$\begin{tabular}{ll} National Taiwan University \\ Ashes \end{tabular}$

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1 Basic

1.1 Fast Integer Input

```
|inline int gtx() {
  const int N = 4096;
  static char buffer[N];
  static char *p = buffer, *end = buffer;
  if (p == end) {
    if ((end = buffer + fread(buffer, 1, N, stdin)) == buffer)
         return EOF;
    p = buffer;
  return *p++;
template <typename T>
inline bool rit(T& x) {
  char c = 0; bool flag = false;
  while (c = getchar(), (c < '0' && c != '-') || c > '9') if (c
       == -1) return false;
  c == '-' ? (flag = true, x = 0) : (x = c - '0');
  while (c = getchar(), c >= '0' && c <= '9') x = x * 10 + c -
  if (flag) x = -x;
  return true;
```

1.2 Increase stack size

```
| const int size = 256 << 20;
| register long rsp asm("rsp");
| char *p = (char*)malloc(size) + size, *bak = (char*)rsp;
| _asm_("movq %0, %%rsp\n"::"r"(p));
| // main
| _asm_("movq %0, %%rsp\n"::"r"(bak));
```

1.3 Default Code

```
#i ncl ude <bits/stdc++. h>
#pragma GCC optimize ("O3")
#pragma GCC target ("sse4")
using namespace std;
#define DBG(x) cout << (#x " = ") << x << endl;
#defi ne ALL(x) x. begi n(), x. end()
#define push_back emplace_back
 #define spl O i os::sync_with_stdio(false);cin.tie(0)
#define mem(arr, value) memset(arr, value, sizeof(arr))
#define for O(i, n) for (int i = 0; i < n; i++)
#define for 1(i, n) for (int i = 1; i \le n; i++)
typedef long long II;
typedef long double ld;
const II mod = 1e9 + 7;
inline II pow2(II target, II p, II MOD = mod)
     II ret = 1;
     while (p)
         if (p & 1)
              (ret *= target) %= MOD;
         p >>= 1:
         (target *= target) %= MOD;
     return ret;
inline II inv(II x, II MDD = mod)
{
     return pow2(x, MDD - 2, MDD);
}
inline II gcd(II x, II y)
{
     if (! y)
         return x:
     if ((x & 1) && (y & 1))
     {
         if (x < y)
              swap(x, y);
         return gcd(y, x - y);
     if (x & 1)
         return gcd(x, y \gg 1);
     if (y & 1)
     return gcd(y, x >> 1);
return 2 * gcd(x >> 1, y >> 1);
inline II sub(II x, II y, II MDD = mod)
| {
     return (x - y < 0 ? x - y + MDD : x - y);
inline II add(II x, II y, II MOD = mod)
{
     return (x + y >= MDD ? x + y - MDD : x + y);
vector<pair<int, int>> directions = {
       \{\,1,\,0\}\,,\,\{\,1,\,1\}\,,\,\{\,0,\,1\}\,,\,\{\,-\,1,\,1\}\,,\,\{\,-\,1,\,0\}\,,\,\{\,-\,1,\,-\,1\}\,,\,\{\,0,\,-\,1\}\,,\,\{\,1,\,-\,1\}\quad \}\,;
```

1.4 Big Integer

```
#i ncl ude <bi ts/stdc++. h>
struct Int {
  static const int inf = 1e9;
  std::vector<int> di g;
  bool sgn;
  Int() {
    di g. push_back(0);
    sgn = true;
  Int(int n) {
    sgn = n >= 0;
    while (n) {
       dig.push_back(n % 10);
       n /= 10:
    if (dig. size() == 0) dig. push_back(0);
  Int(std::string s) {
    int i = 0; sgn = true;
if (s[i] == '-') sgn = false, ++i;
    for (; i < s.length(); ++i) dig. push_back(s[i] - '0');
    reverse(dig. begin(), dig. end());
    if (dig.size() == 1 && dig[0] == '0') sgn = true;
  Int(const std::vector<int>& d, const bool & s = true) {
```

```
dig = std::vector<int>(d. begin(), d. end());
  sgn = s;
Int(const Int& n) {
  sgn = n. sgn;
  dig = n. dig;
bool operator<(const Int&rhs) const {</pre>
  if (sgn && !rhs. sgn) return true;
  if (!sgn && rhs.sgn) return false;
  if (!sgn && !rhs.sgn) return Int(dig) > Int(rhs.dig);
  if (dig. size() < rhs. dig. size()) return true;</pre>
  if (dig. size() > rhs. dig. size()) return false;
  for (int i = dig. size() - 1; i >= 0; --i) {
   if (dig[i] != rhs. dig[i]) return dig[i] < rhs. dig[i];</pre>
  return false;
bool operator ==(const Int& rhs) const {
  if (sgn != rhs.sgn) return false;
  return dig == rhs. dig;
bool operator>(const Int& rhs) const {
  return ! (*this < rhs) && ! (*this == rhs);
bool operator<(const int& n) const {</pre>
 return *this < Int(n);</pre>
bool operator > (const int& n) const {
  return *this > Int(n);
bool operator == (const int & n) const {
  return *this == Int(n);
Int operator-() const {
  return Int(dig, !sgn);
Int operator+(const Int& rhs) const {
  bool res = true;
  if (!sgn && !rhs.sgn) res = false;
  else if (!sgn && rhs.sgn) return rhs - (-*this);
  else if (sgn &&!rhs.sgn) return *this - -rhs;
  std: vector < i nt > v1 = di g, v2 = rhs. di g;
  if (v2. size() > v1. size()) swap(v1, v2);
  int car = 0;
  std::vector<int> nvec;
  for (int i = 0; i < v2. size(); ++i) {
    int k = v1[i] + v2[i] + car;
    nvec. push_back(k % 10);
    car = k / 10
  for (int i = v2. size(); i < v1. size(); ++i) {
    int k = v1[i] + car
    nvec. push_back(k % 10);
car = k / 10;
  return Int(nvec, res);
Int operator-(const Int& rhs) const {
  if (*this < rhs) {</pre>
    std::vector<int> nvec = (rhs - *this).dig;
    return Int(nvec, false);
  if (*this == rhs) return Int(0);
  std: vector < i nt > v1 = di g, v2 = rhs. di g;
  std::vector<int> nvec;
  for (int i = 0; i < v2. size(); ++i) {
    int k = v1[i] - v2[i];
    if (k < 0) {
      for (int j = i + 1; j < v1. size(); ++j) if (v1[j] > 0)
         --v1[j]; k += 10;
        break;
      }
    nvec. push_back(k);
  int rind = v1. size() - 1;
  while (rind \Rightarrow v2. size() && v1[rind] \Rightarrow 0) --rind;
  for (int i = v2. size(); i \le rind; ++i) {
    nvec. push_back(v1[i]);
  return Int(nvec);
Int operator*(const Int& rhs) const {
```

```
if (*this == 0) return Int():
     if (rhs == 0) return Int();
     std: vector < i nt > v1 = di g, v2 = rhs. di g;
     if (v1. size() < v2. size()) swap(v1, v2)
     std::vector<int> res(v1. size() * v2. size(), 0);
     for (int i = 0; i < v2. size(); ++i) {
       int car = 0;
       for (int j = 0; j < v1. size(); ++j) {
         int k = car + v1[j] * v2[i];
         res[j + i] += k \% 10;
         car = k / 10;
       }
     int car = 0;
     for (int i = 0; i < res. size(); ++i) {
       int k = car + res[i];
       res[i] = k \% 10;
       car = k / 10;
     while (car) {
       res. push_back(car % 10);
       car /= 10;
     int ind = res.size() - 1;
     while (ind \Rightarrow 0 && res[ind] == 0) --ind;
     std:: vector <i nt > nvec;
     for (int i = 0; i <= ind; ++i) nvec.push_back(res[i]);</pre>
     return Int(nvec);
   Int operator+(const int& n) const {
    return *this + Int(n);
   Int operator-(const int& n) const {
     return *this - Int(n);
   Int& operator += (const Int& n) {
     *this = (*this + n);
     return *this;
   Int& operator - = (const Int& n) {
     *this = (*this - n);
     return *this;
   Int& operator+=(const int& n) {
     *this += Int(n);
     return *this;
   Int& operator - = (const int& n) {
     *this -= Int(n);
     return *this;
   Int& operator*=(const Int& n) {
     *this = *this
     return *this;
   Int& operator*=(const int& n) {
     *this *= Int(n);
     return *this;
   Int& operator++(int) {
     * this += 1;
     return *this:
   Int& operator--(int) {
     *this -= 1;
     return *this;
   friend std::istream& operator>>(std::istream&in, Int&n) {
     std::string s; in >> s;
     n = Int(s);
     return in;
   friend std::ostream& operator<<(std::ostream& out, const Int&
     if (! n. sgn) out << "-";
     for (int i = n. dig. size() - 1; i >= 0; --i) out << n. dig[i
     return out;
  }
};
```

2 Flows, Matching

2.1 Dinic's Algorithm

```
struct Edge {
  int to, cap, rev;
  Edge(int a, int b, int c) : to(a), cap(b), rev(c) {}
int Flow(vector<vector<Edge>> &g, int s, int t){
 int n = g. size(), res = 0;
vector<int> lev(n, -1), iter(n);
  while(true){
    vector < i nt > que(1, s);
    fill(lev.begin(), lev.end(), -1);
    fill(iter.begin(), iter.end(), 0);
    lev[s] = 0;
    for (int i = 0; i < (int) que. size(); i++){
      int x = que[i];
      for (Edge &e : g[x]){
        if (e. cap > 0 && lev[e. to] == -1){
          lev[e.to] = lev[x] + 1;
           que. push_back(e.to);
      }
    }
    if (lev[t] == -1) break;
    auto Dfs = [\&] (auto dfs, int x, int f = 1e9) {
      if (x == t) return f;
      int res = 0;
      for (int &i = iter[x]; i < (int)g[x].size(); i++){
        Edge &e = g[x][i];
        if (e. cap > 0 \&\& lev[e. to] == lev[x] + 1){
           int p = dfs(dfs, e.to, min(f - res, e.cap));
          res += p;
          e. cap -= p;
          g[e. to][e. rev]. cap += p;
      if (res == 0) lev[x] = -1;
      return res;
    }:
    res += Dfs(Dfs, s);
  return res;
auto Add = [&](int a, int b, int c){
  g[a].emplace_back(b, c, (int)g[b].size());
  g[b].emplace_back(a, 0, (int)g[a].size()-1);
```

2.2 Minimum Cost Maximum Flow

```
struct Edge {
 int to, cap, rev, w,
  Edge(int t, int c, int r, int w) : to(t), cap(c), rev(r), w(w)
pair<int, int> Flow(vector<vector<Edge>> g, int s, int t) {
 int N = g. size();
  vector < int > dist(N), ed(N), pv(N);
  vector <bool > i nque(N);
  int flow = 0, cost = 0;
  while (true) {
   dist.assign(N, klnf);
    inque.assign(N, false);
    pv. assi gn(N, -1);
    dist[s] = 0;
    queue<int> que;
    que. push(s);
    while (!que.empty()) {
      int x = que.front(); que.pop();
      inque[x] = false;
      for (int i = 0; i < g[x].size(); ++i) {
        Edge &e = g[x][i];
        if (e.cap > 0 && dist[e.to] > dist[x] + e. w) {
          dist[e.to] = dist[x] + e.w
          pv[e.to] = x;
          ed[e.to] = i;
          if (!inque[e.to]) {
            inque[e.to] = true;
            que. push(e. to);
         }
       }
      }
    if (dist[t] == kInf) break;
    int f = klnf;
    for (int x = t; x != s; x = pv[x]) f = min(f, g[pv[x]][ed[x])
         ]].cap);
    for (int x = t; x != s; x = pv[x]) {
      Edge &e = g[pv[x]][ed[x]];
```

```
e. cap -= f:
      g[e.to][e.rev].cap += f;
    flow += f;
    cost += f * dist[t];
 return make_pair(flow, cost);
auto AddEdge = [&](int from int to, int cap, int weight) {
  g[from].emplace_back(to, cap, g[to].size(), weight);
  g[to].emplace_back(from 0, g[from].size() - 1, -weight);
```

2.3Bipartite Matching

```
class matching {
 public:
  vector < vector < i nt >> g;
  vector <i nt > pa;
  vector <i nt > pb;
  vector <i nt > was;
  int n, m
  int res;
  int iter;
  matching(int _n, int _m) : n(_n), m(_m)  {
    assert(0 <= n && 0 <= m);
    pa = vector < i nt > (n, -1);
    pb = vector < i nt > (m - 1)
    was = vector < i nt > (n, 0);
    g. resize(n);
    res = 0;
    iter = 0:
  void add(int from int to) {
    assert(0 \le from \&\& from < n \&\& 0 \le to \&\& to < m);
    g[from].push_back(to);
  bool dfs(int v) {
    was[v] = iter;
    for (int u : g[v]) {
       if (pb[u] == -1) {
         pa[v] = u;
         pb[u] = v;
         return true;
      }
    for (int u : g[v]) {
       \begin{tabular}{ll} if & (was[pb[u]] != iter && dfs(pb[u])) & ( \end{tabular} 
         pa[v] = u;
         pb[u] = v;
         return true,
      }
    }
    return false;
  }
  int solve() {
    while (true) {
      iter++;
       int add = 0;
       for (int i = 0; i < n; i ++) {
         if (pa[i] == -1 \&\& dfs(i)) {
           add++;
       if (add == 0) {
         break;
       res += add;
    return res:
  int run_one(int v) {
    if (pa[v] != -1) {
      return O;
    iter++;
    return (int) dfs(v);
```

Weighted Bipartite Match

```
const int klnf = 1e9:
long long KuhnMunkres(vector<vector<int>>> W) {
 int N = Wsize();
  vector < int > fl(N, -1), fr(N, -1), hr(N), hl(N);
 for (int i = 0; i < N; ++i) {
   hl[i] = *max_element(W[i].begin(), W[i].end());
 auto Bfs = [\&](int s) {
    vector <i nt > sl k(N, kl nf), pre(N);
    vector <bool > vI (N, false), vr(N, false);
    queue<i nt > que;
    que. push(s);
    vr[s] = true;
    auto Check = [\&] (int x) -> bool {
     if (vl[x] = true, fl[x] != -1) {
        que. push(fl[x]);
        return vr[fl[x]] = true;
      while (x != -1) swap(x, fr[fl[x] = pre[x]]);
     return false;
    };
    while (true) {
      while (!que.empty()) {
        int y = que.front(); que.pop();
        for (int x = 0, d = 0; x < N; ++x) {
          if (!v[x] \& s[k[x] >= (d = h[x] + hr[y] - W[x][y])
            if (pre[x] = y, d) slk[x] = d;
            else if (!Check(x)) return;
          }
       }
     }
      i nt d = kl nf;
      for (int x = 0; x < N; ++x) {
       if (!vl[x] \&\& d > slk[x]) d = slk[x];
      for (int x = 0; x < N; ++x) {
       if (vl[x]) hl[x] += d;
        else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < N; ++x) {
        if (!vl[x] && !slk[x] && !Check(x)) return;
   }
 }:
 for (int i = 0; i < N; ++i) Bfs(i); long long res = 0;
 for (int i = 0; i < N; ++i) res += W[i][f[[i]];
  return res;
```

Maximum Matching on General Graph

```
const int kN = 500;
namespace matching {
 int fa[kN], pre[kN], match[kN], s[kN], v[kN];
 vector < i nt > g[kN];
 queue<i nt > q;
 void Init(int n) {
    for (int i = 0; i \le n; ++i) match[i] = pre[i] = n;
    for (int i = 0; i < n; ++i) g[i].clear();
 }
 voi d AddEdge(int u, int v) {
    g[u].push_back(v);
    g[v].push_back(u);
 int Find(int u) {
    return \ u == fa[u] \ ? \ u : \ fa[u] = Find(fa[u]); 
 int LCA(int x, int y, int n) {
    static int tk = 0;
    x = Find(x), y = Find(y);
   for (; ; swap(x, y)) {
  if (x != n) {
       if (v[x] == tk) return x;
        v[x] = tk:
        x = Find(pre[match[x]]);
      }
   }
 void Blossom(int x, int y, int l) {
    while (Find(x) != I)
      pre[x] = y, y = match[x];
      if (s[y] == 1) q. push(y), s[y] = 0;
```

```
if (fa[x] == x) fa[x] = I;
if (fa[y] == y) fa[y] = I;
       x = pre[y];
     }
   bool Bfs(int r, int n) {
     for (int i = 0; i <= n; ++i) fa[i] = i, s[i] = -1;
     while (!q.empty()) q.pop();
     q. push(r);
     s[r] = 0;
     while (!q.empty()) {
       int x = q. front(); q. pop();
        for (int u : g[x]) {
  if (s[u] == -1) {
            pre[u] = x, s[u] = 1;
            if (match[u] == n) {
  for (int a = u, b = x, last; b!= n; a = last, b =
                   pre[a])
                last = match[b], match[b] = a, match[a] = b;
              return true;
            q. push(match[u]);
            s[match[u]] = 0
          else if (!s[u] \&\& Find(u) != Find(x)) {
            int I = LCA(u, x, n);
            Blossom(x, u, l);
            Blossom(u, x, 1);
         }
       }
     }
     return false:
   int Sol ve(int n) {
     int res = 0;
     for (int x = 0; x < n; ++x) {
       if (match[x] == n) res += Bfs(x, n);
     return res
   }
}
```

2.6 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x, y, l, u), connect $x \to y$ with capacity u l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v).
 - To maximize, connect $t\to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution.
 - Otherwise, the maximum flow from s to t is the answer To minimize, let f be the maximum flow from S to T. nect $t \to s$ with capacity ∞ and let the flow from S to Tbe f'. If $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
 - 1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise. 2. DFS from unmatched vertices in X.

 - $\begin{array}{ll} 3. & x \in X \text{ is chosen iff } x \text{ is unvisited.} \\ 4. & y \in Y \text{ is chosen iff } y \text{ is visited.} \end{array}$
- $\bullet \quad \text{Maximum density induced subgraph} \\$
 - 1. Binary search on answer, suppose we're checking answer T
 - 2. Construct a max flow model, let K be the sum of all weights 3. Connect source $s\to v,\ v\in G$ with capacity K

 - 4. For each edge (u, v, w) in G, connect $u \to v$ and $v \to u$ with capacity w
 - 5. For $v \in G$, connect it with sink $v \to t$ with capacity K + 2T - $\left(\sum_{e \in E(v)} w(e)\right) - 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u,v).
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v, t) with capacity $-p_v$
 - 2. Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.

• 0/1 quadratic programming

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with capacity c_y . 2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

Data Structure

Disjoint Set Union

```
struct DSU {
  vector < i nt > p, sz;
  int no
  vector <pai r <i nt, i nt >> ops;
  int t:
  DSU(i nt _n): n(_n) {
    p. resi ze(n);
     sz. resi ze(n);
     for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;
  int find(int x) {
    if (p[x] == x) return x;
     else return p[x] = find(p[x]);
  void unite(int x, int y) {
     x = find(x); y = find(y);
     if (x == y) return;
     if (sz[x] > sz[y]) swap(x, y);
     ops. empl ace_back(x, p[x]);
     p[x] = y;
     ops. empl ace_back(~y, sz[y]);
     sz[y] += sz[x];
  voi d save(){
    t = ops. size();
  }
  voi d rol I back(){
     while ((int) ops. size() > t){
       int i = ops.back().first;
       int j = ops. back(). second;
       ops.pop_back();
       if (i >= 0){
         p[i] = j;
       } el se{
         SZ[\sim i] = j;
       }
  }
| };
```

3.2Segment Tree

```
templ ate <cl ass T > struct Seg { // comb(ID, b) = b
  const T ID = 1e18; T comb(T a, T b) { return min(a, b); }
   int n; vector<T> seg;
   void init(int _n) { n = _n; seg. assign(2*n, ID);
   \label{eq:comb} \mbox{void pull(int p) } \{ \mbox{ seg[p] = comb(seg[2*p],seg[2*p+1]); } \}
   void upd(int p, T val) { // set val at position p
  seg[p += n] = val; for (p /= 2; p; p /= 2) pull(p); }
   T query(int I, int r) \{ // min on interval [I, r] \}
       T ra = ID, rb = ID;
      for (I += n, r += n+1; I < r; I /= 2, r /= 2) {
         if (I&1) ra = comb(ra, seg[I++]);
         if (r\&1) rb = comb(seg[--r],rb);
      return comb(ra, rb);
   }
| };
```

Lazy Segment Tree

```
const int maxN = 2e5+5
int a[maxN];
struct node{
  II val;
 II IzAdd
```

```
II IzSet:
  node(){}
} tree[ maxN<<2];
#define Ic p<<1
#define rc (p<<1)+1
inline void pushup(int p){
  tree[p].val = tree[lc].val + tree[rc].val;
void pushdown(int p, int I, int mid, int r){
  //lazy: range set
  if(tree[p].lzSet != 0){
    tree[Ic].IzSet = tree[rc].IzSet = tree[p].IzSet;
tree[Ic].val = (mid-I+1) * tree[p].IzSet;
tree[rc].val = (r-mid) * tree[p].IzSet;
    tree[Ic].IzAdd = tree[rc].IzAdd = O;
    tree[p].lzSet = 0;
  }else if(tree[p].lzAdd != 0){ //lazy: range add
    if(tree[Ic].IzSet == 0)tree[Ic].IzAdd += tree[p].IzAdd;
      tree[Ic].IzSet += tree[p].IzAdd;
      tree[Ic].IzAdd = 0;
    if(tree[rc].lzSet == 0)tree[rc].lzAdd += tree[p].lzAdd;
      tree[rc].lzSet += tree[p].lzAdd;
      tree[rc].IzAdd = 0;
    tree[lc].val += (mid-l+1) * tree[p].lzAdd;
    tree[rc].val += (r-mid) * tree[p].lzAdd;
    tree[p].IzAdd = 0;
  return;
}
void build(int p, int I, int r){
  tree[p].IzAdd = tree[p].IzSet = 0;
  if(l == r){
    tree[p].val = a[l];
    return;
  i nt mid = (I + r) >> 1;
  build(Ic, I, mid);
  bui I d(rc, mi d+1, r);
  pushup(p);
  return;
void add(int p, int I, int r, int a, int b, II val){
  if(a > r \mid | b < I)return;
  if(a <= | && r <= b){
tree[p].val += (r-|+1) * val;
    if(tree[p].lzSet == 0) tree[p].lzAdd += val;
    else tree[p].lzSet += val;
  int mid = (1+r) >> 1;
  pushdown(p, I, mid, r);
  add(Ic, I, mid, a, b, val);
  add(rc, mid+1, r, a, b, val);
  pushup(p);
void st(int p, int I, int r, int a, int b, II val){
  if(a > r || b < l)return;
  if(a \le 1 \& r \le b){
    tree[p].val = (r-l+1) * val;
    tree[p].IzAdd = 0;
    tree[p].lzSet = val;
  int mid = (1+r) >> 1;
  pushdown(p, I, mid, r);
  st(Ic, I, mid, a, b, val)
  st(rc, mid+1, r, a, b, val);
  pushup(p);
  return;
}
II query(int p, int I, int r, int a, int b){
  if(a > r \mid \mid b < I)return 0;
  if(a <= I && r <= b)return tree[p].val;</pre>
  int mid = (1+r) >> 1;
  pushdown(p, l, mid, r);
```

```
return query(lc, l, mid, a, b) + query(rc, mid+1, r, a, b);
3.4 Treap
typedef struct item * pitem
struct item {
  int prior, value, cnt;
bool rev;
  pitem I, r;
int cnt (pitemit) {
 return it ? it->cnt : 0;
}
void upd_cnt (pitemit) {
 if (it)
    it->cnt = cnt(it->l) + cnt(it->r) + 1;
void push (pitemit) {
 if (it && it->rev) {
    it->rev = false;
    swap (it->I, it->r);
    if (it->I) it->I->rev ^= true;
    if (it->r) it->r->rev ^= true;
 }
}
void merge (pitem & t, pitem I, pitem r) {
  push (I):
  push (r);
  if (!| || !r)
t = | ? | : r;
  else if (I->prior > r->prior)
    merge (1->r, 1->r, r), t = 1;
    merge (r->1, l, r->l), t=r;
  upd_cnt (t);
}
void split (pitem t, pitem & I, pitem & r, int key, int add =
  if (!t)
   return void( I = r = 0 );
  push (t);
  int cur_key = add + cnt(t->I);
  if (key <= cur_key)</pre>
    split (t->1, l, t->1, key, add), r = t;
    split (t->r, t->r, r, key, add + 1 + cnt(t->l)), l = t;
  upd_cnt (t);
}
void reverse (pitem t, int I, int r) {
  pi tem t1, t2, t3;
  split (t, t1, t2, l);
  split (t2, t2, t3, r-l+1);
  t2->rev ^= true;
  merge (t, t1, t2);
  merge (t, t, t3);
void output (pitemt) {
 if (!t) return;
  push (t);
  output (t->I);
printf ("%d", t->value);
  output (t->r);
3.5 LiChaoTree
const int maxc = 1e6 + 10;
namespace lichao {
  struct line {
    long long a, b;
line() : a(0), b(0) {}
    line(long long a, long long b) : a(a), b(b) {}
```

const int maxc = 1e6 + 10; namespace lichao { struct line { long long a, b; line() : a(0), b(0) {} line(long long a, long long b) : a(a), b(b) {} long long operator()(int x) const { return a * x + b; } }; line st[maxc * 4]; int sz; void init() { sz = 0; } void add(int l, int r, line tl, int o) {

```
bool | cp = st[o](l) > tl(l);
bool | mcp = st[o]((l + r) / 2) > tl((l + r) / 2);
if (mcp) | swap(st[o], tl);
if (r - l == 1) | return;
if (| cp ! = mcp) {
    add(l, (l + r) / 2, tl, 2 * o + 1);
}
else {
    add((l + r) / 2, r, tl, 2 * o + 2);;
}

long | long | query(int | l, int | r, int | x, int | o) {
    if (r - l == 1) | return | st[o](x);
    if (x < (l + r) / 2) {
        return | min(st[o](x), | query(l, (l + r) / 2, | x, | 2 * o + 1))
    ;
}
else {
    return | min(st[o](x), | query((l + r) / 2, | r, | x, | 2 * o + 2))
    ;
}
</pre>
```

4 Graph

4.1 Bi-Connected Component

```
| vector < i nt > dfn(n), low(n);
int timer = 0, bcc = 0;
vector <i nt > i d(n);
stack<int> stk:
function<void(int, int) > tarjan = [&](int u, int fa){
     dfn[u] = Iow[u] = timer++;
     stk.push(u);
     for (auto e: adj [u]){
         int v = e.first;
         int w = e. second:
         if (w == fa) continue:
         if (!dfn[v]){
             tarjan(v, w)
             low[u] = min(low[u], low[v]);
         } el se{
             low[u] = min(low[u], dfn[v]);
     if (low[u] == dfn[u]){
         while (true) {
             int v = stk.top();
       stk.pop();
             id[v] = bcc;
             if (v == u) break;
         bcc++;
     }
};
```

4.2 Strongly Connected Component

```
| vector < i nt > dfn(n), low(n), ins(n);
int timer = 0, scc = 0;
vector < i nt > i d(n);
stack<int> stk;
voi d tarj an(i nt u){
  low[u] = dfn[u] = ++timer;
  ins[u] = 1
  stk.push(u);
  for (int v: adj [u]){
    if (!dfn[v]){
       tarjan(v);
       low[u] = min(low[u], low[v]);
     }else if (ins[v]){
       low[u] = min(low[u], dfn[v]);
    }
   if (Iow[u] == dfn[u]) \{ \\
     int v;
     do {
       v = stk. top(); stk. pop();
       id[v] = scc;
       ins[v] = 0;
     } while(v != u);
  }
```

}

4.3 Lowest Common Ancestor

```
const int mxN = 1e5+5, LOG = 18;
vector <i nt > adj [ mxN]
int dep[mxN], up[mxN][LOG];
void dfs(int u, int p){
 for (int v:adj[u]){
    if (v == p)
     conti nue;
    dep[v] = dep[u] + 1;
    up[v][0] = u;
    for (int j = 1; j < LOG; j + +) {
      up[v][j] = up[up[v][j - 1]][j - 1];
    dfs(v, u);
 }
}
int lca(int u, int v){
  if(dep[u] < dep[v])</pre>
    swap(u, v);
  int k = dep[u] - dep[v];
  for (int j = LOG - 1; j >= 0; j --) {
    if(k & (1 << j)) {
      u = up[u][j];
    }
  if(u == v) {
    return u;
  for (int j = LOG - 1; j >= 0; j --) {
    if(up[u][j] != up[v][j]) {
      u = up[u][j];
      v = up[v][j];
    }
  }
  return up[u][0];
```

5 String

5.1 Knuth-Morris-Pratt Algorithm

```
vector<int> Failure(const string &s) {
  vector < i nt > f(s. size(), 0);
   // f[i] = length of the longest prefix (excluding s[0:i])
        such that it coincides with the suffix of s[0:i] of the
        same Length
  // i + 1 - f[i] is the length of the smallest recurring
       period of s[O:i]
  int k = 0:
  for (int i = 1; i < (int)s. size(); ++i) {
     while (k > 0 \&\& s[i] != s[k]) k = f[k - 1];
     if (s[i] == s[k]) ++k;
    f[i] = k;
  }
  return f;
vector<int> Search(const string &s, const string &t) {
  // return O-indexed occurrence of t in s
  vector < i nt > f = Failure(t), res;
  for (int i = 0, k = 0; i < (int)s. size(); ++i) {
     while (k > 0 \&\& (k == (int)t.size() || s[i] != t[k])) k = f
         [k - 1];
     if(s[i] == t[k]) ++k;
    if (k == (int)t.size()) res. push_back(i - t.size() + 1);
  return res:
| }
```

5.2 Z Algorithm

```
| int z[maxn];
// z[i] = LCP of suffix i and suffix 0
| void z_function(const string& s) {
| memset(z, O, sizeof(z));
| z[O] = (int)s.length();
| int l = O, r = O;
| for (int i = 1; i < s.length(); ++i) {
| z[i] = max(O, min(z[i - l], r - i + 1));
| while (i + z[i] < s.length() && s[z[i]] == s[i + z[i]]) {
| l = i; r = i + z[i];
| ++z[i];
| }
| }
| }</pre>
```

5.3 Suffix Automaton

```
struct SAM {
  static const int maxn = 5e5 + 5;
  int nxt[maxn][26], to[maxn], len[maxn];
  int root, last, sz;
  int gnode(int x) {
    for (int i = 0; i < 26; ++i) nxt[sz][i] = -1;
    to[sz] = -1:
    len[sz] = x;
    return sz++;
  voi d i ni t() {
    sz = 0;
    root = gnode(0);
    last = root;
  voi d push(int c) {
    int cur = last:
    last = gnode(len[last] + 1);
    for (; \simcur && nxt[cur][c] == -1; cur = to[cur]) nxt[cur][c
        1 = Last:
    if (cur == -1) return to[last] = root, void();
    int link = nxt[cur][c];
    if (len[link] == len[cur] + 1) return to[last] = link, void
    int tlink = gnode(len[cur] + 1);
    for (; ~cur && nxt[cur][c] == link; cur = to[cur]) nxt[cur
        ][c] = tlink;
    for (int i = 0; i < 26; ++i) nxt[tlink][i] = nxt[link][i];
    to[tlink] = to[link];
    to[link] = tlink;
    to[last] = tlink;
  voi d add(const string &s) {
    for (int i = 0; i < s. size(); ++i) push(s[i] - 'a');
  bool find(const string &s) {
    int cur = root;
    for (int i = 0; i < s. size(); ++i) {
      cur = nxt[cur][s[i] - 'a'];
      if (cur == -1) return false;
    return true;
  int solve(const string &t) {
    int res = 0, cnt = 0;
    int cur = root;
    for (int i = 0; i < t. size(); ++i) {
     if (~nxt[cur][t[i] - 'a']) {
        ++cnt:
        cur = nxt[cur][t[i] - 'a'];
      } el se {
        for (;
              ~cur && nxt[cur][t[i] - 'a'] == -1; cur = to[cur
             1):
        if (\sim cur) cnt = len[cur] + 1, cur = nxt[cur][t[i] - 'a']
        else cnt = 0, cur = root;
      res = max(res, cnt);
    return res;
```

5.4 Suffix Array

```
// sa[i]: sa[i]-th suffix is the i-th lexigraphically smallest
     suffix.
// lcp[i]: longest common prefix of suffix sa[i] and suffix sa[
     i - 11.
namespace sfx {
vector<int> Build(const string &s) {
  int n = s. size();
  vector < int> str(n * 2), sa(n * 2), c(max(n, 256) * 2), x(max(
       n, 256), p(n), q(n * 2), t(n * 2);
  for (int i = 0; i < n; ++i) str[i] = s[i];
  auto Pre = [&](int *sa, int *c, int n, int z) {
    memset(sa, 0, sizeof(int) * n);
    memcpy(x.data(), c, sizeof(int) * z);
  auto Induce = [&] (int *sa, int *c, int *s, int *t, int n, int
        z) {
    memcpy(x. data() + 1, c, sizeof(int) * (z - 1));
    for (int i = 0; i < n; ++i) if (sa[i] &&!t[sa[i] - 1]) sa[
         x[s[sa[i] - 1]] ++] = sa[i] - 1;
    memcpy(x.data(), c, sizeof(int) * z);
```

```
for (int i = n - 1; i >= 0; --i) if (sa[i] && t[sa[i] - 1])
           sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
  };
   auto SALS = [&] (auto self, int *s, int *sa, int *p, int *q,
       int *t, int *c, int n, int z) -> void {
     bool uni q = t[n - 1] = true;
     int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n, last =
          - 1:
     memset(c, 0, sizeof(int) * z);
     for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
     for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
     if (uni q) {
       for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
       return;
     for (int i = n - 2; i >= 0; --i) t[i] = (s[i] == s[i + 1] ?
          t[i + 1] : s[i] < s[i + 1]);
     Pre(sa, c, n, z);
     for (int i = 1; i <= n - 1; ++i) if (t[i] \&\&!t[i - 1]) sa
          [--x[s[i]]] = p[q[i] = nn++] = i;
     Induce(sa, c, s, t, n, z);
     for (int i = 0; i < n; ++i) if (sa[i] && t[sa[i]] && !t[sa[
          i] - 1]) {
       bool neq = last < 0 \mid \mid memcmp(s + sa[i], s + last, (p[q[
            sa[i]] + 1] - sa[i]) * sizeof(int));
       ns[q[last = sa[i]]] = nmxz += neq;
     self(self, ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz +
          1):
     Pre(sa, c, n, z);
     for (int i = nn - 1; i >= 0; --i) sa[--x[s[p[nsa[i]]]]] = p
          [nsa[i]];
     Induce(sa, c, s, t, n, z);
   SAI S(SAI S, str. data(), sa. data(), p. data(), q. data(), t. data
        (), c. data(), n + 1, 256);
   return vector <i nt > (sa. begi n() + 1, sa. begi n() + n + 1);
 vector <i nt > BuildLCP(const vector <i nt > &sa, const string &s) {
  int n = s. size()
   vector <i nt > lcp(n), rev(n);
  for (int i = 0; i < n; ++i) rev[sa[i]] = i; for (int i = 0, ptr = 0; i < n; ++i) {
    if (!rev[i]) {
       ptr = 0;
       continue;
     while (i + ptr < n \&\& s[i + ptr] == s[sa[rev[i] - 1] + ptr]
         ]) ptr++;
    lcp[rev[i]] = ptr ? ptr-- : 0;
  return lcp:
|}} // namespace sfx
```

6 Math

6.1 Compute Primes

```
void compPrimes(int N) {
   vector < int > primes, leastFac;
   for (int i = 0; i < N; i++) {
     leastFac.push_back(0);
   }
   leastFac[0] = 1; leastFac[1] = 1;
   for (int i = 2; i < N; i++) {
     if (leastFac[i] == 0) {
        primes.push_back(i);
        leastFac[i] = i;
   }
   for (int j = 0; j < primes.size() && i*primes[j] < N &&
        primes[j] <= leastFac[i]; j++) {
        leastFac[i*primes[j]] = primes[j];
   }
   }
}</pre>
```

6.2 Extended GCD

```
templ ate <typename T> tupl e<T, T, T> extgcd(T a, T b) {
   if (!b) return make_tupl e(a, 1, 0);
   T d, x, y;
   tie(d, x, y) = extgcd(b, a % b);
   return make_tupl e(d, y, x - (a / b) * y);
}
```

6.3 Primitive Root

```
int powmod (int a, int b, int p) {
  int res = 1;
   while (b)
    if (b & 1)
      res = int (res * 111 * a % p), --b;
     el se
       a = int (a * 111 * a % p), b >>= 1;
  return res;
int generator (int p) {
  vector <i nt > fact;
  int phi = p-1, n = phi;
  for (int i = 2; i * i <= n; ++ i)
     if (n \% i == 0) {
       fact.push_back (i);
       while (n \% i == 0)
        n /= i:
  if (n > 1)
     fact.push_back (n);
   for (int res=2; res<=p; ++res) {</pre>
     bool ok = true;
     for (size_t i =0; i <fact. size() && ok; ++i)</pre>
       ok &= powmod (res, phi / fact[i], p) ! = 1;
     if (ok) return res;
  return -1;
```

6.4 Discrete Logarithm

```
1// Returns minimum x for which a ^ x % m = b % m
int solve(int a, int b, int m) {
   a %= m, b %= m
   int k = 1, add = 0, g;
   while ((g = gcd(a, m)) > 1) {
     if (b == k)
        return add;
     if (b % g)
        return -1;
     b /= g, m /= g, ++add;
k = (k * 11 I * a / g) % m
   int n = sqrt(m) + 1;
   int an = 1;
   for (int i = 0; i < n; ++i)
     an = (an * 111 * a) % m
   unordered_map<int, int> vals;
   for (int q = 0, cur = b; q <= n; ++q) {
  vals[cur] = q;
  cur = (cur * 1|| * a) % m
   for (int p = 1, cur = k; p <= n; ++p) { cur = (cur * 111 * an) % m
     if (vals.count(cur)) {
        int ans = n * p - vals[cur] + add;
        return ans:
     }
   return -1;
```

7 Geometry

7.1 Basic

```
bool same(double a, double b) { return abs(a - b) < eps; }

struct P {
    double x, y;
    P() : x(0), y(0) {}
    P(double x, double y) : x(x), y(y) {}
    P operator + (P b) { return P(x + b.x, y + b.y); }
    P operator - (P b) { return P(x - b.x, y - b.y); }
    P operator * (double b) { return P(x * b, y * b); }
    P operator * (double b) { return P(x / b, y / b); }
    double operator * (P b) { return x * b.x + y * b.y; }
    double operator ^ (P b) { return x * b.y - y * b.x; }
    double abs() { return hypot(x, y); }
    P unit() { return *this / abs(); }
    P rot(double o) {
        double c = cos(o), s = sin(o);
```

```
return P(c * x - s * y, s * x + c * y);
     double angle() { return atan2(y, x); }
};
 struct L {
      // ax + by + c = 0
      double a, b, c, o;
      P pa, pb;
     L() : a(0), b(0), c(0), o(0), pa(), pb() {}
     L(P pa, P pb) : a(pa. y - pb. y), b(pb. x - pa. x), c(pa ^ pb), o
                    (atan2(-a, b)), pa(pa), pb(pb) {}
      P project(P p) { return pa + (pb - pa). unit() * ((pb - pa)
                    (p - pa) / (pb - pa).abs()); }
      P \ reflect(P \ p) \ \{ \ return \ p + (project(p) - p) * 2; \}
      double get_ratio(P p) { return (p - pa) * (pb - pa) / ((pb -
                    pa).abs() * (pb - pa).abs()); }
}:
 bool SegmentIntersect(P p1, P p2, P p3, P p4) {
     if (max(p1. x, p2. x) < min(p3. x, p4. x) | | max(p3. x, p4. x) <
                    min(p1. x, p2. x)) return false;
     if (max(p1. y, p2. y) < min(p3. y, p4. y) | | max(p3. y, p4. y) <
                   min(p1. y, p2. y)) return false;
      return si gn((p3 - p1) ^{\circ} (p4 - p1)) ^{*} si gn((p3 - p2) ^{\circ} (p4 -
                    p2)) <= 0 &&
            sign((p1 - p3) ^ (p2 - p3)) * sign((p1 - p4) ^ (p2 - p4))
                           <= 0:
}
 bool parallel(L x, L y) { return same(x.a * y.b, x.b * y.a); }
PIntersect(L x, L y) \{ return P(-x.b * y.c + x.c * y.b, x.a * y.
             y.c - x.c * y.a) / (-x.a * y.b + x.b * y.a); }
```

7.2 Half Plane Intersection

```
bool jizz(L | 1, L | 2, L | 3) {
  P p=Intersect(I2,I3);
  return ((I 1. pb-I 1. pa) ^(p-I 1. pa)) <- eps;</pre>
bool cmp(const L &a, const L &b){
 return same(a. o, b. o)?(((b. pb-b. pa) ^(a. pb-b. pa)) >eps): a. o<b. o;
// availble area for L I is (I.pb-I.pa)^(p-I.pa)>0
vector <P> HPI (vector <L> &ls){
  sort(Is. begin(), Is. end(), cmp);
  vector <L> pls(1, ls[0]);
  for(int i = 0; i < (int) | s. size(); ++i) if(! same(|s[i].o, pls. back().</pre>
       o))pls.push_back(ls[i]);
  deque<i nt > dq; dq. push_back(0); dq. push_back(1);
#define meow(a, b, c) while(dq. size() > 1u && jizz(pls[a], pls[b],
     pls[c]))
  for(int i = 2; i < (int) pls. size(); ++i){
    meow(i, dq. back(), dq[dq. si ze()-2])dq. pop_back();
    meow(i, dq[0], dq[1]) dq. pop_front();
    dq. push_back(i);
  meow(dq. front(), dq. back(), dq[dq. si ze()-2])dq. pop_back();
  meow(dq. back(), dq[0], dq[1]) dq. pop_front();
  if(dq. size() < 3u) return vector < P > (); // no solution or
       solution is not a convex
  vector <P> rt:
  for(int i =0; i <(int) dq. size(); ++i)rt. push_back(Intersect(pls[</pre>
       dq[i]], pl s[dq[(i+1)%dq. size()]]));
```

7.3 Slope & Fraction

```
| struct P {
    int x, y, i;
    P() : x(0), y(0), i(-1) {}

|};
|struct Frac {
    int u, d;
    void norm() {
        if (d == 0) {
            u = u > 0 ? 1 : u < 0 ? -1 : 0;
            return;

    }
    int g = __gcd(u, d);
    u /= g;
    d /= g;
    if (d < 0) {
        d *= -1;
}</pre>
```

```
u^* = -1:
      }
  }
};
bool operator > (const Frac &a, const Frac &b) {
  return 1II * a.u * b.d > 1II * b.u * a.d;
bool operator >= (const Frac &a, const Frac &b) {
  return 1|| * a. u * b. d >= 1|| * b. u * a. d;
bool operator < (const Frac &a, const Frac &b) {
  return 1|| * a. u * b. d < 1|| * b. u * a. d;</pre>
bool operator <= (const Frac &a, const Frac &b) {
  return 1|| * a. u * b. d <= 1|| * b. u * a. d;</pre>
ostream& operator << (ostream &o, const Frac &f) {
   o << f. u << "/" << f. d;
   return o:
Frac Slope(P &a, P &b) {
   Frac f;
   f. u = b. y - a. y;
   f. d = b. x - a. x;
   f. nor m();
   return f:
```

7.4 Convex order

```
int quard(P p) {
    if (p. x > 0 \&\& p. y >= 0) return 1;
    if (p. x \le 0 \&\& p. y > 0) return 2;
    if (p. x < 0 && p. y <= 0) return 3;
    if (p. x >= 0 \&\& p. y < 0) return 4;
    return -1:
P getcenter(vector<P>& p) {
    P res(0, 0); double n = (double) p. size();
     for (Pit : p) res. x += it. x, res. y += it. y;
    res. x /= n; res. y /= n;
    return res;
voi d convex_order(vector <P>& p) {
    P center = getcenter(p);
    auto cmp = [&](Pa, Pb) {
P tmpa = a - center, tmpb = b - center;
         int qa = quard(tmpa), qb = quard(tmpb);
         if (qa != qb) return qa < qb;</pre>
         return (tmpa ^ tmpb) > 0;
    sort(ALL(p), cmp);
```

7.5 Area

```
double area(const vector < P>& fig) {
    double res = 0;
    for (unsigned i = 0; i < fig. size(); i++) {
        P p = i ? fig[i - 1] : fig. back();
        P q = fig[i];
        res += (p.x - q.x) * (p.y + q.y);
    }
    return fabs(res) / 2;
}</pre>
```