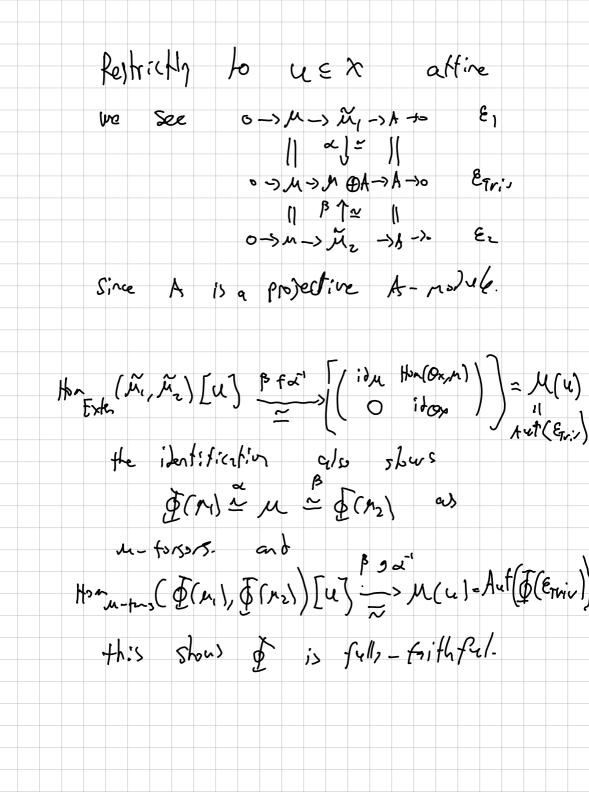
Proposition If H'(Sper A, I) =0 for any ring A Gold any quesir-herent shout I => formal smoothness is local on source. H' vs. Ext Let X be a schene and ME Q Coh. Definition - we let Extense be the Citesory of extensions where objects one exact sequences: 0 -> M -> M -> Ox ->0 and maps one 0 -> 1 -> M, -> 0x -> 0 id I idox 0 3 M -> Mz -> Ox ->s - We let Exf'(Ox, U) Inste the set of Isomorphism classes st extension.

Park Exti(Oxu) is a pinter set with the frivial extarian 0-) N-) MOOx -> Ox -> O being the distinstished element. (Late Ext'(Oxn) will be a group). Proposition: H'(x, u) ~ Ext'(Ox, u). Prost we construct an equivalence of categories

= = = M - forsers. If  $\varepsilon \sim [0 \rightarrow n \rightarrow \tilde{n} \stackrel{\rho_{\varepsilon}}{\rightarrow} 0_{\kappa} \rightarrow 0]$ \$( ε) [u] = \ se μ(u) (s) = 1 \ M acts on  $\overline{\Phi}(E')$   $G_7$ (m, s) 1......> m+s This well defined since pcm)=0.

Since taking global seeflows is left exact

o-> M(u) -> M(u) -> Ox(u) if there is  $t, s \in \mathcal{M}(a)$  with P(s) = 1 = P(t) then P(t-s) = 0 so t= s+m. Moreover, if u=speek is affine then  $0 \rightarrow M \rightarrow \tilde{M} \rightarrow A \rightarrow 0$ is exact so  $\Phi(M)$  (by  $\neq \phi$ ). Clain: \$ is fully-faithful: We have Homerton  $(\tilde{\mathcal{U}}_{i}, \tilde{\mathcal{U}}_{z}) \longrightarrow Hom_{Tors}(\tilde{\Phi}(\tilde{\mathcal{V}}_{i}), \tilde{\Phi}(\tilde{\mathcal{V}}_{z}))$ and 60th promote to sheaves ot To show of is an isomorphyn me can do this locally,



Essortal Sarjectivity: Objects in the atesomies Extern M-Tars. glue. The is it Use is a ser of x a G-forser p can be Sperified by finding ({Pui }iEL, ~is: Pui ) ui > Pui | ui 3.1. Likodij = dik. (similarly for Extern). Since Extens > en-tos is fully-fuithful to show it is essentially surjective car le pavel locallz. But locally every untorse is trivial and the frivial unforser is (0-> hes h@k-> k->0).

Towards Jerived functors Proposition Let T be a topological sple and 0->F1->F2->P3->0 an exact sequence of steames one T. Then une set an exact sequence 0-> 1'(T,F,)-> 1'(T,F2) -> 1'(T,F3)-> H'(T,F,) Post Given SET(T, F3) we attach an Fi-torsa- Ps defined as Ps: 4 -> { te fz(u) | f(t) = 9 | 4 } with F, -action siven by F, xF2 -> F2 (t,t2) ) (t, t2). We let S(s) = [Ps] & \$ Fi - tory - 5 \$ /2.  $S(s) = * \langle - \rangle P_s \simeq F' \langle - \rangle P_s(T) \neq \emptyset$ S e f(H(T,F3)).

Definition A function of abelian cutegories fic-)D is additive if # ABEC the map of soups. (F(A), F(B))
is a map of soups. Definition: Let F: C > D be a left exact functor between abelian categories. A cohomological &-function extendin, F ;) a sequence of additive function fi: C -> D tosether with fractist maps Si: Fi(c) >> Pi+1(A) A all short exact sequences  $0 \rightarrow A \rightarrow B \rightarrow C \rightarrow 0.$  s.  $\}.$ -.. 5i-1 Fi(A) -> Fi(B) -> Fi(c) 5h Fi+(A)-1... is exact. and Fo=5

Ruk Functoriality of si mens tkt tor 0-) A-) B-) C->> 2 / / 0-) A'-> B'-> C'->> the indeb diagram Fi(c) & Fill(A)  $F^{i}(C') \xrightarrow{S^{i}} F^{ifl}(A^{l})$ Commutes. we can form a tesson st deta (Fi, Si) of netural transformations compatible with S. Definition Let F: E-> D be a left exact functor. A coh. &-functor extending F is universal if it is initial in the catesoly at coh. S-Fark

Definition: Let F: C-> D be lett exict (Fi, si) extending F is erasable if tiz1 the e there exists a monomorphism  $A \longrightarrow B$  s.t.  $\mathcal{F}^{i}(\mathcal{B}) = 0$ . Theorem Every ersable S-furcher extending & is universal. Prot Supple (Fi, Si) is erasable and (Gi, Gi) also extends J-. We want to constract unique \$i: fi -> ci. we induct on i 9°: 7° -> 6° F 18 11 Given AEC there is BEC w:fh 0→ A → B → C → 0 and . 2, (B) =0

We have by inaction -> Fi-1 (B) -> Fi(c) -> Fi(A) -> Fi(B)=0  $-> G^{i-1}(B) -> G^{i-1}(C) -> G^{i}(A) -> G^{i}(B)$ Since ( o o Fi-(6) =0 +his produces a unique map \$ : F^(A) -> G'(A). Book Keepings: a) The construction of pri doesn't repend on P. b) \$\dagger{gaing} \text{\$\text{\$\text{\$\interlightarrow}} \text{\$\end{\end{\end{\$\text{\$\text{\$\text{\$\text{\$\text{\$\text{\$\text{\$\exittit{\$\tex{\$\$\text{\$} functoris! c) We have communitie pin(L) { for exact server Gi(L) -> Ci(k) 0-> k-> -> L->0

