## Elliptic curves and their moduli spaces Exercise sheet 2

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## Problem 1.

a) From algebraic geometry 1 we know that for every scheme /k T, we have an isomorphism

$$\operatorname{Hom}_{Sch/k}(T, \mathbb{A}^n_k) \to \operatorname{Hom}_k(k[x_1, \dots, x_n], \mathcal{O}_T(T)).$$

which is natural in T. After picking the basis for A, we have a simple identification

$$\operatorname{Hom}_{k}(k[x_{1},\ldots,x_{n}],\mathcal{O}_{T}(T)) \to \mathcal{O}_{T}(T) \otimes_{k} A$$

$$\varphi \mapsto \sum_{i} \varphi(x_{i}) \otimes x_{i},$$

where we use  $x_i$  to denote basis as well. This identification is simply on the level of sets, because although  $\mathcal{O}_T(T) \otimes_k A$  has the structure of a (maybe non-commutative) k-algebra,  $\operatorname{Hom}_k(k[x_1,\ldots,x_n],\mathcal{O}_T(T))$  doesn't seem to have any structure (is there any? Set of maps of k-algebras is nothing more than a set, right? We're assuming it has to map 1 to 1 and be k-linear).

For second part, we want some sort of criterion for when is  $\sum_i f_i \otimes x_i$  invertible. We can represent  $\mathcal{O}_T(T) \otimes_k A$  by embedding it into algebra of endomorphisms  $\operatorname{End}(\mathcal{O}_T(T)^n)$  which is isomorphic to matrix algebra  $M_{n \times n}(\mathcal{O}_T(T))$ . Then we can simply calculate if the element is invertible by evaluating its determinant. In our case determinant is a polynomial in  $n = \dim A$  variables (matrix depends on n coefficients, for example first column, all others are implicitly given by that). So for the ring we take  $k[x_1, \ldots, x_n, \det(x_1, \ldots, x_n)^{-1}]$ .

Maybe it would be easier to define ring from scratch and not look at open subschemes of  $\mathbb{A}^n_k$ . In which case we take  $k[t_{ij}, i, j \in \{1, \dots, n\}]$ , invert  $\det(t_{ij})$  and then quotient by relations  $t_{ij} = \sum_k a^k_{ij} t_{k1}$  where  $a^k_{ij}$  represents k-th coordinate in expansion of the product  $(t_{11}x_1 + \dots + t_{n1}x_n)x_j$  as  $i \in \{1, \dots, n\}$  and  $j \in \{2, \dots, n\}$ . In example of  $\mathbb{R}$ -algebra  $\mathbb{C}$  that would mean

$$\mathbb{R}[x_{11}, x_{21}, x_{12}, x_{22}, (x_{11}x_{22} - x_{12}x_{21})^{-1}]/(x_{12} = -x_{21}, x_{22} = x_{11})$$

b)

c) Let  $k = \mathbb{R}$ ,  $A = \mathbb{C}$  and  $G = \underline{A}^{\times}$ . In part a) of problem 1 we argued that  $\underline{A}^{\times}$  is the affine open subscheme of  $\mathbb{A}^n_k$  we get by inverting the determinant of embedding  $A \hookrightarrow \operatorname{End}(k^n) \cong M_{n \times n}(k)$ . In this case (one possible) embedding is  $a + bi \mapsto \begin{bmatrix} a & -b \\ b & a \end{bmatrix}$ . So we have to invert  $a^2 + b^2$ . So now we need to define a scheme morphism. Since both G and  $\mathbb{G}_{m,\mathbb{R}}$  are affine, we can give map of rings:

$$\begin{split} \mathbb{R}[t,t^{-1}] &\to \mathbb{R}[x,y,(x^2+y^2)^{-1}] \\ t &\mapsto x^2+y^2 \end{split}$$

So if we have a map  $\operatorname{Spec}(\mathbb{R}) \to G$  (i.e. a  $\mathbb{R}$ -valued point z) and post-compose it with  $G \to \mathbb{G}_{m,\mathbb{R}}$ , we get  $\operatorname{Spec}(\mathbb{R}) \to \mathbb{G}_{m,\mathbb{R}}$  corresponding to  $\mathbb{R}$ -valued point  $z\overline{z}$ .

**Problem 2.** Let k be a field and  $\mathbb{G}_{a,k} = \operatorname{Spec}(k[t])$  with  $a^* : t \mapsto t \otimes 1 + 1 \otimes t$ .

a) We want to find maps  $\mathbb{G}_{a,k} \to \mathbb{G}_{a,k}$  that will respect operation a. Since  $\mathbb{G}_{a,k}$  is affine, this is equivalent to finding all maps  $f^* \colon k[t] \to k[t]$  such that

$$k[t] \otimes_k k[t] \xleftarrow{a^*} k[t]$$

$$f^* \otimes f^* \uparrow \qquad \qquad f^* \uparrow$$

$$k[t] \otimes_k k[t] \xleftarrow{a^*} k[t]$$

commutes. Since we are working over schemes over k, these are maps of k-algebras, which means that  $f^*$  is uniquely defined by its value at t. From commutativity we get the condition

$$f^*(t) \otimes 1 + 1 \otimes f^*(t) = f^*(t \otimes 1 + 1 \otimes t)$$

Writing  $f^*(t) = \sum_i a_i t^i$  we get

$$\sum_{i} a_i t^i \otimes 1 + 1 \otimes \sum_{i} a_i t^i = \sum_{i} a_i \sum_{j} \binom{i}{j} (t^{i-j} \otimes t^j)$$

Since char(k) = 0, none of the elements on right vanish. So by comparing terms we obtain  $a_i = 0$  for i > 1. So  $f^*(t) = a_0 + a_1 t$ .

$$(a_0 + a_1 t) \otimes 1 + 1 \otimes (a_0 + a_1 t) = a_0 + a_1 (t \otimes 1 + 1 \otimes t)$$

Compare again and get that  $2a_0 = 0$ , so  $a_0 = 0$  and  $a_1$  can be anything. Therefore endomorphisms  $\operatorname{End}(\mathbb{G}_{a,k})$  are parametrized by  $a_1 \in k$ . b) Let now char(k) = p. Same as before, but at the step when we have

$$\sum_{i} a_{i} t^{i} \otimes 1 + 1 \otimes \sum_{i} a_{i} t^{i} = \sum_{i} a_{i} \sum_{j} {i \choose j} (t^{i-j} \otimes t^{j})$$

when i is a power of p, all terms  $\binom{i}{j}$  vanish in k, but if i is not a power of p, then there exists j such that  $\binom{i}{j}$  does not vanish, meaning  $a_i$  has to be 0 for i not a power of p. We also get  $a_0=0$  by comparing terms. So  $f^*(t)=\sum_k a_{p^k}t^{p^k}$ . We see that all endomorphisms of  $\mathbb{G}_{a,k}$  are of this form and clearly all morphisms of above form in fact define endomorphisms of  $\mathbb{G}_{a,k}$