CSCI 305, Homework # 3

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Due date: Tuesday, May 8, midnight.

- Find $\Theta(T(n))$ for each case.
- For problem 1, find the solution three ways:
 - The master theorem
 - The substitution method
 - Iterate, cancel, and solve the summation to find an exact solution.
- For all other problems, use just the master theorem.
- When using the master theorem, be explicit about $n^{\log_b a}$, and about which case of the master theorem applies, about making sure the special condition applies in case 3 of the master theorem, and about proving f's correct asymptotic relationship to $n^{\log_b a}$.
- 1. Solve this one three different ways: master theorem, substitution method, and the iterate/cancel/summation method.

$$T(n) = 2T(n/2) + n^4$$

Master theorem

$$f(n) \quad vs. \quad n^{log_ba}$$

$$n^4 \quad vs. \quad n^{log_22}$$

$$n^4 \quad vs. \quad n^1$$

$$f(n) = \Omega(n^{log_ba+\epsilon})$$

$$T(n) = \Theta(n^4)$$

Substitution method Guess $T(n) \le c(n^4)$

$$T(n) = 2T(n/2) + n^4$$

 $T(n) \le 2(cn/2)^4 + n^4$
 $T(n) \le c(n^4) + n^4$
 $T(n) < c(n^4)$

We can use this same proof for $T(n) = \Omega(n^4)$ by flipping the signs, So $T(n) = \Theta(n^4)$.

1

Iterate/cancel/summation method

$$T(n) = 2T(n/2) + n^4$$

$$2T(n/2) = 2^2T(n/2^2) + 2((n/2)^4)$$

$$2^2T(n/2^2) = 2^3T(n/2^3) + 2^2((n/2^2)^4)$$

$$2^{k-1}T(n/2^{k-1}) = 2^kT(n/2^k) + 2^{k-1}((n/2^{k-1})^4)$$

$$T(n) = 2^{\log_2 n}T(n/2^{\log_2 n}) + 2^{(\log_2 n)-1}((n/2^{(\log_2 n)-1})^4)$$

$$T(n) = nT(n/n) + \sum_{i=0}^{(\log_2 n)-1} 2^i(n/2^i)^4$$

$$T(n) = nT(1) + n^4 \sum_{i=0}^{(\log_2 n)-1} 2^i/16^i$$

$$T(n) = \Theta(n^4)$$

2. Solve this one using the master theorem.

$$T(n) = 16T(n/4) + n^2$$

3. Solve this one using the master theorem.

$$T(n) = 2T(n/4) + \sqrt{n}$$

4. Solve this one using the master theorem.

$$T(n) = 4T(n/3) + n \lg n$$

5. Solve this one using the master theorem.

$$T(n) = 4T(n/2) + n^2\sqrt{n}$$