

# OPERATING SYSTEM

## Page Table Entries (PTE)

### Definition

A **Page Table Entry (PTE)** is a **single record in the page table** that stores information about **where a page of a process is located in physical memory** and **how it can be accessed**.

Each page in a process has **one corresponding page table entry**.

### What Does a Page Table Entry Contain?

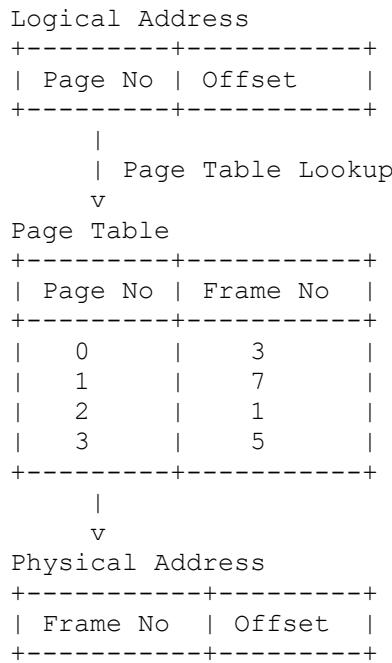
A typical **PTE** includes the following fields:

1. **Frame Number**
  - Indicates **which physical frame** holds the page
2. **Valid / Invalid Bit**
  - 1 → Page is present in memory
  - 0 → Page is not present (page fault)
3. **Protection Bits**
  - Read / Write / Execute permissions
4. **Modified (Dirty) Bit**
  - 1 → Page has been modified
  - 0 → Page not changed
5. **Reference (Accessed) Bit**
  - Set when the page is accessed
  - Used in page replacement algorithms

### Simple Page Table Entry Format

Frame No	Valid	R/W/X	Dirty	Reference

### Paging with Page Table – Diagram



## Step-by-Step Working

1. CPU generates a **logical address**
2. Logical address is divided into:
  - a. **Page number**
  - b. **Offset**
3. Page number indexes the **page table**
4. Page table entry gives the **frame number**
5. Frame number + offset forms the **physical address**

## Why Page Table Entries Are Important

- Enable **address translation**
- Support **memory protection**
- Help detect **page faults**
- Used in **page replacement algorithms**

**Offset** is the part of a logical address that specifies the exact location within a page. During address translation, the offset remains unchanged and is combined with the frame number to form the physical address.

A **page table entry (PTE)** stores information about a page, including its **frame number, validity, protection bits, and status bits**. It is used by the operating system to translate logical addresses into physical addresses and manage memory efficiently.