

4 Closure Properties of Regular Languages

[Linz § 4.1]

We show **closure** under some **simple set operations**

(1) L_1, L_2 **regular** $\implies L_1 \cup L_2$ **regular**

Proof : Suppose for some nfa's, M_1, M_2 :

$$L_1 = L(M_1), L_2 = L(M_2)$$

Let $M_1 \sqcup M_2 =$ the “**union**” of M_1 and M_2
shown on p. 3-8 (bottom).

Then $L_1 \cup L_2 = L(M_1 \sqcup M_2)$ is regular.

(2) L_1, L_2 **regular** $\implies L_1 L_2$ **regular**

Proof : Suppose $L_1 = L(M_1), L_2 = L(M_2)$ for nfa's M_1, M_2

Let $M_1 \cdot M_2 =$ the “**concatenation**” of M_1 and M_2
Shown on p. 3-9 (top).

Then $L_1 L_2 = L(M_1 \cdot M_2)$ is regular.

(3) L regular $\implies L^*$ regular

Proof : Suppose $L = L(M)$ for some nfa M .

Let M^* = the “**star**” of M
shown on p. 3-9 (bottom).

Then $L^* = L(M^*)$ is regular.

Given $L \subseteq \Sigma^*$,

$\bar{L} = \Sigma^* \setminus L$ = the **complement** of L in Σ .

Then:

(4) L regular $\implies \bar{L}$ regular

To prove this, we must first define the “**complement**” of a Σ -nfa.

How?

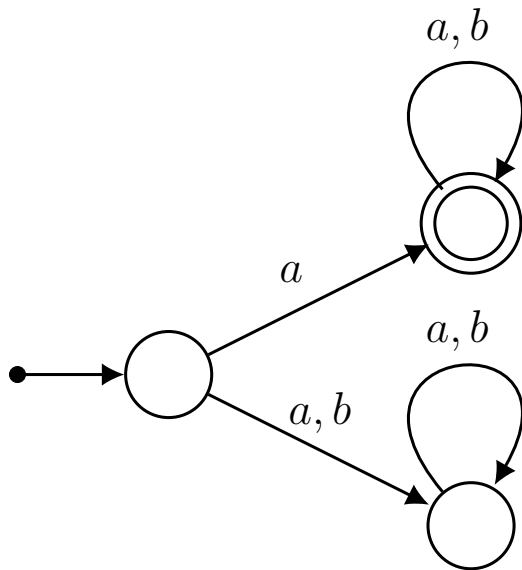
Proposal: Given $M = (Q, \Sigma, \delta, q_0, F)$

Define $\bar{M} = (Q, \Sigma, \delta, q_0, \bar{F})$

where $\bar{F} = Q \setminus F$

Problem Consider e.g. $M : \Sigma = \{a, b\}$

\overline{M} :



$$L(M) = \{au \mid u \in \Sigma^*\}$$

$$L(\overline{M}) = \Sigma^*$$

Solution Replace M by **equivalent dfa** M^D (p. 2-13)

Then if
$$L = L(M),$$

then
$$\overline{L} = L(\overline{M^D})$$

From now on we write, for any **nfa** M ,

$$\overline{M} \text{ means } \overline{M^D}$$

so that if $L = L(M)$

then $\overline{L} = L(\overline{M})$

is regular.

(5) L_1, L_2 **regular** $\implies L_1 \cap L_2$ **regular**

Why?

Suppose $L_1 = L(M_1)$, $L_2 = L(M_2)$ (**nfa**'s).

Define the nfa $M_1 \sqcap M_2 = \overline{\overline{M_1} \sqcup \overline{M_2}}$

Then

$$\begin{aligned} L_1 \cap L_2 &= \overline{\overline{L_1} \cup \overline{L_2}} \\ &= \underline{L(M_1 \sqcap M_2)} \end{aligned}$$