DataBases Design & Relational Algebra

DataBase Foundations

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Basic Concepts

Normalization

Relational Algebra





Basic Concepts

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Relational Algebra





DataBases Design Foundations

- In the context of **databases**, the **design** of a database is the process of producing a **detailed** data model of a database.
- This data model contains all the needed logical and physical design choices and physical storage parameters needed to generate a design in a data definition language, which can then be used to create a database.
- A fully attributed data model contains detailed attributes for each entity.
- Relational Data Models avoid redundancy and inconsistency by ensuring that data is normalized.





Set Theory in Databases

- The **set theory** is a branch of **hmathematical logic** that studies sets, which are collections of objects.
- The set theory is applied in databases to define the relational model and the relational algebra.
- The relational model is a mathematical model of data for large shared data banks and it has a solid theoretical foundation.
- The **relational algebra** is a procedural query language, which takes relations as input and produces relations as output.





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Normalization in Databases

- Normalization is the process of organizing the columns (attributes) and tables (relations) of a relational database to minimize data redundancy.
- Normalization involves decomposing a table into smaller tables and defining relationships between them.
- The objective is to isolate data so that additions, deletions, and modifications of a field can be made in just one table and then propagated through the rest of the database using the defined relationships.





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Onthologies

- An ontology is a formal naming and definition of the types, properties, and interrelationships of the entities that really or fundamentally exist for a particular domain of discourse.
- Ontologies are used in databases to define the schema of the database.
- The schema of a database is a formal definition of the structure of the database: the types of data that are stored, the relationships between the data, and the constraints on the data.





Normal Forms

- First normal form (1NF): The table is a two-dimensional table with rows and columns. Each column contains atomic values, and there are no repeating groups or arrays.
- Second normal form (2NF): The table is in first normal form and all the non-key attributes are fully functionally dependent on the primary key.
- Third normal form (3NF): The table is in second normal form and all the non-key attributes are non-transitively dependent on the primary key.
- Fourth normal form (4NF): The table is in third normal form and there are no multi-valued dependencies.





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What is relational algebra?

- The relational algebra is a procedural query language, which takes relations as input and produces relations as output.
- The relational algebra is a set of operations that can be performed on a relation. Also, it is used to define the relational model, which is a mathematical model of data for large shared data banks.
- Let's take a look at the basic operations of the **relational algebra**. First, remember next table called **Students**:

ID	Name	Lastname	Address	Phone	Age
1	John	Doe	123 Fake St	555-1234	25
2	Jane	Smith	456 Elm St	555-5678	30
3	Mike	Johnson	789 Evergreen St	555-9012	35





Select Operation

Definition

Select: $\sigma_{\text{condition}}(R)$, is a unary operation that returns the rows (subset) of R that satisfy the condition.

For example, the following expression selects the students whose age is greater than 25:

 $\sigma_{\mathsf{Age}>25}(\mathsf{Students})$





Project Operation

Definition

Project: $\pi_{\text{column_list}}(R)$, is a unary operation that returns the columns (subset) of R that are specified in the column list.

For example, the following expression projects the name and lastname of the students:

 $\pi_{\mathsf{Name, Lastname}}(\mathsf{Students})$





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Union Operation

Definition

Union: $R \cup S$, is a binary operation that returns the rows that are in R or in S.

For example, the following expression returns the students whose age is greater than 25 or whose lastname is Johnson:

$$\sigma_{\mathsf{Age}>25}(\mathsf{Students}) \cup \sigma_{\mathsf{Lastname}=\mathsf{Johnson}}(\mathsf{Students})$$





Set Different Operation

Definition

Set Different: R-S, is a binary operation that returns the rows that are in R but not in S.

For example, the following expression returns the students whose age is greater than 25 but not whose lastname is Johnson:

$$\sigma_{\mathsf{Age}>25}(\mathsf{Students}) - \sigma_{\mathsf{Lastname} = \mathsf{Johnson}}(\mathsf{Students})$$





Cartesian Product Operation

Definition

Cartesian Product: $R \times S$, is a binary operation that returns the Cartesian product of R and S. A formal definition is:

$$R \times S = \{r \cup s \mid r \in R \land s \in S\}$$

For example, the following expression returns the Cartesian product of the students and the courses:

 $Students \times Courses$





Rename Operation

Definition

Rename: $\rho_{\text{new_name}}(R)$, is a unary operation that returns the relation Rwith the name R changed to new_name .

For example, the following expression returns the students relation with the name changed to People:

 $\rho_{\mathsf{People}}(\mathsf{Students})$





Exercises

- Select the students whose age is greater than 25 and whose lastname is Johnson.
- Project the name and lastname of the students whose age is greater than 25.
- Select the students whose age is greater than 25 and whose lastname is Johnson, and project the name and lastname of the students, and rename the relation to People.





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Normalization





Thanks!

Questions?



Repo: https://github.com/EngAndres/ud-public/tree/main/courses/databases-foundations



