

CSCI 305, Homework # 2

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Due date: Tue, May 1, midnight

In all cases, we require that $f(n)$ and $g(n)$ be positive functions, *i.e.* $f(n) > 0$ and $g(n) > 0$ for all $n > 0$. Prove or disprove each of the following conjectures.

1. $f(n) = O((f(n))^2)$
 $f(n) \leq c(f(n))^2$
let $c = 1$
 $f(n) \leq f(n) * f(n)$

Provided the result of $f(n)$ is equal to or greater than 1 this is always true

2. $f(n) = \Theta(f(n/2))$
take $f(n) = 2^n$
 $2^n = \Theta(2^{n/2})$
for $n = 2$ $f(n) = 4$ and $f(n/2) = 2$ difference = 2
for $n = 4$ $f(n) = 16$ and $f(n/2) = 4$ difference = 12
for $n = 6$ $f(n) = 64$ and $f(n/2) = 8$ difference = 46
the difference between the two functions is increasing meaning as $n \rightarrow \infty$ the difference between goes to infinity and $f(n)$ is $o(n/2)$ for this equation.
3. $f(n) + o(f(n)) = \Theta(f(n))$
 $cg(n) \leq f(n) \leq dg(n)$ for all $n \geq n_0$
 $f(n) + o(f(n)) \leq df(n)$
 $o(f(n))$ is always smaller than $f(n)$, so we can substitute that in.
 $f(n) + f(n) \leq df(n)$
 d can be 2

 $cf(n) \leq f(n) + o(f(n))$
let $c = 1$
 $f(n) \leq f(n) + o(f(n))$
 $f(n)$ will always be less than $f(n) + \text{something}$.
4. If $f(n) = O(g(n))$ then $f(n) + g(n) = O(f(n))$.
Suppose $f(n) = O(n)$ and $g(n) = O(n^2)$
 $f(n) = O(g(n))$ would be true,
but $f(n) + g(n) = O(f(n))$ would not be true because $g(n) < f(n)$.