

## 1 Sets

**Sets** are collections of values of one type with no internal structure. Sets have no repeated items and order does not matter.

### 1.1 Set Operations

- Union -  $A \cup B = \{x : x \in A \text{ or } x \in B\}$
- Intersect -  $A \cap B = \{x : x \in A \text{ and } x \in B\}$
- Concatenation -  $AB = \{ab : a \in A \text{ and } b \in B\}$
- Cartesian Product -  $A \times B = \{(a, b) : a \in A \text{ and } b \in B\}$
- Complement -  $A^c = \{x : x \notin A\}$
- Powerset -  $P(A) = \{A' : A' \subset A\}$

### 1.2 Finite and Infinite Sets

Set  $A$  is **finite** if it can be put in 1-to-1 correspondence with an initial segment of  $\mathbb{N}$ , or is the empty set  $\emptyset$ . More intuitively, a set is finite if we can list the elements of  $A$  in finite time.

Set  $A$  is **countably infinite** if it can be put in 1-to-1 correspondence with *all* of  $\mathbb{N}$ . Intuitively, a procedure can be devised to list all elements, but this procedure will never finish.

Set  $A$  is **uncountably infinite** if it is not countable.

### 1.3 Relations and Functions

A **function** can be defined as a binary operation between two sets that given an input in set  $A$  produces an output in set  $B$ . There are three types of functions:

1. Total function:  $A \rightarrow B$  means for every input  $A$  there is exactly one output  $B$
2. Partial function:  $A \rightarrow B$  means for every input  $A$  there is at most one output  $B$
3. Multi function:  $A \multimap B$  means for every input  $A$  there are 0, 1, or many outputs in  $B$

We also define the **identity function** such that for any type  $A$ ,  $f(A) = A$ .

### 1.4 Cardinality of Sets

The cardinality of a set  $A$ , or  $|A|$ , is the total number of elements in  $A$ .

- Union -  $|A \cup B| = |A| + |B| - |A \cap B|$
- Intersect -  $|A \cap B| = |A| + |B| - |A \cup B|$
- Concatenation -

- Cartesian Product -  $|A \times B| = |A| \cdot |B|$
- Powerset -  $P(A) = 2^{|A|}$
- Total function -  $|A \rightarrow B| = |B|^{|A|}$
- Partial function -  $|A \rightharpoonup B| = |B + 1|^{|A|}$
- Multi function -  $|A \multimap B| = |P(B)|^{|A|}$

## 2 Alphabets and Languages

**Alphabet**  $A$  is a set of single characters. **Word**  $w$  in alphabet  $A$  is a string of characters from  $A$ .  $A^*$  denotes the set of all words in alphabet  $A$ . This operation is called *Kleene star*.

### 2.1 Regular Languages

A language is **regular** if it can be defined as follows.

- The empty language  $\emptyset$  is regular
- For each  $a \in \epsilon$  where  $\epsilon$  is the alphabet in question, the singleton language  $\{a\}$  is regular
- If  $A$  and  $B$  are regular, then  $A \cup B$  is regular
- If  $A$  and  $B$  are regular, then  $A \cdot B$  is regular
- If  $A$  is regular, then  $A^*$  is regular.

Regular languages can be expressed by **regular expressions**. For each rule above, we have a corresponding regular expression.

- $e \rightarrow \emptyset$
- $a \rightarrow \{a\}$
- $a \vee b \rightarrow \{a\} \cup \{b\} = \{a, b\}$
- $ab \rightarrow \{a\} \cdot \{b\} = \{ab\}$
- $a^* \rightarrow \{a\}^* = \{\emptyset, a, a^2, a^3, \dots\}$

### 2.2 Pumping Lemma

All finite languages are regular. For infinite languages, we use the Pumping lemma to determine whether or not they are regular. Specifically, the Pumping lemma states that if  $L$  is regular and infinite, then there exists an  $N > 0$  such that for all words  $\{w : w \in L \wedge |w| \geq N\}$ ,  $w$  can be decomposed into  $w = xyz$  such that

- $y \neq e$
- $|xy| \leq N$
- $wy^kz \in L$  for all  $k \geq 0$

This is typically used to show something is *not* regular through proof by contradiction.<sup>1</sup>

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<sup>1</sup>See worksheet problem 5 for an example.

## 3 Finite Automata

A **deterministic finite automata** consists of

- alphabet  $A$
- set of states  $S = S_0, S_1, S_2 \cdots S_n$
- transition function  $f : S \times A \rightarrow S$

The transition function given a state  $S_i$  and input  $a$  changes the machine to state  $S_j$ . By adding

- start state  $S_0$
- final states  $F \subset S$

we create a f.a. that can be said to *accept a language*. The language accepted by a fa can be defined as  $LM = \{w | S_0 \xrightarrow{w} \text{stops in final state}\}$

### 3.1 Deterministic versus Non-Deterministic

A deterministic fa is one where the number of arrows leaving each state is exactly equal to  $|A|$ . In a **non-deterministic** fa the number of arrows leaving a state is  $\geq 0$ . Additionally, the same word may lead to different states from the same state. In terms of functions, we consider the transition function as  $f : S \times A \rightarrow S$  for a dfa and  $f : S \times A \rightarrow \mathcal{P}(S)$  for a ndfa.

We describe a procedure for converting a ndfa into a dfa.<sup>2</sup>

1. List  $ES_0$ , the set including the initial state and all states reachable without consuming any characters
2. For  $a \in A$ , list the set of states reachable from  $ES_0$  by consuming  $a$
3. Repeat 2. until no new sets are produced (including  $\emptyset$ )
4. Draw the new dfa by treating each unique set as a state, where  $ES_0$  is the initial state and any set containing a final state from the ndfa is a final state

### 3.2 Regular Expressions to fa

The collection of languages accepted by fa is exactly equal to the regular languages. We can therefore define a simple fa for each of the regular expressions in section 2.1.

- $e$
- $a$
- $a \vee b$
- $ab$
- $a^*$

These can be used as building blocks to construct fa for more complex regular expressions. When two of these are chained together, a new state must be added where all final states can reach this new state using  $e$ .

### 3.3 Simplifying

### 3.4 Finding Language of fa

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<sup>2</sup>See worksheet problem 7 for example.