

Network Access Layer

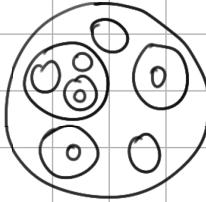
Hop 2 hop

Link layer: Responsible for transferring datagrams from one node to physically adjacent node.

layer 2

Link layer

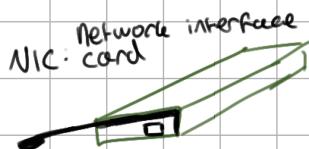
→ Framing (^{Encapsulates} datagram _{into a frame})



Router/host
↑

Definitions:

Flow Control: Pacing between adjacent sending and receiving nodes



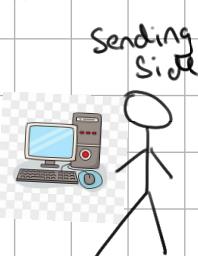
Error Detection: noise + errors caused by signal attenuation

↳ Receiver detects errors + signals transmission + drops frame

Error Correction: Receiver corrects bit errors w/o transmission.

Half-duplex and full-duplex: nodes can transmit but not at same time

Interface Communicating

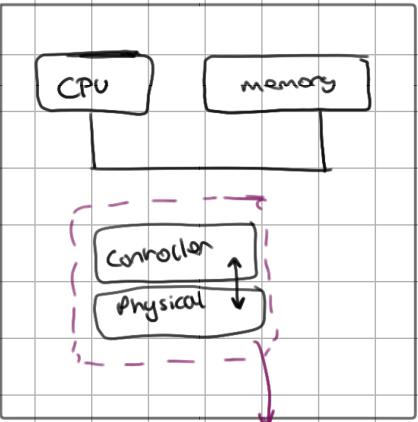


Receiving End



- ① puts datagram into a frame

- ② Adds error checking bits and reliable data transfer + flow control



① Looks for errors, reliable data, flow control

② Extracts datagram and sends everything to the upper layer.

ECE: Error Detection and correction bits.

Error Detection

= not 100% reliable !!

Datagram
↓

Sensitive Link

Everything in order ??

no →

Datagram

Addressing Scheme — MAC addresses.

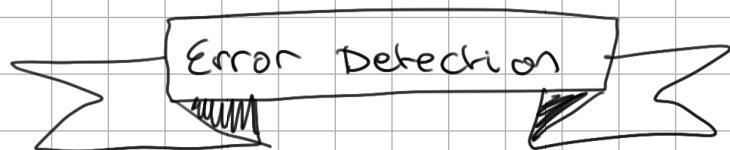
↳ 12 hex digits

Data Link → switches
→ NIC

Every NC has its own unique MAC address

Link Layer Services:

- **Framing**: wraps up data w/ header + trailer
 - **Link Access**: determines when node can send data
 - **MAC addressing**: Frame uses this to identify sender/receiver LAN.
 - **Reliable Delivery**
 - **Flow Control**: Prevents sender from overwhelming receiver
 - **Error Detection**: identifies if bits are corrupted during transmission
 - **Error Correction**: Fixes errors w/o retransmission
 - **Duplex Modes**: Half or Full communication
(one at a time) (both at once)



- Corrects 1 bit error
- Parity Bit : Adds 1 bit to detect 1 bit error
- 2D Parity : uses grid w/ rows & columns → checksum : Adds numbers in a row & column → TCP / UDP Protocols
- CRC : Polynomial Division to detect b bits of error
(Cyclic Redundancy Check)

Multiple Access protocols : when multiple devices share the same channel.

channel divides
Radio access channel into a set of parallel channels (TDMA, FDMA, CDMA)

Partitioning: chrm
Random devices send

Access whenever

turns: Devices take turns (e.g. token passing)

Random Access Protocols

Pure ALOHA : Sends anytime

slotted ALOHA : sends only at time slots

CSMA : sense channel before transmitting
If busy = waits

CSMA/CD : collision detection. (Ethernet)
Arbiter or collision

CSMA/CA : Collision (with) avoidance. Avoid collisions

Wands

however unused slots = waste bandwidths

FDMA → Each cell a frequency band

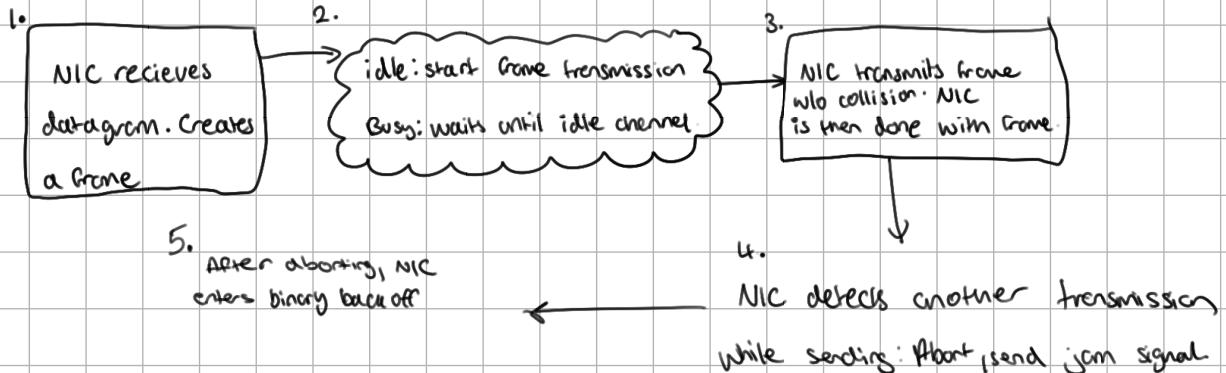
F DMA \rightarrow each gen a Frequency band.

Good isolation but can waste unused

www.english-test.net

bonds

Ethernet: CSMA/CD ~~Priority~~ Algorithm



$$\text{efficiency} = \frac{1}{1 + 5t_{\text{prop}}/t_{\text{trans}}} \quad \begin{matrix} \text{Time between to transmit max-size frame} \\ \text{max prop delay between 2 nodes in LAN.} \end{matrix}$$

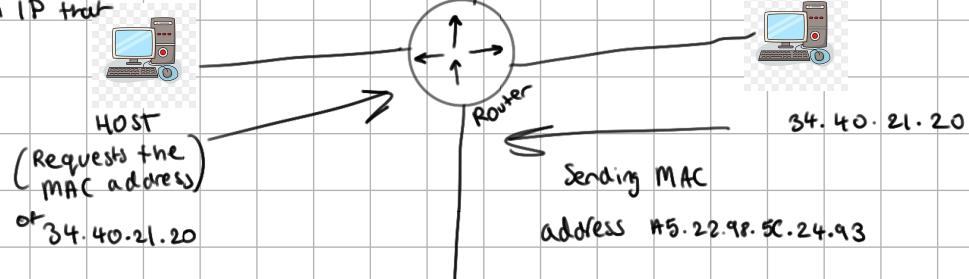
Switches Vs Routers

	Switch	Router
Layer addressing tables.	Data Link MAC address	Network IP address
	Forwarding	Calculated Routing Algorithm.

LAN = Local Area Network.

When a new PC joins, it needs an IP.

So it asks ARP to find an IP that matches the MAC address.



ARP vs DHCP

- DHCP prevents two IP addresses from being the same.

- Client connects:
 - uses DHCP to get IP
 - DHCP uses IP/Ethernet/UDP
- DNS Query:
 - needs MAC of router → uses ARP
 - sends DNS query via router to DNS server
- HTTP requests:

→ Opens TCP connection to web server

→ Send HTTP Get over TCP

→ Web server responds w/ HTTP

Routing to another Subnet: Device A wants to talk to B

Device A



Sends frame.

Gives core information to find B



default gateway

Device B



sent!

