

But, by definition, $\text{height}((\varphi_1 \vee \varphi_2)) = 1 + \max(\text{height}(\varphi_1), \text{height}(\varphi_2))$ and, as $\max(a, b) \leq a + b$ (for any naturals $a, b \in \mathbb{N}$), we have that $\text{height}((\varphi_1 \vee \varphi_2)) \leq 1 + \text{height}(\varphi_1) + \text{height}(\varphi_2)$. But $\text{height}(\varphi_i) \leq \text{size}(\varphi_i)$ ($1 \leq i \leq 2$) by our hypothesis, and therefore $\text{height}((\varphi_1 \vee \varphi_2)) \leq 1 + \text{size}(\varphi_1) + \text{size}(\varphi_2)$. But, by definition, $\text{size}((\varphi_1 \vee \varphi_2)) = 1 + \text{size}(\varphi_1) + \text{size}(\varphi_2)$, and therefore $\text{height}((\varphi_1 \vee \varphi_2)) \leq \text{size}((\varphi_1 \vee \varphi_2))$, what we had to prove.

4. (Inductive Case iii) Similar to Inductive Case ii.

q.e.d.

4.4 Exercise Sheet

Exercise 19. Compute, using the function *subf*, the set of subformulae of the following formulae:

1. $((p \wedge \neg q) \wedge r)$; 2. $((p \vee \neg q) \wedge r)$; 3. $\neg((p \vee \neg q) \wedge r)$.

Exercise 20. Compute the abstract syntax trees of the following formulae:

1. $((p \wedge \neg q) \wedge r)$;
2. $((p \vee \neg q) \wedge r)$;
3. $\neg((p \vee \neg q) \wedge r)$;
4. $(\neg(p \vee \neg q) \wedge r)$.

Exercise 21. Recall the recursive definition of the function $\text{height} : \mathbb{PL} \rightarrow \mathbb{N}$, which computes, given a formula, the height of its abstract syntax tree. Compute the height of the formulae shown in Exercise 20.

Exercise 22. Recall the recursive definition of the function $\text{size} : \mathbb{PL} \rightarrow \mathbb{N}$, which computes the number of nodes in abstract syntax tree of a formula. Compute the size of the formulae shown in Exercise 20.

Exercise 23. Recall the recursive definition of the function $\text{prop} : \mathbb{PL} \rightarrow 2^A$, which computes, given a formula, the set of propositional variables occurring in the formula. Compute the set of propositional variables occurring in the formulae shown in Exercise 20.

Exercise 24. Show by structural induction that $\text{height}(\varphi) < \text{size}(\varphi) + 1$ for any formula $\varphi \in \mathbb{PL}$.