

Graphing the Mermin-Peres Magic Square

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SUnMaRC
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Department of Mathematics and Statistics

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4. Define Graphs and Coloring

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- built a graph based on the game algebra of the Mermin-Peres Magic Square Game
- We chose the Mermin-Peres game because it is the smallest synchronous game that has a winning quantum strategy but no classical strategy
- encoded the graph into SageMath and ran a multitude of constructions to identify which vertices could be removed without changing the colorabilities of the graph

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- i.e. $P_A(a|x) = \sum_{b \in \mathcal{O}} P(a, b|x, y)$ is independent of y , and

$$P_B(b|y) = \sum_{a \in \mathcal{O}} P(a, b|x, y) \text{ is independent of } x.$$

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 - i.e. If they give different answers they lose

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Definition

The Mermin-Peres Magic Square Game is a synchronous non-local game. In this game, Alice and Bob will each receive one equation from the following list:

$$x_1 + x_2 + x_3 \equiv 0 \pmod{2}$$

$$x_4 + x_5 + x_6 \equiv 0 \pmod{2}$$

$$x_7 + x_8 + x_9 \equiv 0 \pmod{2}$$

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Alice and Bob fill out their assigned equations with 0's and 1's simultaneously without knowing how the other fills theirs out. Their goal is to satisfy each of their own equations while producing the same value for any shared variable.

Magic Square Representation

The Magic Square game can be represented by a 3x3 grid by labeling each cell like so:

x_1	x_2	x_3
x_4	x_5	x_6
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Magic Square Walkthrough

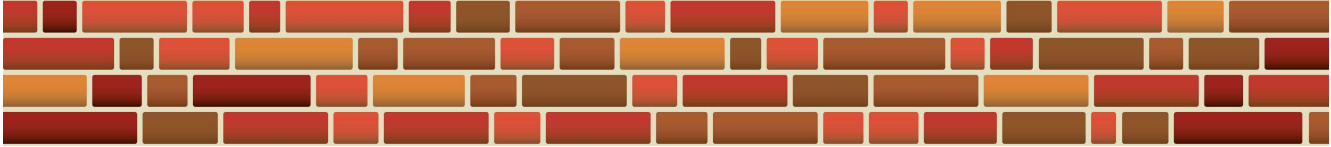
Alice

x ₁	x ₂	x ₃
x ₄	x ₅	x ₆
x ₇	x ₈	x ₉

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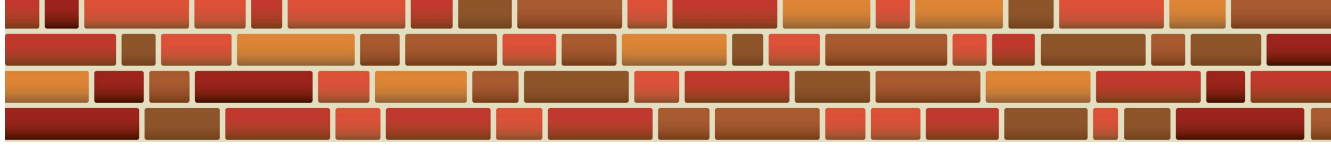
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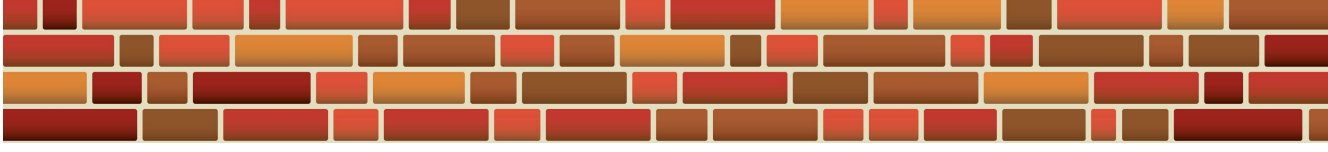
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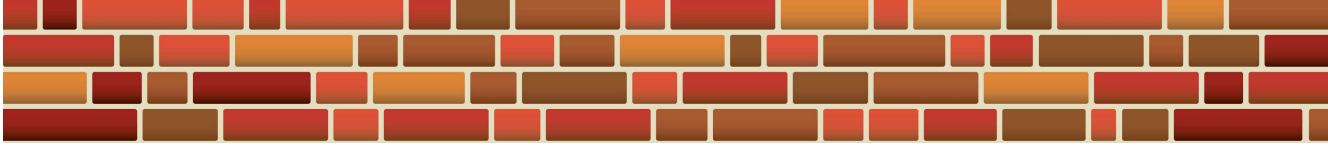
$$0 = x_4 + x_5 + x_6$$

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$$x_3 + x_6 + x_9 = 1$$

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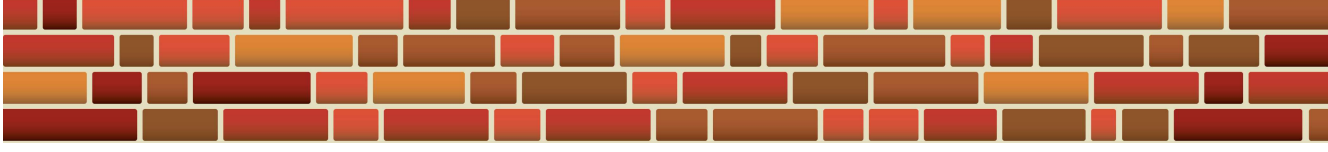
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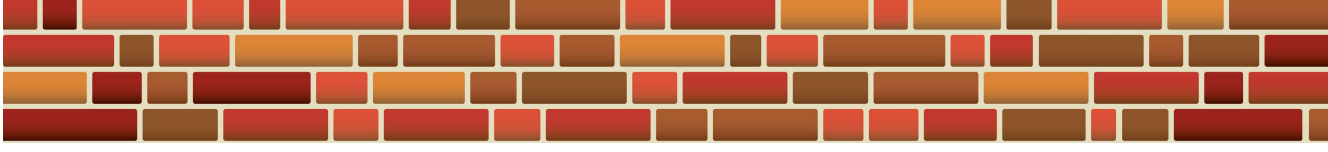
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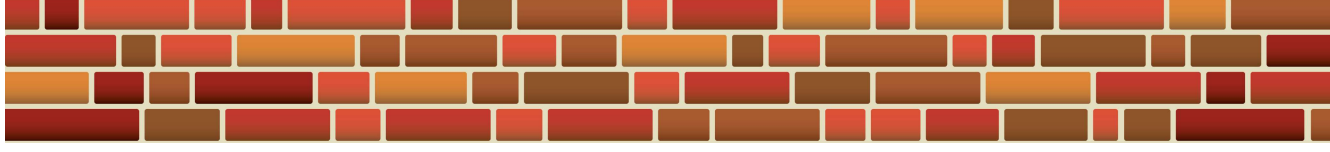
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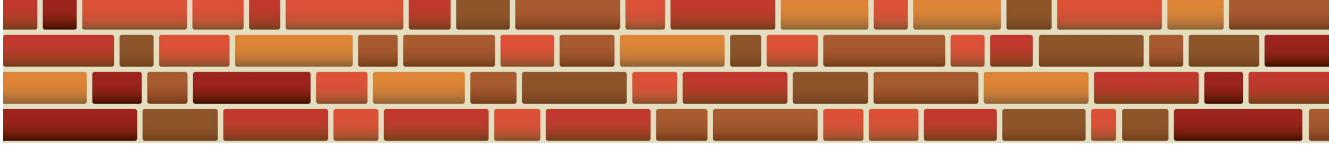
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Definition

A classical strategy for a Non-Local Game is represented by a function $f : \mathcal{I} \rightarrow \mathcal{O}$.

For the Magic Square Game, we can define a classical strategy using the grid below. The best strategy has them agree upon a set table of answers, this allows them to win most of the time. It is impossible to create a solution set that satisfies all nine equations at the same time, otherwise they would be able to win every single game.

x_1 1	x_2 0	x_3 1	$1 + 1 = 0$
x_4 1	x_5 1	x_6 0	$1 + 1 = 0$
x_7 0	x_8 1	x_9 0	$1 \neq 0$

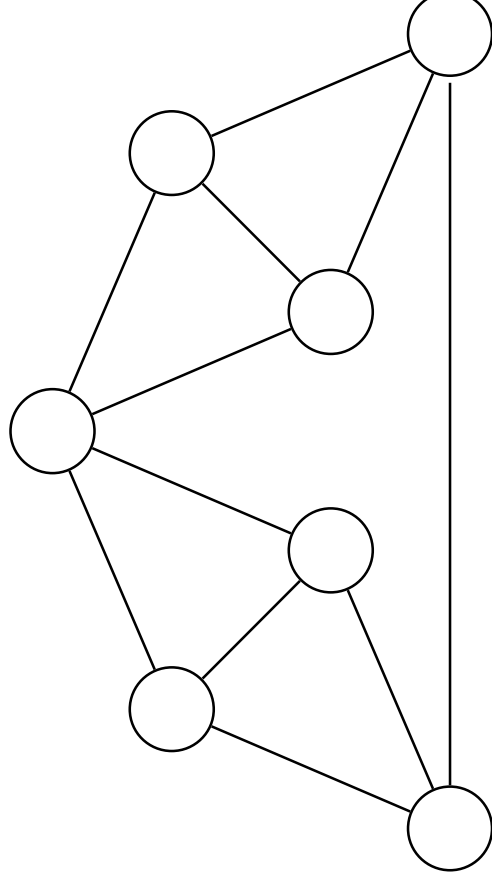
$$\begin{array}{ccccccc} 1 & 1 & 1 & & & & \\ + & + & + & & & & \\ 1 & 1 & 0 & & & & \\ = & = & = & & & & \\ 0 & 0 & 1 & & & & \end{array}$$

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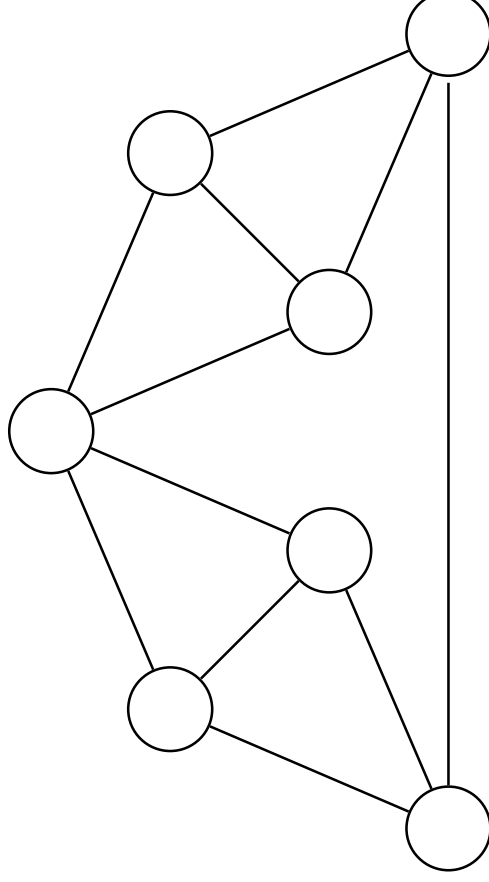
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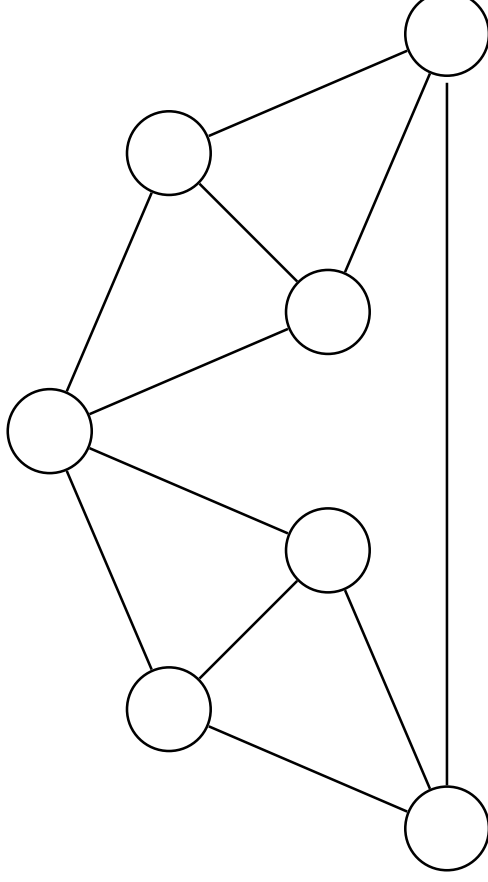
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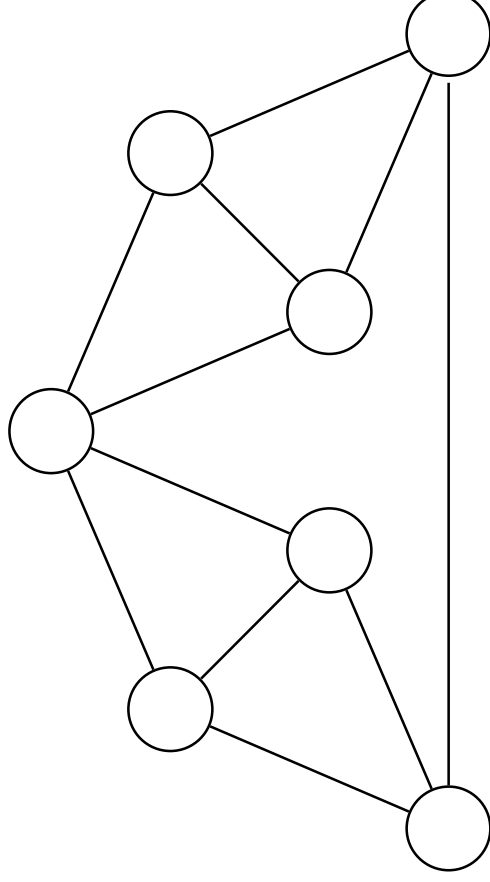
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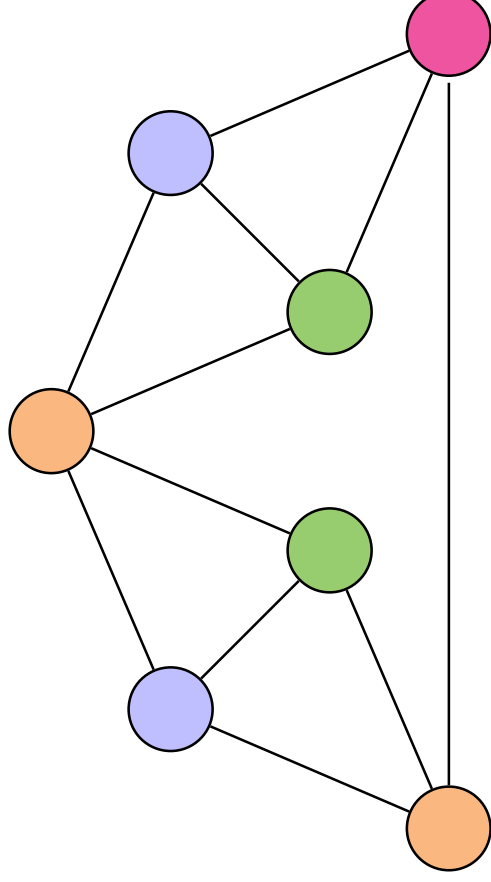
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- We say that a graph is k -colorable if the minimum amount of colors needed to color the graph is at most k .



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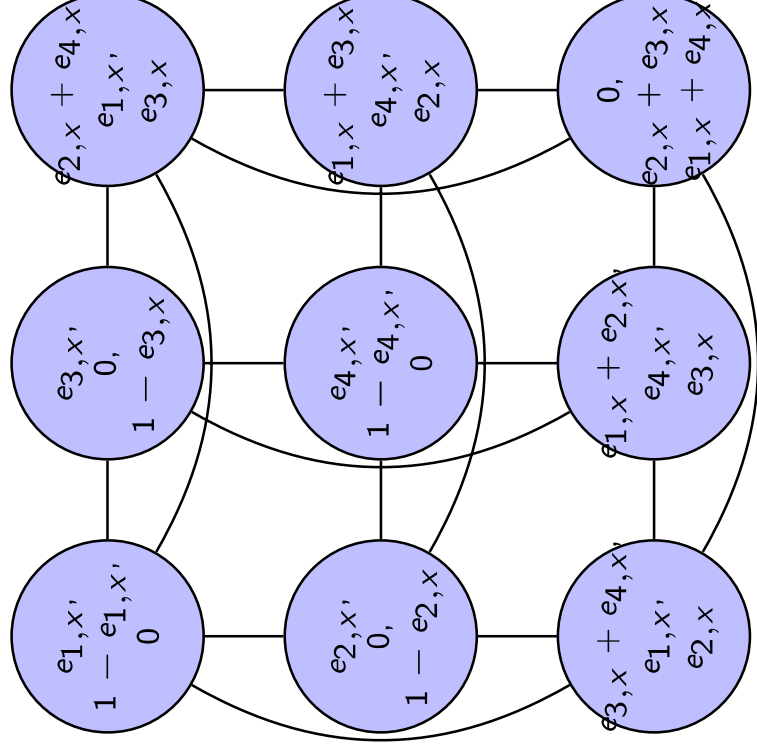
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 - If $1 \leq x, y \leq 6$ and a is a solution to equation x , b is a solution to equation y , and a, b are inconsistent, then $e_{a,x} \cdot e_{b,y} = 0$.

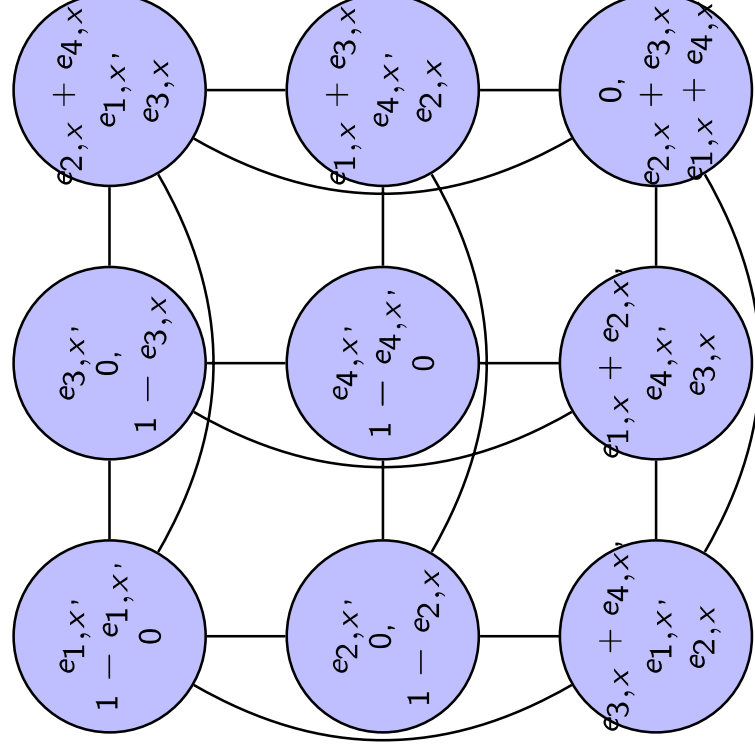
Building a Graph from Mermin-Peres

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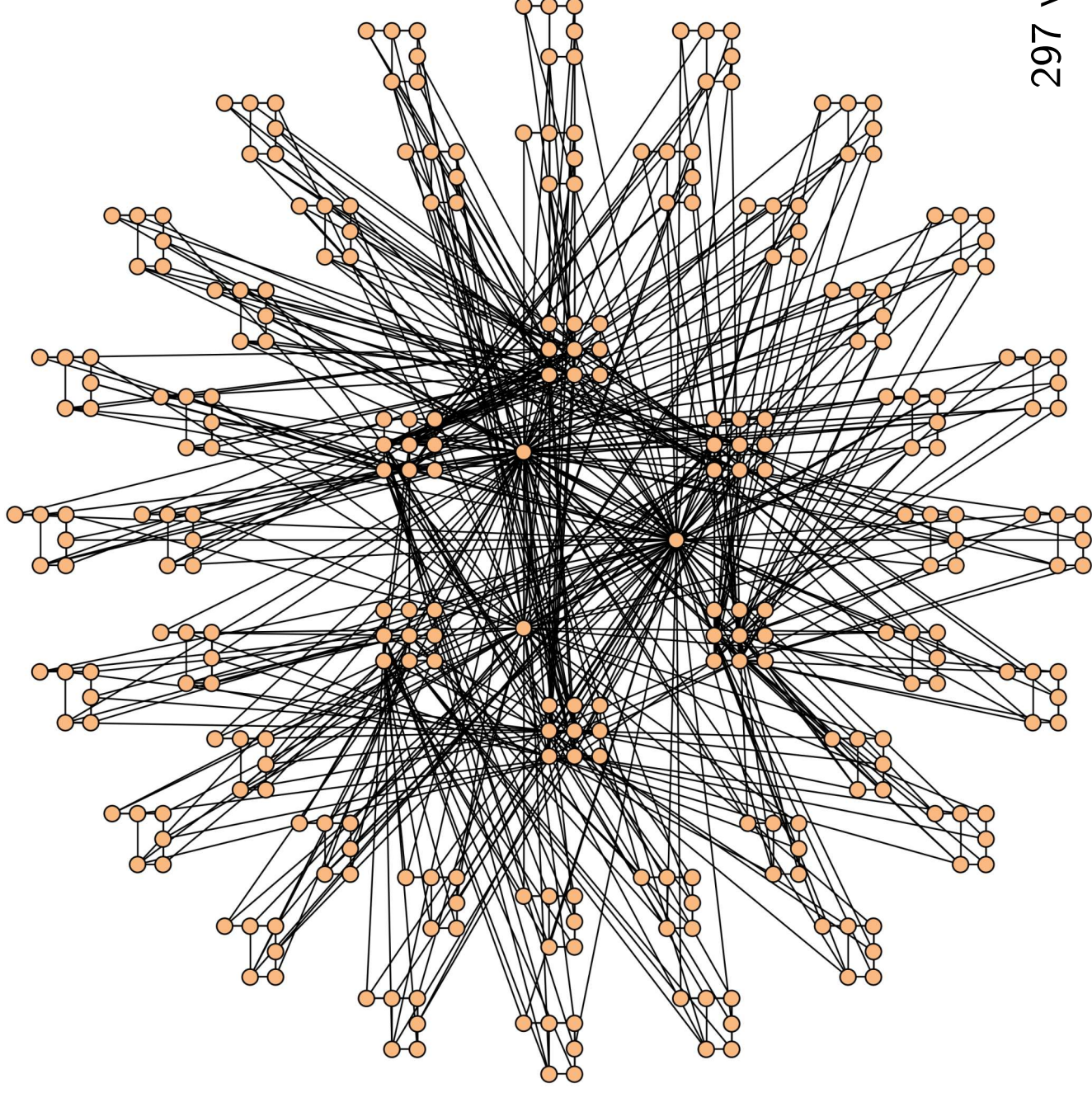
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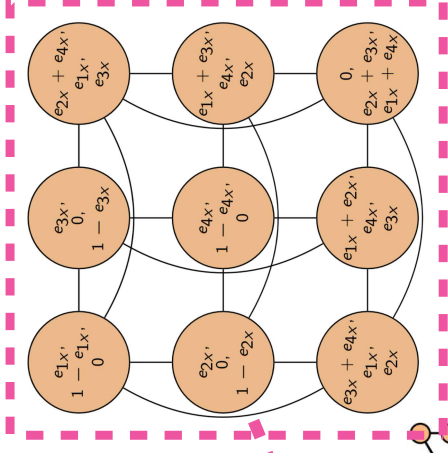
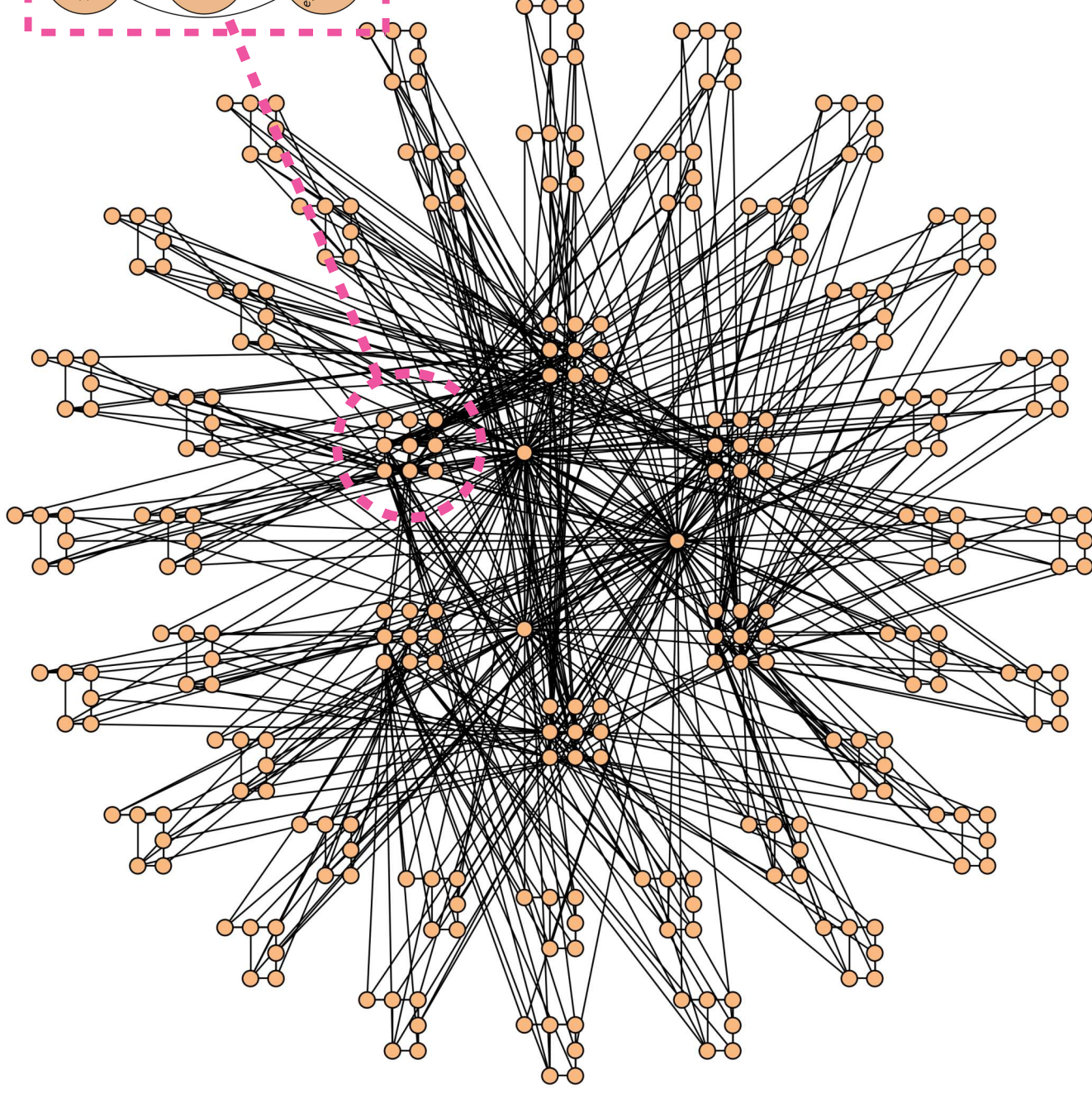
In quantum coloring, each vertex is colored with 3 projections that sum to 1 instead of a single color. These act on a shared entangled state, allowing Alice and Bob to coordinate their responses even when they are separated.

Mermin-Peres Graph



297 vertices

Mermin-Peres Graph



Thank you!

Questions?