# Lab 5 - Karatsuba polynomial multiplication

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#### I. ASSIGNMENT TASK

Perform the multiplication of 2 polynomials. Use both the regular  $O(n^2)$  algorithm and the Karatsuba algorithm, and each in both the sequential form and the parallelized form. Compare the 4 variants.

### II. IMPLEMENTATIONS

The sequential regular multiplication algorithm is the trivial, normal algorithm for multiplying numbers and polynomials. The parallelized version takes chunks from the polynomials, passes them to threads inside a custom implemented thread pool and multiplies them. (Complexity:  $O(n^2)$ )

For the Karatsuba algorithm, we split each polynomial in half recursively and the number of multiplications is reduced by 1. In a classic example of normal multiplication, 4 multiplications are required, while Karatsuba's algorithm only requires 3 multiplications. The parallelized version uses C++ futures with the std :: async function which runs under the std :: async launch strategy to spawn new threads for each of the 3 recursive calls to the Karatsuba function. (Complexity:  $O(n^{\log_2 3})$ )

#### III. RESULTS

The most obvious and noticeable detail of the results is that Karatsuba's algorithm only starts to become significantly faster when the polynomial degree is very large. The sequential versions of both the  $O(n^2)$  and  $O(n^{\log_2 3})$  algorithms have the lowest performance, whereas the parallel versions seem to be the leading methods in terms of execution time.

Another possible reason why the Karatsuba algorithm seemed to underperform when running tests on polynomials with small degrees is that the variables used for the polynomials inside the function doing the computation are copied in recursive calls, but larger polynomial sizes show a clear advantage of Karatsuba's algorithm over the traditional methods.

Unfortunately, the Karatsuba implementation works only for polynomials of degree  $2^n - 1$ 

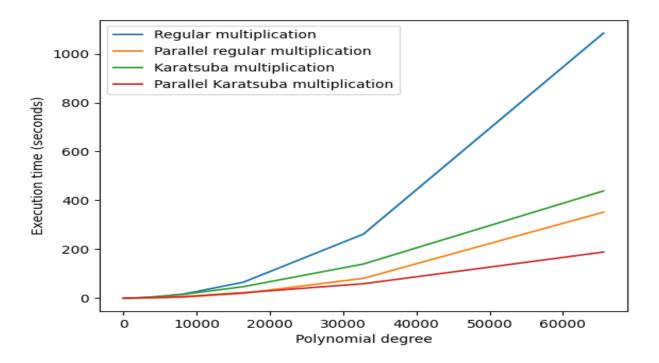


Fig. 1: Line plot containing the polynomial degree on the X-axis and the execution time in seconds on the Y-axis

Polynomial degree	$O(n^2)$ seq.	$O(n^2)$ par.	$O(n^{\log_2 3})$ seq.	$O(n^{\log_2 3})$ par.
31	0.00012	0.00137	0.00221	0.00267
63	0.00046	0.00169	0.00746	0.00546
127	0.00242	0.00259	0.02091	0.02643
255	0.01338	0.00154	0.06089	0.02689
511	0.05748	0.00176	0.18336	0.07137
1023	0.24838	0.5049	0.57698	0.21373
2047	1.01859	0.5021	1.70614	0.69265
4095	4.06426	1.0635	5.12932	1.94257
8191	16.4522	4.1569	15.497	5.88652
16383	65.1912	19.7005	46.8556	22.0241
32767	261.86	80.9225	139.531	58.7718
65535	1064.32	351.596	438.625	188.38

TABLE I: Table with execution times (seconds)