



Notebook - Maratona de Programação

Cabo HDMI, VGA, USB

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1 Graph

1.1 Bellman Ford

Complexity: $O(n * m)$ | $n = |\text{nodes}|$, $m = |\text{edges}|$
Finds shortest paths from a starting node to all nodes of the graph. The node can have negative cycle and belman-ford will detected

```
1 // TITLE: Bellman Ford
2 // COMPLEXITY:  $O(n * m)$  |  $n = |\text{nodes}|$ ,  $m = |\text{edges}|$ 
3 // DESCRIPTION: Finds shortest paths from a starting
  node to all nodes of the graph. The node can have
  negative cycle and belman-ford will detected
4
5 // a and b vertices, c cost
6 // [{a, b, c}], {a, b, c}]
7 vector<tuple<int, int, int>> edges;
8 int N;
9
10 void bellman_ford(int x){
11     for (int i = 0; i < N; i++){
12         dist[i] = oo;
13     }
14     dist[x] = 0;
15
16     for (int i = 0; i < N - 1; i++){
17         for (auto [a, b, c]: edges){
18             if (dist[a] == oo) continue;
19             dist[b] = min(dist[b], dist[a] + w);
20         }
21     }
22 }
23 // return true if has cycle
24 bool check_negative_cycle(int x){
25     for (int i = 0; i < N; i++){
26         dist[i] = oo;
27     }
28     dist[x] = 0;
29
30     for (int i = 0; i < N - 1; i++){
31         for (auto [a, b, c]: edges){
32             if (dist[a] == oo) continue;
33             dist[b] = min(dist[b], dist[a] + w);
34         }
35     }
36
37     for (auto [a, b, c]: edges){
38         if (dist[a] == oo) continue;
39         if (dist[a] + w < dist[b]){
40             return true;
41         }
42     }
43     return false;
44 }
45 ""
```

1.2 Topological Sort

Complexity: $O(N + M)$, N : Vertices, M : Arestas
Retorna no do grapho em ordem topologica, se a quantidade de nos retornada nao for igual a quantidade de nos e impossivel

```
1 // TITLE: Topological Sort
2 // COMPLEXITY:  $O(N + M)$ ,  $N$ : Vertices,  $M$ : Arestas
3 // DESCRIPTION: Retorna no do grapho em ordem
  topologica, se a quantidade de nos retornada nao
  for igual a quantidade de nos e impossivel
4
5 typedef vector<vector<int>> Adj_List;
```

```
6 typedef vector<int> Indegree_List; // How many nodes
  depend on him
7 typedef vector<int> Order_List; // The order in
  which the nodes appears
8
9 Order_List kahn(Adj_List adj, Indegree_List indegree)
10 {
11     queue<int> q;
12     // priority_queue<int> q; // If you want in
  lexicografic order
13     for (int i = 0; i < indegree.size(); i++) {
14         if (indegree[i] == 0)
15             q.push(i);
16     }
17     vector<int> order;
18
19     while (not q.empty()) {
20         auto a = q.front();
21         q.pop();
22
23         order.push_back(a);
24         for (auto b: adj[a]) {
25             indegree[b]--;
26             if (indegree[b] == 0)
27                 q.push(b);
28         }
29     }
30     return order;
31 }
32
33 int32_t main()
34 {
35
36     Order_List = kahn(adj, indegree);
37     if (Order_List.size() != N) {
38         cout << "IMPOSSIBLE" << endl;
39     }
40     return 0;
41 }
```

2 String

2.1 Z function

Complexity: Z function complexity
z function

```
1 // TITLE: Z function
2 // COMPLEXITY: Z function complexity
3 // DESCRIPTION: z function
4
5 void z_function(string& s)
6 {
7     return;
8 }
```

3 Set

3.1 Ordered Set

Complexity: $O(\log(n))$

```
1 // TITLE: Ordered Set
2 // COMPLEXITY:  $O(\log(n))$ 
3 // DESCRIPTION: Set but you can look witch elements is
  in position (k)
4
```

```

5 #include <ext/pb_ds/assoc_container.hpp>
6 #include <ext/pb_ds/tree_policy.hpp>
7 using namespace __gnu_pbds;
8
9 #define ordered_set tree<int, null_type, less<int>,
  rb_tree_tag, tree_order_statistics_node_update>
10
11 int32_t main() {
12     ordered_set o_set;
13
14     o_set.insert(5);
15     o_set.insert(1);
16     o_set.insert(2);
17     // o_set = {1, 2, 5}
18     5 == *(o_set.find_by_order(2));
19     2 == o_set.order_of_key(4); // {1, 2}
20 }

```

3.2 Set

Complexity: Insertion $\log(n)$

Keeps elements sorted, remove duplicates, upper_bound, lower_bound, find, count

```

1 // TITLE: Set
2 // COMPLEXITY: Insertion Log(n)
3 // Description: Keeps elements sorted, remove
  duplicates, upper_bound, lower_bound, find, count
4
5 int main() {
6     set<int> set1;

```

```

7
8     set1.insert(1);           //  $O(\log(n))$ 
9     set1.erase(1);           //  $O(\log(n))$ 
10
11     set1.upper_bound(1);      //  $O(\log(n))$ 
12     set1.lower_bound(1);      //  $O(\log(n))$ 
13     set1.find(1);             //  $O(\log(n))$ 
14     set1.count(1);            //  $O(\log(n))$ 
15
16     set1.size();              //  $O(1)$ 
17     set1.empty();             //  $O(1)$ 
18
19     set1.clear()              //  $O(1)$ 
20     return 0;
21 }

```

3.3 Multiset

Complexity: $O(\log(n))$

Same as set but you can have multiple elements with same values

```

1 // TITLE: Multiset
2 // COMPLEXITY:  $O(\log(n))$ 
3 // DESCRIPTION: Same as set but you can have multiple
  elements with same values
4
5 int main() {
6     multiset<int> set1;
7 }

```