

Sebastiaan_Bye_Julius_Watenaar_homework2

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1 Homework Set 2

2 Exercise 0

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```
[3]: import numpy as np
import scipy.linalg as la
import matplotlib.pyplot as plt
import random
```

3 Exercise 1

The goal of this problem is to show that apparently harmless looking systems of linear equations may be very difficult to solve. Some functions that may be useful are `numpy.triu`, `numpy.tril`, `numpy.eye`, `random.randrange`. ## (a) Generate an $n \times n$ matrix B with random integer elements in the range $b_{ij} \in [-10, 10]$. Choose for instance $n = 20$.

```
[4]: n = 20
B = np rint(np.random.rand(n,n)*20 - 10)
```

The cell above generates a 20×20 matrix with integer elements in the range $b_{ij} \in [-10, 10]$

3.1 (b)

Remove the diagonal of B , save the upper triangular part in U and the lower triangular part in L , and put ones on the diagonals $l_{ii} = u_{ii} = 1$.

```
[5]: L = np.zeros((n,n))
U = np.zeros((n,n))
for row in range(n):
    for column in range(n):
        if row == column:
            # diagonal
            B[row][column] = 0

for row in range(n):
```

```

for column in range(n):
    if row < column:
        # lower
        L[row][column] = B[row][column]
    elif row > column:
        #upper
        U[row][column] = B[row][column]

np.fill_diagonal(L, 1)
np.fill_diagonal(U, 1)

```

Assuming that removing the diagonal of B means setting the elements to 0, the first loop removes the diagonal of B . The second loop saves the elements of the triangular matrices and the final command sets the diagonals to 1.

3.2 (c)

Compute $A = L \cdot U$. What is the value of $\det(A)$ and why? Compute the determinant using the appropriate python command and confirm your prediction. In case that you have doubts about the result, compute separately $\det(L)$ and $\det(U)$.

```
[6]: A = np.matmul(L,U)
      np.linalg.det(A)
```

```
[6]: 804835905.521148
```

The above cell computes the determinant of A which is in this case 804835905.521148. This appears to be wrong. The determinant should be 1 as the determinants of L and U are one and $\det(A) = \det(L) * \det(U)$. This happens as finding the determinant is prone to error.

```
[7]: print(np.linalg.det(L),np.linalg.det(U))
```

```
1.0 0.9999986129845544
```

The cell above computes the determinates of L and U seperately. They are respectively 1.0 and 0.9999986129845544. This result is in line with the theory as the determinant of a triangular matrix is the product the diagonal. In this case the elements of the diagonal are all 1 hence, this result is correct.

```
[8]: np.linalg.slogdet(A)
```

```
[8]: (1.0, 20.506148970530703)
```

The above cell uses the natural log to compute the determinant of A . Now the determinant is 1.

3.3 (d)

Choose now an exact solution, for instance $x_e = \text{numpy.ones}(n)$, and compute the corresponding right hand side $b = Ax_e$.

```
[9]: b = np.matmul(A, np.ones(n))
```

```
b
```

```
[9]: array([ 636.,  331.,  320.,  326.,  669., -406.,  325.,  -86.,  548.,  
          156., -251., -236., -422., -349.,  -86.,   62., -189., -297.,  
          -136.,  -27.])
```

The above cell generates a vector with all elements one and calculates what b is. The resulting vector b is

```
[ 636.  
  331.  
  320.  
  326.  
  669.  
 -406.  
  325.  
  -86.  
  548.  
  156.  
 -251.  
 -236.  
 -422.  
 -349.  
  -86.  
   62.  
 -189.  
 -297.  
 -136.  
  -27.]
```

3.4 (e)

Solve $Ax = b$ using `scipy.linalg.lu_factor` and `scipy.linalg.lu_solve` and compare the solution with the exact x_e .

```
[10]: lu, piv = la.lu_factor(A)
```

```
[11]: la.lu_solve((lu, piv), b)
```

```
[11]: array([ 1.          ,  1.          ,  1.          ,  1.          ,  1.          ,  
            0.99999998,  0.99999983,  0.99999869,  0.99999922,  0.99999149,  
            1.00003146,  1.00018874,  0.99983995,  0.99855031,  0.99709845,  
            1.02101795,  1.10214885,  1.86904292, -1.00410593,  6.41994475])
```

Using the suggested libraries, the solution of x_e is

```

1.
1.
1.
1.
1.
0.99999998
0.99999983
0.99999869
0.99999922
0.99999149
1.00003146
1.00018874
0.99983995
0.99855031
0.99709845
1.02101795
1.10214885
1.86904292
-1.00410593
6.41994475

```

3.5 (f)

Explain the bad results by computing the condition number of A .

```
[12]: np.linalg.cond(A)
```

```
[12]: 1.2260241100456064e+18
```

The condition number is very high. It is $1.2260241100456064e + 18$. This means that the matrix is very sensitive to input errors and that can be seen in the last 5 entries of the vector in exercise (e).

4 Exercise 2

4.1 (a)

Verify that

$$\begin{bmatrix} A & B \\ C & D \end{bmatrix} = \begin{bmatrix} I_p & 0 \\ CA^{-1} & I_q \end{bmatrix} \begin{bmatrix} A & 0 \\ 0 & D - CA^{-1}B \end{bmatrix} \begin{bmatrix} I_p & A^{-1}B \\ 0 & I_q \end{bmatrix}$$

The first step:

$$\begin{bmatrix} I_p & 0 \\ CA^{-1} & I_q \end{bmatrix} \begin{bmatrix} A & 0 \\ 0 & D - CA^{-1}B \end{bmatrix} = \begin{bmatrix} (I_p * A) + (0 * 0) & (I_p * 0) + (0 * (D - CA^{-1}B)) \\ (CA^{-1} * A) + (I_q * 0) & (CA^{-1} * 0) + (I_q * (D - CA^{-1}B)) \end{bmatrix} = \begin{bmatrix} A & 0 \\ C & D - CA^{-1}B \end{bmatrix}$$

The second step:

$$\begin{bmatrix} A & 0 \\ C & D - CA^{-1}B \end{bmatrix} \begin{bmatrix} I_p & A^{-1}B \\ 0 & I_q \end{bmatrix} = \begin{bmatrix} (A * I_p) + (0 * 0) & (A * A^{-1}B) + (0 * I_q) \\ (C * I_p) + (0 * (D - CA^{-1}B)) & (C * A^{-1}B) + (D - CA^{-1}B * I_q) \end{bmatrix} = \begin{bmatrix} A & B \\ C & D \end{bmatrix}$$

Thus

$$\begin{bmatrix} A & B \\ C & D \end{bmatrix} = \begin{bmatrix} I_p & 0 \\ CA^{-1} & I_q \end{bmatrix} \begin{bmatrix} A & 0 \\ 0 & D - CA^{-1}B \end{bmatrix} \begin{bmatrix} I_p & A^{-1}B \\ 0 & I_q \end{bmatrix}$$

4.2 (b)

Describe how a system $Mx = b$, with x and b in \mathbb{R}^{p+q} , can be solved by applying matrix-vector products with C and B and solves with A and $(D - CA^{-1}B)$.

$$Mx = b$$

$$\begin{bmatrix} A & B \\ C & D \end{bmatrix} x = b$$

$$\begin{bmatrix} A & B \\ C & D \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \end{bmatrix}$$

From the question it is known that A is invertible. Hence, x_1 can be eliminated by

$$Ax_1 + Bx_2 = b_1$$

$$x_1 = A^{-1}(b_1 - Bx_2).$$

For the second equation

$$Cx_1 + Dx_2 = b_2$$

$$C(A^{-1}(b_1 - Bx_2)) + Dx_2 = b_2$$

$$(D - CA^{-1}B)x_2 = b_2 - CA^{-1}b_1.$$

Now let $S = D - CA^{-1}B$, then

$$x_2 = S^{-1}(b_2 - CA^{-1}b_1).$$

Substituting this expression into the expression for x_1 gives

$$x_1 = A^{-1}(b_1 - B(S^{-1}(b_2 - CA^{-1}b_1)))$$

$$x_1 = (A^{-1} + A^{-1}BS^{-1}CA^{-1})b_1 - A^{-1}BS^{-1}b_2$$

.

Now $Mx = b$ can be solved with x, b by

$$\begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} A^{-1} + A^{-1}BS^{-1}CA^{-1} & -A^{-1}BS^{-1} \\ -S^{-1}CA^{-1} & S^{-1} \end{bmatrix} \begin{bmatrix} b_1 \\ b_2 \end{bmatrix}$$

4.3 (c)

What is the cost, to highest order, of LU-factorizing A and of computing and LU-factorizing $D - CA^{-1}B$?

The cost to the highest order of LU-factorizing A is as follows:

Eliminating the first column requires p additions and p multiplications on $p-1$ rows. Thus $2p(p-1)$ operations.

For the second column there are $p-1$ additions and $p-1$ multiplications on $p-2$ rows. Thus $2(p-1)(p-2)$ operations.

Thus $\sum_{i=1}^p 2(p-1)(p-i+1)$ operations per column. Now let $j = p-i+1$.

$$2 \sum_{j=0}^{p-1} j(j+1)$$

$$2 \sum_{j=0}^{p-1} (j^2 + j)$$

$$2\left(\frac{p^3}{3} - \frac{p}{3}\right)$$

$$\frac{2p^3}{3} - \frac{2p}{3}.$$

So in the worst case the cost of LU-factorizing A is $\frac{2p^3}{3} - \frac{2p}{3}$. So the cost to the highest order of LU-factorizing A is $\frac{2p^3}{3}$.

Computing $D - CA^{-1}B$. Inverting a matrix has a time-complexity of $O(n^3)$, in this case $O(p^3)$.

First, compute CA^{-1} . C is $q \times p$ and A^{-1} is $p \times p$, which computing the matrix multiplication is of order $O(pqp) = O(p^2q)$. This is because adding and multiplying p terms for each pq element. The resulting matrix is of size $q \times p$.

Second, compute CA^{-1} times B which is $q \times p$ times $p \times q$. This has a time complexity of $O(qpq) = O(q^2p)$. This is because adding and multiplying q terms for each qp element. The resulting matrix is of size $q \times q$.

The final operation is subtracting a $q \times q$ matrix from a $q \times q$ matrix which is a total of q^2 operations.

The total time complexity is $O(p^3 + p^2q + q^2p + q^2)$.

The LU-factorizing $D - CA^{-1}B$ cost follows the same logical as the LU-factorizing cost of A

Thus $\sum_{i=1}^q 2(q-1)(q-i+1)$ operations per column. Now let $j = q-i+1$.

$$2 \sum_{j=0}^{q-1} j(j+1)$$

$$2 \sum_{j=0}^{q-1} (j^2 + j)$$

$$2 \left(\frac{q^3}{3} - \frac{q}{3} \right)$$

$$\frac{2q^3}{3} - \frac{2q}{3}.$$

So in the worst case the cost of LU-factorizing $D - CA^{-1}B$ is $\frac{2q^3}{3} - \frac{2q}{3}$. So the cost to the highest order of LU-factorizing $D - CA^{-1}B$ is $\frac{2q^3}{3}$.

[]: