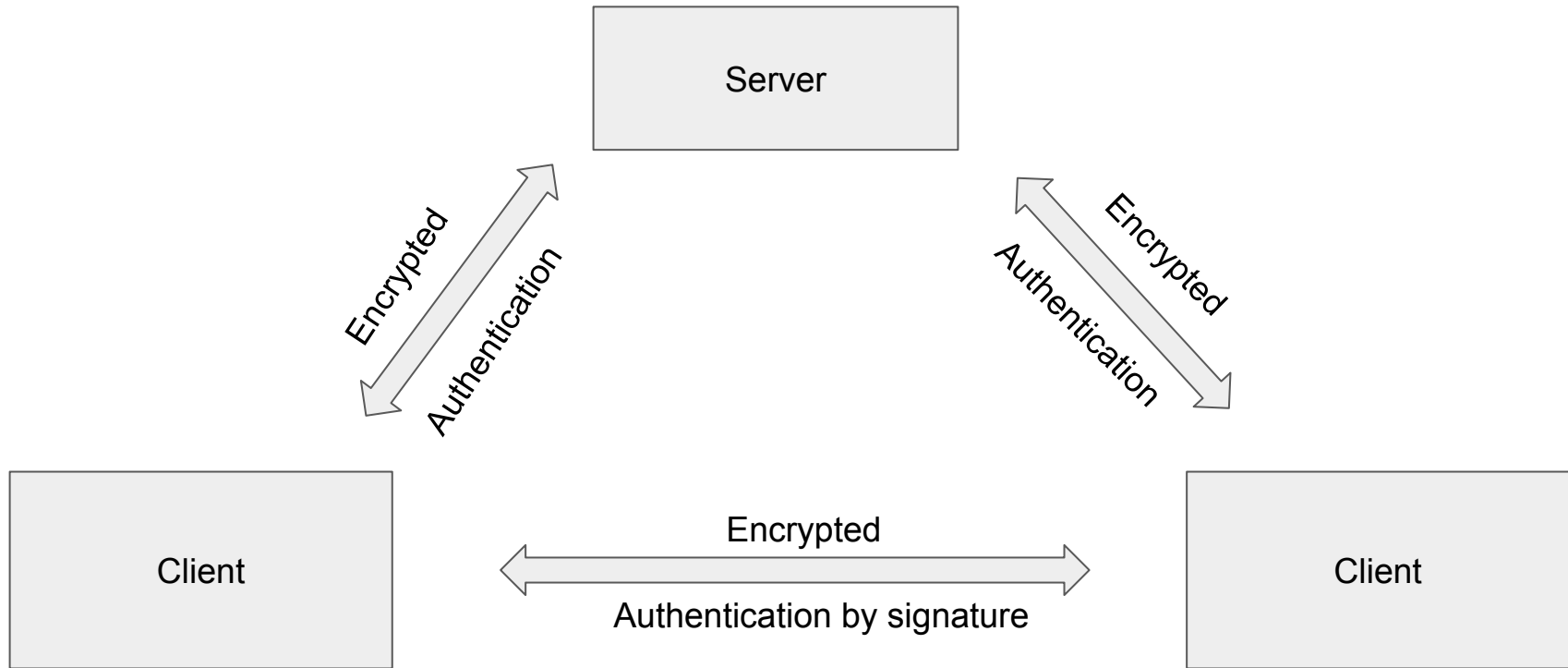


Secure Instant Messaging System

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Architecture



Before Authentication (assumptions)

Client A:

- Password P_a

Client can compute: W, W'

Server:

- W : derived from P_a using PBKDF2
- A's public key
- $Y = W' \{A's\ private\ key\}$

Note: All symmetric encryption use AES-GCM to provide confidentiality and data integrity

During Client-Server Authentication

Augmented Strong Password Protocol:

A->S: "Alice"

S->A: $c = \text{Hash}(A's\ ip, secret)$

A->S: $c, W\{g^a \bmod p\}$

S->A: $W\{g^b \bmod p\}, (g^{ab} \bmod p)\{W'\{A's\ private\ key\}\}, c1$

A->S: $[\text{hash}(g^{ab} \bmod p, c1)]\text{sign-A}$

Client to Server Communication (after authentication)

Both party: Session key: $K_s = g^{ab} \bmod p$

Client: private RSA key

A -> Server: $K_s\{\text{"List"}, \text{Timestamp}\}$

Server -> A: $K_s\{\text{"List of log in users"}, \text{Timestamp}\}$

A -> Server: $K_s\{\text{"Tell me about B"}, \text{Timestamp}\}$

Server -> A: $K_s\{B\text{'s public key}, B\text{'s ip address}, \text{Timestamp}\}$

A -> Server: $K_s\{\text{"Log me out"}, \text{Timestamp}\}$

Protocol Security

1. Mutual authentication.
2. Resistant to offline attack
3. If the server is compromised, hard to brute-force the password
4. If the server is compromised, attacker cannot impersonate the user, since he does not know W' .
5. Resistant to DoS attack by using cookies.
6. Resistant to man-in-the-middle by encrypting with W .
7. Perfect Forward Secrecy using Diffie-Hellman
8. Resistant to replay attack by using timestamp
9. Provides confidentiality and integrity using AES-GCM

Client to Client Key Establishment

A gets B's ip address from the server

A→B: $g^a \bmod p$

B→A: $g^b \bmod p$

A→B: $g^{ab} \bmod p$ {“A”, $[g^a \bmod p]\text{sign-A}$ }

B verifies the signature using A's public key retrieved from the server

B→A: $g^{ab} \bmod p$ {“B”, $[g^b \bmod p]\text{sign-B}$ }

A verifies the signature using B's public key retrieved from the server

Client to Client Communication

Both party: Session key: $K_s = g^{ab} \bmod p$

A -> B: $K_s\{\text{"message"}, \text{Timestamp}\}$

B-> A: $K_s\{\text{"message"}, \text{Timestamp}\}$

Protocol Security

1. Resistant to man-in-the-middle by using signature
2. Identity hiding
3. Perfect forward Secrecy (Deffie-hellmen)
4. Mutual authentication by signature
5. Provide both confidentiality and data integrity by using AES-GCM
6. Resistant to replay attack by using timestamp.