

Logic for Computer Science 2020-2021 *Alexandru Ioan Cuza* University

But, by definition, $height((\varphi_1 \vee \varphi_2)) = 1 + \max(height(\varphi_1), height(\varphi_2))$ and, as $\max(a, b) \leq a + b$ (for any naturals $a, b \in \mathbb{N}$), we have that $height((\varphi_1 \vee \varphi_2)) \leq 1 + height(\varphi_1) + height(\varphi_2)$. But $height(\varphi_i) \leq size(\varphi_i)$ ($1 \leq i \leq 2$) by our hypothesis, and therefore $height((\varphi_1 \vee \varphi_2)) \leq 1 + size(\varphi_1) + size(\varphi_2)$. But, by definition, $size((\varphi_1 \vee \varphi_2)) = 1 + size(\varphi_1) + size(\varphi_2)$, and therefore $height((\varphi_1 \vee \varphi_2)) \leq size((\varphi_1 \vee \varphi_2))$, what we had to prove.

4. (Inductive Case iii) Similar to Inductive Case ii.

q.e.d.

4.4 Exercise Sheet

Exercise 19. *Compute, using the function `subf`, the set of subformulae of the following formulae:*

1. $((p \wedge \neg q) \wedge r)$;
2. $((p \vee \neg q) \wedge r)$;
3. $\neg((p \vee \neg q) \wedge r)$.

Exercise 20. *Compute the abstract syntax trees of the following formulae:*

1. $((p \wedge \neg q) \wedge r);$

2. $((p \vee \neg q) \wedge r);$

3. $\neg((p \vee \neg q) \wedge r);$

4. $(\neg(p \vee \neg q) \wedge r).$

Exercise 21. Recall the recursive definition of the function $\text{height} : \text{PL} \rightarrow \mathbb{N}$, which computes, given a formula, the height of its abstract syntax tree. Compute the height of the formulae shown in Exercise 20.

Exercise 22. Recall the recursive definition of the function $\text{size} : \text{PL} \rightarrow \mathbb{N}$, which computes the number of nodes in abstract syntax tree of a formula. Compute the size of the formulae shown in Exercise 20.

Exercise 23. Recall the recursive definition of the function $\text{prop} : \text{PL} \rightarrow 2^A$, which computes, given a formula, the set of propositional variables occurring in the formula. Compute the set of propositional variables occurring in the formulae shown in Exercise 20.

Exercise 24. *Show by structural induction that $\text{height}(\varphi) < \text{size}(\varphi) + 1$ for any formula $\varphi \in \text{PL}$.*