

CSCI 305, Homework # 7

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Due date: Midnight, Tuesday, May 29

Quadratic probing. This is problem 11-3 in the book.

Suppose that we are given a key k to search for in a hash table with positions $0, 1, \dots, m-1$, and suppose that we have a hash function h mapping the key space into the set $\{0, 1, \dots, m-1\}$. The search scheme is as follows:

1. Compute the value $j = h(k)$ and set $i = 0$.
2. Probe in position j for the desired key k . If you find it, or if this position is empty, terminate the search.
3. Set $i = i + 1$. If i now equals m , the table is full, so terminate the search. Otherwise, set $j = (i + j) \bmod m$ and return to step 2.

Assume that m is a power of 2.

- a. Show that this scheme is an instance of the general “quadratic probing” scheme by exhibiting the appropriate constants c_1 and c_2 for equation (11.5).

The equation above can be written as

$$\begin{aligned}h(k, 0) &= j \\h(k, 1) &= j + 1 \bmod m \\h(k, 2) &= j + 1 + 2 \bmod m \\h(k, 3) &= j + 1 + 2 + 3 \bmod m \\h(k, i) &= j + \sum_{q=1}^i q \bmod m \\h(k, i) &= j + \frac{i(i+1)}{2} \bmod m \\h(k, i) &= j + \frac{i^2}{2} + \frac{i}{2} \bmod m\end{aligned}$$

The general form of the equation is

$$h(k, i) = (h(k) + c_1 i + c_2 i^2) \bmod m$$

Choose $c_1 = 1/2$ and $c_2 = 1/2$

$$h(k, i) = (h(k) + i/2 + i^2/2) \mod m$$

$$h(k, i) = j + \frac{i}{2} + \frac{i^2}{2} \mod m$$

b. Prove that this algorithm examines every table position in the worst case.

There are 3 ways to end a search. 1. We find the key. 2. We find an empty entry. 3. $i = m$. In the worst case we don't find the key and we have a full table. So case 1 and 2 are never met. Leaving only case 3. Case 3 guarantees that we will not look at more than m entries because we stop when $i = m$.

Now to prove this claim we must show that over the course of those m tries we don't look at the same item twice. In order to visit the same node twice we would have to get two unique numbers such that $h(k, i) = h(k, q)$ where both i and q are less than $m-1$ and greater than 0.

$$j + \frac{i}{2} + \frac{i^2}{2} \mod m = j + \frac{q}{2} + \frac{q^2}{2} \mod m$$

$$\frac{i}{2} + \frac{i^2}{2} \mod m = \frac{q}{2} + \frac{q^2}{2} \mod m$$

$$\frac{i}{2} + \frac{i^2}{2} - (\frac{q}{2} + \frac{q^2}{2}) \mod m = 0$$

$$\frac{(i + i^2) - (q + q^2)}{2} \mod m = 0$$

$$\frac{i(1 + i) - q(1 + q)}{2} \mod m = 0$$

$$\frac{(i - q)(i + q + 1)}{2} \mod m = 0$$

To get mod m to equal 0, $(i - q)$ must equal m , $(i + q + 1)$ must equal m or $(i - q)(i + q + 1)$ must equal m .

Both i and q are less than m , so $i - q$ is definitely not m .

$(i + q + 1)$ could equal m , but if it does then $(i - q)$ is forced to be odd. Because of the divide by two can't be taken out of the odd, it will force $(i + q + 1)$ to be equal to $m/2$, not m .

$(i + q + 1)$ can't equal $2m$ because of i and q are less than $m - 1$.

$(i - q)(i + q + 1)$ can't equal m because one will be odd and the other even, so can't make a power of two.