



Online Learning

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LECTURE 13

Slides adapted from Mohri

Motivation

- PAC learning: distribution fixed over time (training and test), IID assumption.
- On-line learning:
 - no distributional assumption.
 - worst-case analysis (adversarial).
 - mixed training and test.
 - Performance measure: mistake model, regret.

General Online Setting

- For $t = 1$ to T :
 - Get instance $x_t \in X$
 - Predict $\hat{y}_t \in Y$
 - Get true label $y_t \in Y$
 - Incur loss $L(\hat{y}_t, y_t)$
- Classification: $Y = \{0, 1\}$, $L(y, y') = |y' - y|$
- Regression: $Y \subset \mathbb{R}$, $L(y, y') = (y' - y)^2$

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- **Objective:** Minimize total loss $\sum_t L(\hat{y}_t, y_t)$

Plan

Experts

Perceptron Algorithm

Prediction with Expert Advice

- For $t = 1$ to T :
 - Get instance $x_t \in X$ and advice $a_t, i \in Y, i \in [1, N]$
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- **Objective:** Minimize regret, i.e., difference of total loss vs. best expert

$$\text{Regret}(T) = \sum_t L(\hat{y}_t, y_t) - \min_i \sum_t L(a_{t,i}, y_t) \quad (1)$$

Mistake Bound Model

- Define the maximum number of mistakes a learning algorithm L makes to learn a concept c over any set of examples (until it's perfect).

$$M_L(c) = \max_{x_1, \dots, x_T} |\text{mistakes}(L, c)| \quad (2)$$

- For any concept class C , this is the max over concepts c .

$$M_L(C) = \max_{c \in C} M_L(c) \quad (3)$$

Halving Algorithm

```
 $H_1 \leftarrow H;$   
for  $t \leftarrow 1 \dots T$  do  
    Receive  $x_t$ ;  
     $\hat{y}_t \leftarrow \text{Majority}(H_t, \vec{a}_t, x_t);$   
    Receive  $y_t$ ;  
    if  $\hat{y}_t \neq y_t$  then  
        |  $H_{t+1} \leftarrow \{c \in H_t : c(x_t) = y_t\};$   
return  $H_{T+1}$ 
```

Algorithm 1: The Halving Algorithm (Mitchell, 1997)

Halving Algorithm Bound (Littlestone, 1998)

- For a finite hypothesis set

$$M_{\text{Halving}(H)} \leq \lg |H| \quad (4)$$

- After each mistake, the hypothesis set is reduced by at least by half

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- For a finite hypothesis set

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- After each mistake, the hypothesis set is reduced by at least by half
- Consider the optimal mistake bound $\text{opt}(H)$. Then

$$\text{VC}(H) \leq \text{opt}(H) \leq M_{\text{Halving}(H)} \leq \lg |H| \quad (5)$$

- For a fully shattered set, form a binary tree of mistakes with height $\text{VC}(H)$

Weighted Majority (Littlestone and Warmuth, 1998)

```

for  $t \leftarrow 1 \dots N$  do
  |  $w_{1,i} \leftarrow 1$ ;
for  $t \leftarrow 1 \dots T$  do
  | Receive  $x_t$ ;
  |  $\hat{y}_t \leftarrow \mathbb{1} \left[ \sum_{y_{t,i}=1} w_t \geq \sum_{y_{t,i}=0} w_t \right]$ ;
  | Receive  $y_t$ ;
  | if  $\hat{y}_t \neq y_t$  then
  |   | for  $t \leftarrow 1 \dots N$  do
  |     | if  $\hat{y}_t \neq y_t$  then
  |       |  $w_{t+1,i} \leftarrow \beta w_{t,i}$ ;
  |     | else
  |       |  $w_{t+1,i} \leftarrow w_{t,i}$ 
return  $w_{T+1}$ 

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- **Weights for every expert**
- Classifications in favor of side with higher total weight ($y \in \{0, 1\}$)
- Experts that are wrong get their weights decreased ($\beta \in [0, 1]$)
- If you're right, you stay unchanged

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Weighted Majority

- Let m_t be the number of mistakes made by WM until time t
- Let m_t^* be the best expert's mistakes until time t

$$m_t \leq \frac{\log N + m_t^* \log \frac{1}{\beta}}{\log \frac{2}{1+\beta}} \quad (6)$$

- Thus, mistake bound is $O(\log N)$ plus the best expert
- Halving algorithm $\beta = 0$

Proof: Potential Function

- Potential function is the sum of all weights

$$\Phi_t \equiv \sum_i w_{t,i} \quad (7)$$

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Weights are nonnegative, so $\sum_i w_{t,i} \geq w_{t,i}$

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Each error multiplicatively reduces weight by β

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- If an algorithm makes an error at round t

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$$\Phi_1 = N \quad (10)$$

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$$\Phi_1 = N \quad (10)$$

- After m_T mistakes after T rounds

$$\Phi_T \leq \left[\frac{1 + \beta}{2} \right]^{m_T} N \quad (11)$$

Weighted Majority Proof

- Put the two inequalities together, using the best expert

$$\beta^{m_T^*} \leq \Phi_T \leq \left[\frac{1 + \beta}{2} \right]^{m_T} N \quad (12)$$

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- Take the log of both sides

$$m_T^* \log \beta \leq \log N + m_T \log \left[\frac{1 + \beta}{2} \right] \quad (13)$$

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$$m_T^* \log \beta \leq \log N + m_T \log \left[\frac{1 + \beta}{2} \right] \quad (13)$$

- Solve for m_T

$$m_T \leq \frac{\log N + m_T^* \log \frac{1}{\beta}}{\log \left[\frac{2}{1 + \beta} \right]} \quad (14)$$

Weighted Majority Recap

- Simple algorithm
- No harsh assumptions
- Depends on best learner
- Downside: Takes a long time to do well in worst case (but okay in practice)
- Solution: Randomization

Plan

Experts

Perceptron Algorithm

Perceptron Algorithm

- Online algorithm for classification
- Very similar to logistic regression (but 0/1 loss)
- But what can we prove?

Perceptron Algorithm

```
 $\vec{w}_1 \leftarrow \vec{0};$   
for  $t \leftarrow 1 \dots T$  do  
    Receive  $x_t$ ;  
     $\hat{y}_t \leftarrow \text{sgn}(\vec{w}_t \cdot \vec{x}_t)$ ;  
    Receive  $y_t$ ;  
    if  $\hat{y}_t \neq y_t$  then  
        |  $\vec{w}_{t+1} \leftarrow \vec{w}_t + y_t \vec{x}_t$ ;  
    else  
        |  $\vec{w}_{t+1} \leftarrow \vec{w}_t$ ;  
return  $\vec{w}_{T+1}$ 
```

Algorithm 2: Perceptron Algorithm (Rosenblatt, 1958)

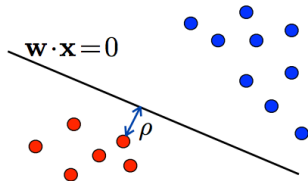
Objective Function

- Optimizes

$$\frac{1}{T} \sum_t \max(0, -y_t(\vec{w} \cdot x_t)) \quad (15)$$

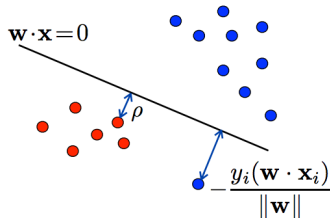
- Convex but not differentiable

Margin and Errors



- If there's a good margin ρ , you'll converge quickly

Margin and Errors



- If there's a good margin ρ , you'll converge quickly
- Whenever you see an error, you move the classifier to get it right
- Convergence only possible if data are separable

How many errors does Perceptron make?

- If your data are in a R ball and there is a margin

$$\rho \leq \frac{y_t(\vec{v} \cdot \vec{x}_t)}{\|\vec{v}\|} \quad (16)$$

for some \vec{v} , then the number of mistakes is bounded by R^2/ρ^2

- The places where you make an error are support vectors
- Convergence can be slow for small margins

Why study Perceptron?

- Simple algorithm
- Bound independent of dimension and tight
- Foundation of deep learning
- Proof techniques helped usher in SVMs
- Generalizes to structured prediction