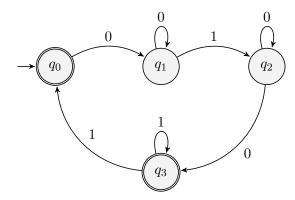
Computer Science III 2024-III WorkShop No. 1 — The Old Times

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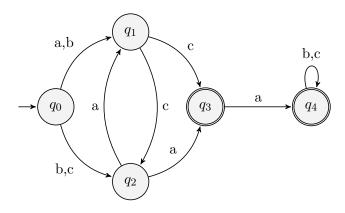
- 1. For each of the following languages, define the corresponding finite-state machine:
 - (i) $\Sigma = \{0, 1, 2\}$. $L = (01^*2 \cup 2102)^*101(01 \cup 12 \cup 20)^*$.
 - (ii) $\Sigma = \{a, b, c\}$. $L = (abc \cup bca \cup cab)(abc \cup bca \cup cab)^*$.
 - (iii) $\Sigma = \{a, b, c\}$. $L = (abc \cup bca \cup cab)^*(abc \cup bca \cup cab)$.
 - (iv) $\Sigma = \{0, 1, 2\}$. $L = (01^2 \cup 10^2 \cup 21^0)^*(01 \cup 12 \cup 20)^*101$.
- 2. For each one of the following finite-state machines, define the corresponding regular expression and a generative grammar:
 - (i) $\Sigma = \{0, 1\}.$



(ii) $\Sigma = \{a, b, c\}.$

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Any comment or concern related to this document could be send to Carlos A. Sierra at e-mail: cavir-guezs@udistrital.edu.co



- 3. For each of the following regular expressions, define the corresponding generative grammar (all over the alphabet $\Sigma = \{a, b, c, d\}$):
 - (i) $\{a^i b^j c^j d^i : i, j \ge 1\}.$
 - (ii) $\{a^i b^i c^j d^j : i, j \ge 1\}.$
 - (iii) $\{a^ib^jc^jd^i: i, j \ge 1\} \cup \{a^ib^ic^jd^j: i, j \ge 1\}.$
 - (iv) $\{a^ib^jc^{i+j}: i \ge 0, j \ge 1\}.$
- 4. Be G a context-free grammar with the following productions:

$$G = \left\{ \begin{array}{l} S \rightarrow ABC \mid BaC \mid aB \\ A \rightarrow Aa \mid a \\ B \rightarrow BAB \mid bab \\ C \rightarrow cC \mid \lambda \end{array} \right.$$

Found derivation trees for the following strings:

- (i) $w_1 = abab$.
- (ii) $w_2 = babacc$.
- (iii) $w_3 = ababababc$.
- 5. As follows there is a context-free grammar to generate real numbers without sign, the alphabet is $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, ., +, -, E\}$:

$$\begin{array}{lll} < real > & \rightarrow & < digits > < decimal > < exp > \\ < digits > & \rightarrow & < digits > < digits > | \ 0 \ | \ 1 \ | \ 2 \ | \ 3 \ | \ 4 \ | \ 5 \ | \ 6 \ | \ 7 \ | \ 8 \ | \ 9 \\ < decimal > & \rightarrow & < digits > | \ \lambda \\ < exp > & \rightarrow & E < digits > | \ E + < digits > | \ E - < digits > | \ \lambda \\ \end{array}$$

Define the derivation tree for the following strings:

- (i) $w_1 = 47.236$
- (ii) $w_2 = 321.25E + 35$

- (iii) $w_3 = 0.8E9$
- (iv) $w_4 = 0.8E + 9$
- 6. As follows there is a context-free grammar to generate identifiers, identifiers are strings of letters and digits, starting with a letter:

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 \begin{array}{lll} < identifier > & \rightarrow & < letter > < lsds > \\ < lsds > & \rightarrow & < letter > < lsds > | < digit > < lsds > | \lambda \\ < letter > & \rightarrow & a \mid b \mid c \mid \dots \mid x \mid y \mid z \mid A \mid B \mid C \mid \dots \mid X \mid Y \mid Z \\ < digit > & \rightarrow & 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \end{array}
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Define the derivation tree for the following names:

- (i) $w_1 = MyVariable$
- (ii) $w_2 = temp2$
- (iii) $w_3 = string2int$
- (iv) $w_4 = 2NotAVariable$

Deadline: Saturday, 5th of October, 2024, 18:00 (local time).