid, branch\_no}(Duns \bowtie Account) \right) \right)$$

all cid + branch-no that do exist

TRC =

$$\{ R \mid \exists C \in Customer, \exists D \in Duns, \exists A \in Account, \forall B \in Branch ( \underbrace{(20000 < B.budget < 200000 \rightarrow (C.cid = D.cid \wedge D.accno = A.accno \wedge A.branch\_no = B.branch\_no))}_{\text{all possible combination}} ) \wedge R.cid = C.cid ) \}$$