Regular languages and nfa's

Theorem : [Linz, Theorem 3.1]

Every **regular language**, L, is **accepted** by some **nfa**.

I.e., For every **Σ-regular expression**, r, there is a **Σ-nfa** M, s.t.

$$L(r) = L(M).$$

Proof. By **structural induction** on $r \in RegExp(\Sigma)$,

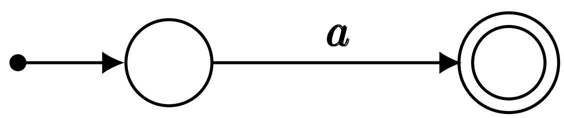
i.e. according to the **inductive definition** of $RegExp(\Sigma)$ (p. 3-1). □

(***Note***: By note (2) on p. 3-6, we can assume that every **nfa** we construct in the course of the proof has exactly one **final state**.)

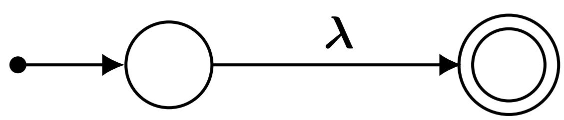
We proceed with the **inductive proof**:

Base cases:

- $r = a$ $(\in \Sigma)$. Then $L(r) = \{a\}$ is accepted by:



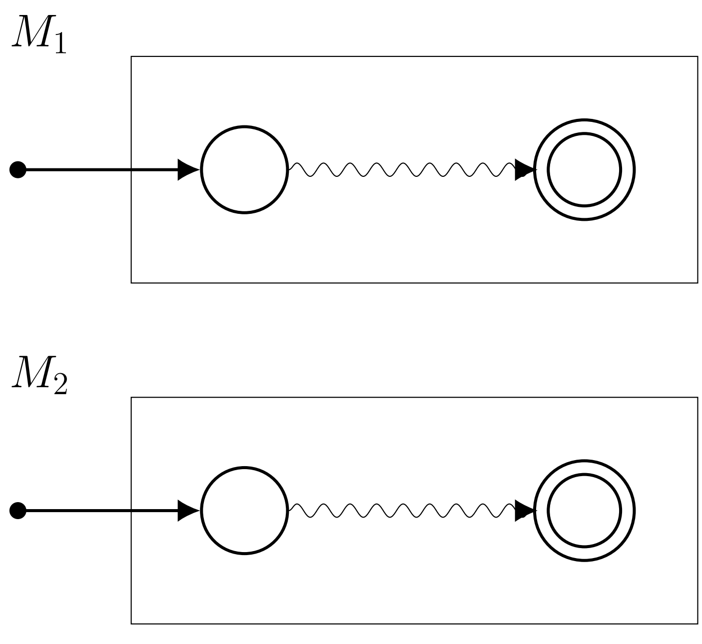
- $r = \lambda$. Then $L(r) = \{\lambda\}$ is accepted by:



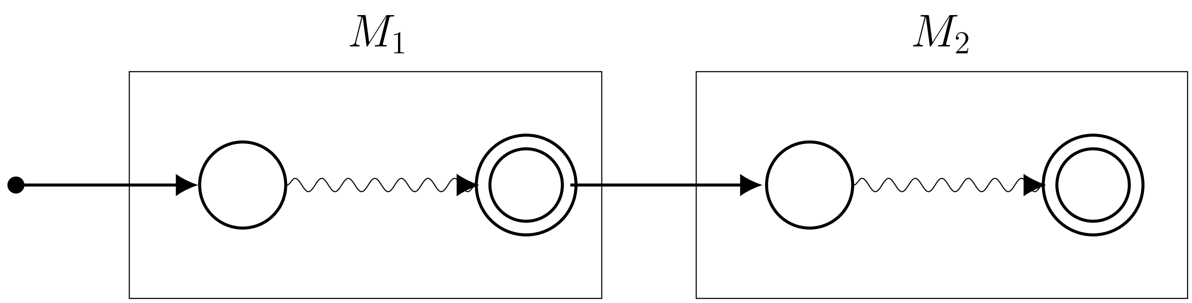
- $r = \emptyset$. Then $L(r) = \emptyset$ is accepted by:

Recursive steps:

- $r = r_1 + r_2$. Suppose (**induction hypothesis**) we have **nfa's** M_1 and M_2 which accept $L(r_1)$, $L(r_2)$ respectively. Then $L(r) = L(r_1) \cup L(r_2)$ is accepted by:

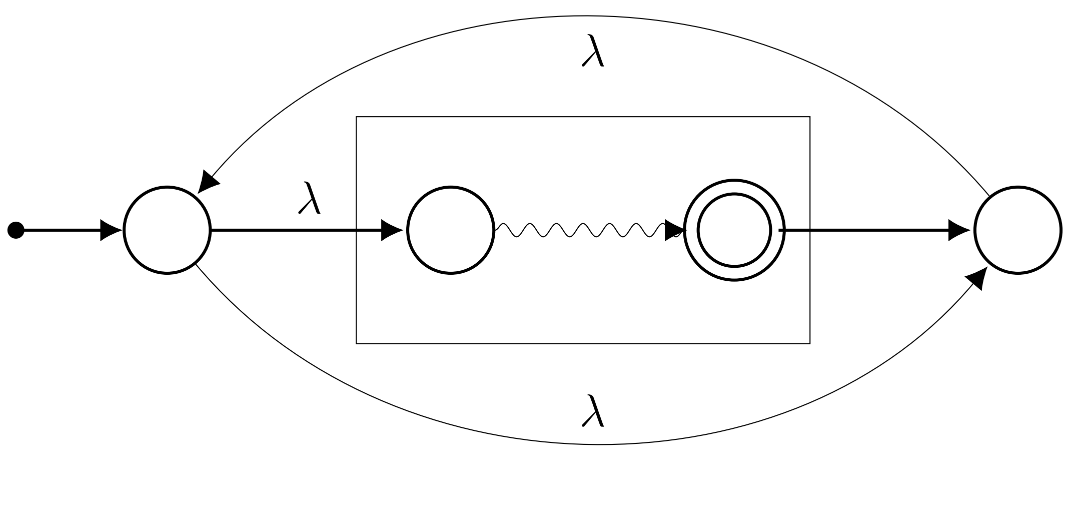


- $r = r_1 \cdot r_2$. Suppose (**induction hypothesis**) we have **nfa's** M_1 and M_2 which accept $L(r_1)$, $L(r_2)$ respectively. Then $L(r) = L(r_1) \cdot L(r_2)$ is accepted by:



- $r = r^*$. Suppose (**induction hypothesis**) we have an **nfa** M_1 which accepts $L(r_1)$. Then, $L(r) = L(r_1)^*$ is accepted by:

M



Q.   Is the following true?

$\forall r \in \boldsymbol{RegExp}(\Sigma),\ \exists\ \Sigma$-**dfa** which accepts $\boldsymbol{L(r)}$

Q.   Is the **converse** true?

$\forall\ \boldsymbol{\Sigma}$-**nfa** $\boldsymbol{M},\ \exists\, r \in \boldsymbol{RegExp(\Sigma)} :\ L(r) = L(M)$

Theorem :   [Linz, Theorem 3.2]

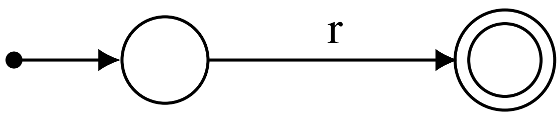
For every nfa, M, $L(M)$ is **regular**.
I.e. $\forall\ \boldsymbol{\Sigma}$-**nfa** $\boldsymbol{M},\ \exists\, r \in \boldsymbol{RegExp(\Sigma)} :\ L(r) = L(M)$

Proof :   (**outline**)

The proof uses the concept of **generalized transition graphs**, i.e. graphs where the edges are labeled not just by **elements** of Σ, but by Σ-**regular expressions**.

The graph for M is **transformed** into an equivalent graph with **fewer nodes**, and edges with **more complex** Σ-expressions.

This is repeated until we arrive at a graph of the form:



Then $\boldsymbol{L(M) = L(r)}$, i.e. $\boldsymbol{L(M)}$ is **regular**.

Combining Theorems 3.1 and 3.2, we have proven:

Theorem : For any Σ-language $\boldsymbol{L \subseteq \Sigma^*}$,

L is **regular** $\iff$ L is accepted by some **nfa**.

Note:

We have already seen:

For every **$\boldsymbol{\Sigma}$-nfa**, N, we can find an **equivalent** Σ-dfa, M.

i.e. with $L(M) = L(N)$

So we can **rewrite** the above theorem:

Theorem : For any Σ-language $L \subseteq \Sigma^*$,

L is **regular** $\iff$ L is accepted by some **dfa**

Note: Linz **defines** a language to be "regular" if it is accepted by some **dfa** [Linz, Def. 2.3]. We have not done that here – but it comes to the same thing by the theorems above.

Example: "**Reversing**" an **nfa**

Show if L is **regular**, so is L^{R}.

Use theorem on p. 3-11

Suppose L is regular.

Then, L, is accepted by some **nfa**:

$M = (Q, \Sigma, \delta, q_0, F)$

Define $M^{\mathsf{R}} = (Q, \Sigma, \delta^{\mathsf{R}}, q_0^{\mathsf{R}}, F^{\mathsf{R}})$ as follows:

1) **Reverse all arrows**

i.e. $q_1 \in \delta^{\mathsf{R}}(q_2, a) \Longleftrightarrow q_2 \in \delta^{\mathsf{R}}(q_1, a)$

2) **Interchange starting and final states**

i.e. $F^{\mathsf{R}} = \{q_0\}$

Let $q_0^{\mathsf{R}} =$ **new start state**, joined to **all old final states** by λ- transitions. (see p. 2-17)

Then, $L(M^{\mathsf{R}}) = (L(M))^{\mathsf{R}}$.