

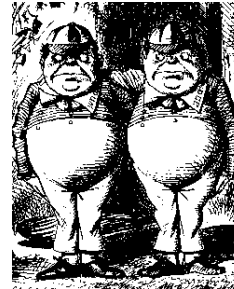
Introduction to Cryptanalysis

Good Guys and Bad Guys

- Alice and Bob are the good guys



- Trudy is the bad guy



- Trudy is our generic “intruder”

Good Guys and Bad Guys

- ❑ Alice and Bob want to communicate securely
 - Typically, over a network
- ❑ Alice or Bob might also want to store their data securely
- ❑ Trudy wants to read Alice and Bob's secrets
- ❑ Or Trudy might have other devious plans...
 - Cause confusion, denial of service, etc.

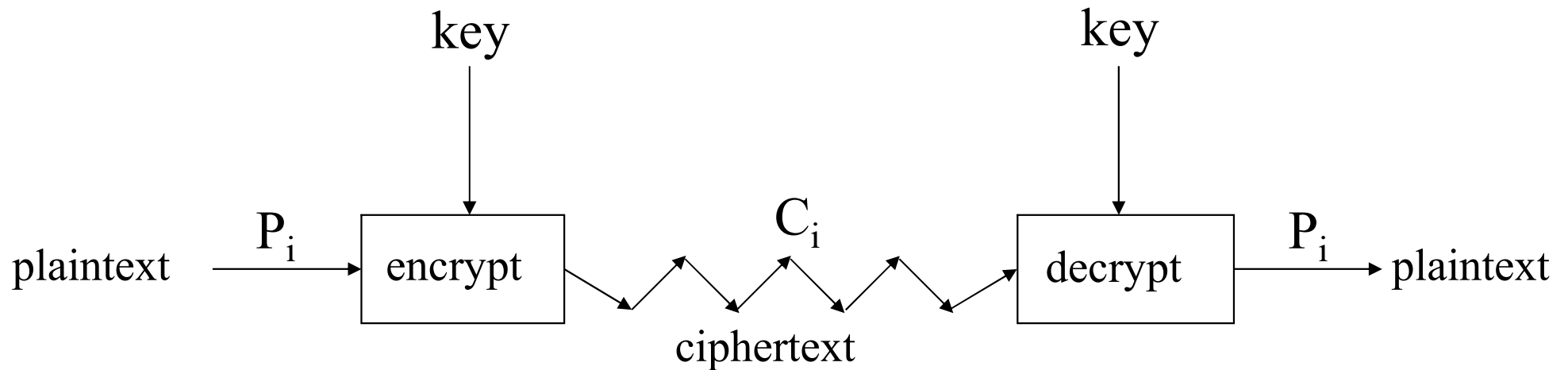
CIA

- ❑ Confidentiality, Integrity and Availability
- ❑ **Confidentiality:** prevent unauthorized reading of information
- ❑ **Integrity:** prevent unauthorized writing of information
- ❑ **Availability:** data is available in a timely manner when needed
 - Availability is a “new” security concern
 - Due to denial of service (DoS) threats

Crypto

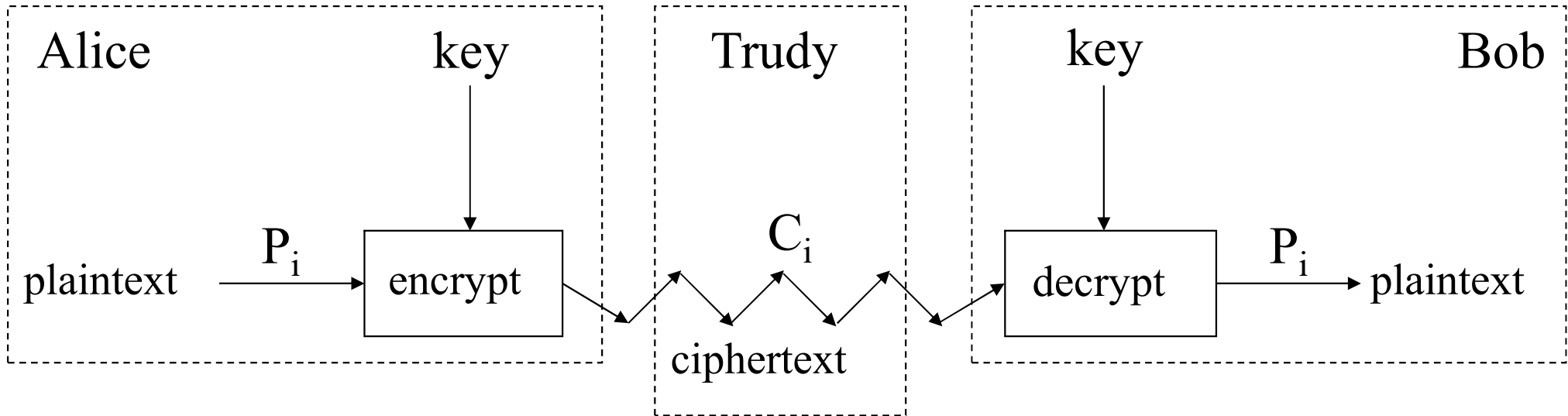
- ❑ **Cryptology** — The art and science of making and breaking “secret codes”
- ❑ **Cryptography** — making “secret codes”
- ❑ **Cryptanalysis** — breaking “secret codes”
- ❑ **Crypto** — all of the above (and more)

Crypto as a Black Box



- ❑ Note P_i is i^{th} “unit” of plaintext
- ❑ And C_i is corresponding ciphertext
- ❑ “Unit” may be bit, letter, block of bits, etc.

Who Knows What?



- ❑ Trudy knows the ciphertext
- ❑ Trudy knows the cipher and how it works
- ❑ Trudy might know a little more
- ❑ Trudy does **not** know the key

Cryptanalysis

- ❑ This course focused on cryptanalysis
- ❑ Trudy wants to recover key or plaintext
- ❑ Trudy is not bound by any rules
 - For example, Trudy might attack the implementation, not the algorithm itself
 - She might use “side channel” info, etc.

Attacking Block Ciphers

- ❑ Standard attacks
 - exhaustive key search
 - dictionary attack
 - differential cryptanalysis
 - linear cryptanalysis
- ❑ Side channel attacks against implementations.
 - Timing attacks
 - Power consumption attacks
 - Fault injection attacks

Exhaustive Key Search

- ❑ How can Trudy attack a cipher?
- ❑ She can simply try all possible keys and test each to see if it is correct
 - Exhaustive key search
- ❑ To prevent an exhaustive key search, a cryptosystem must have a large **keyspace**
 - Must be too many keys for Trudy to try them all in any reasonable amount of time

Beyond Exhaustive Search

- ❑ A large key space is necessary for security
- ❑ But a large key space is not sufficient
- ❑ Shortcut attacks might exist
- ❑ We'll see many examples of shortcut attacks
- ❑ In cryptography we can (almost) never prove that no shortcut attack exists
- ❑ This makes cryptography interesting...

Chosen-Plaintext Dictionary Attacks Against Block Ciphers

- ❑ Construct a table with the following entries
 - $(K, E_K[0])$ for all possible key K
 - Sort based on the second field (ciphertext)
 - How much time does this take?
- ❑ To attack a new key K (under chosen message attacks)
 - Choose 0, obtain the ciphertext C , looks up in the table, and finds the corresponding key
 - How much time does this step take?
- ❑ Trade off space for time

Differential Cryptanalysis

□ Main idea:

- This is a **chosen plaintext attack**,
- The attacker knows many (plaintext, ciphertext) pairs
- Difference $\Delta_P = P_1 \oplus P_2$, $\Delta_C = C_1 \oplus C_2$
- **Distribution of Δ_C 's given Δ_P may reveal information about the key (certain key bits)**
- After finding several bits, use brute-force for the rest of the bits to find the key.

Basic idea of linear cryptanalysis

- Suppose that
- (*) $\Pr [\begin{array}{l} M_{i1} \oplus M_{i2} \oplus \dots \oplus M_{iu} \\ \oplus C_{j1} \oplus C_{j2} \oplus \dots \oplus C_{jv} \\ \oplus K_{p1} \oplus K_{p2} \oplus \dots \oplus K_{pw} = 1 \end{array}] = 0.5 + \varepsilon$
- Then one can recover some key bits given large number of PT/CT pairs
- For DES, exists (*) with $\varepsilon=2^{-21}$
- Using this method, one can find 14 key bits using $(2^{21})^2$ PT/CT pairs

DES Strength Against Various Attacks

Attack Method	Known	Chosen	Storage complexity	Processing complexity
Exhaustive precomputation	-	1	2^{56}	1
Exhaustive search	1	-	negligible	2^{55}
Linear cryptanalysis	2^{43} 2^{38}	- -	For texts	2^{43} 2^{50}
Differential cryptanalysis	- 2^{55}	2^{47} -	For texts	2^{47} 2^{55}

The weakest point of DES remains the size of the key (56 bits)!

Taxonomy of Cryptanalysis

- ❑ Ciphertext only — always an option
- ❑ Known plaintext — possible in many cases
- ❑ Chosen plaintext
 - “Lunchtime attack”
 - Protocols might encrypt chosen text
- ❑ Adaptively chosen plaintext
- ❑ Related key
- ❑ Forward search (public key crypto only)
- ❑ “Rubber hose”, bribery, etc., etc., etc.

Definition of Secure

- ❑ A cryptosystem is **secure** if the best know attack is to try all possible keys
- ❑ Cryptosystem is **insecure** if **any** shortcut attack is known
- ❑ By this definition, an insecure system might be harder to break than a secure system!

Definition of Secure

- ❑ Why do we define **secure** this way?
- ❑ The size of the keyspace is the “advertised” level of security
- ❑ If an attack requires less work, then false advertising
- ❑ A cipher must be secure (by our definition) and have a “large” keyspace
 - Too big for an exhaustive key search

Theoretical Cryptanalysis

- ❑ Suppose that a cipher has a 100 bit key
 - Then keyspace is of size 2^{100}
- ❑ On average, for exhaustive search Trudy tests $2^{100}/2 = 2^{99}$ keys
- ❑ Suppose Trudy can test 2^{30} keys/second
 - Then she can find the key in about 37.4 trillion years

Theoretical Cryptanalysis

- ❑ Suppose that a cipher has a 100 bit key
 - Then keyspace is of size 2^{100}
- ❑ Suppose there is a shortcut attack with “work” equal to testing about 2^{80} keys
- ❑ If Trudy can test 2^{30} per second
 - Then she finds key in 36 million years
 - Better than 37 trillion, but not practical

Applied Cryptanalysis

- ❑ In this class, we focus on attacks that produce plaintext
 - Not interested in attacks that just show a theoretical weakness in a cipher
- ❑ We call this **applied cryptanalysis**
- ❑ Why applied cryptanalysis?
 - Because it's a lot more fun...
 - And it's a good place to start

Applied Cryptanalysis: Overview

- ❑ Classic (pen and paper) ciphers
 - Transposition, substitution, etc.
 - Same principles appear in later sections
- ❑ World War II ciphers
 - Enigma, Purple, Sigaba
- ❑ Stream ciphers
 - Shift registers, correlation attack, ORYX, RC4, PKZIP

Applied Cryptanalysis: Overview

❑ Block ciphers

- Hellman's TMTO, CMEA, Akelarre, FEAL

❑ Hash functions

- Nostradamus attack, MD4, MD5

❑ Public key crypto

- Knapsack, Diffie-Hellman, Arithmetica, RSA, Rabin, NTRU, ElGamal
- Factoring, discrete log, timing, glitching

Side Channel Analysis

- Time
 - Does the number of CPU cycles depend on exact values used in the operation? ex. RSA exponent
 - Memory access – do exact values impact tables used, time to read from a table and/or number of memory accesses? ex. AES using tables of 32-bit values
 - Acoustics
 - Impacted by operations or exact values used?
 - Memory
 - Can intermediate values be read from memory by another process?
-

Timing - Toy Example

k: array of n key bytes

d: 16 byte data

Suppose encryption is a series of n rounds

n = 16;

d = plaintext;

for (i=0; i < n; ++i) {

 d = f(d,k[i]); // do something to the data with k, but
 // whose time does not depend on k

 d[i] = d[i] ^{int(k[i])} mod 256; // alter one byte, time depends on k

}

Timing - Toy Example

What if use a table lookup instead?

table(a,b): function retrieves table a, entry b

```
d = plaintext;
```

```
x = 0;
```

```
for (i=0; i < n; ++i) {
```

```
    // do something to the data with k where time does not depend on k
```

```
    d = f(d,k[i]);
```

```
    // memory lookup - was table already in cache?
```

```
    // (k[i] same as a previous key byte)
```

```
    x= table(k[i], d[i]);
```

```
}
```

Why Study Cryptography?

- ❑ Information security is a big topic
 - Crypto, Access control, Protocols, Software
 - Real world info security problems abound
- ❑ Cryptography is the part of information security that works best
- ❑ Using crypto correctly is important
- ❑ The more we make other parts of security behave like crypto, the better

Why Study Cryptanalysis?

- ❑ Study of cryptanalysis gives insight into all aspects of crypto
- ❑ Gain insight into attacker's mindset
 - “black hat” vs “white hat” mentality
- ❑ Cryptanalysis is more fun than cryptography
 - Cryptographers are boring
 - Cryptanalysts are cool
- ❑ But cryptanalysis is hard