# Problem Set 1 - ECON 880

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# Problem 1

(1a)

We can evaluate the polynomial as follow:

$$f(x,y) = 83521y^8 + 578x^2y^4 - 2x^4 + 2x^6 - x^8$$
  
=  $y^4(83521y^4 + 578x^2) + x^4(-2 + 2x^2 - x^4)$ 

(1b)

We estimate the order of magnitude (number of digits) of

$$83521y^8,$$
 (1)

where y = 2298912. We can rewrite the expression as follows:

$$83521y^8 = 83521 \cdot 2298912^8$$
  
=  $8.3521 \cdot 10^4 \cdot (2.298912 \cdot 10^6)^8$ 

Then, the digits can be estimated by using the following formula

$$\begin{array}{lll} 1 + \lfloor \log_{10}(83521 \cdot 2298912^8) \rfloor & = & 1 + \lfloor \log_{10}(83521 \cdot 2298912^8) \rfloor \\ & = & 1 + \lfloor \log_{10}(8.3521 \cdot 10^4 \cdot (2.298912 \cdot 10^6)^8) \rfloor \\ & = & 1 + \lfloor \log_{10}8.3521 + \log_{10}10^4 + \log_{10}2.298912^8 + \log_{10}(10^6)^8 \rfloor \\ & = & 56 \end{array}$$

Now, we compute the last digits of that term, first by plugging in only the last digit of y into expression (1): our calculation gives 21312256 as a result. However, printing out the complete term using the command fprintf((%.0f)n',  $83251 * y^8$ ), we obtain

2131225664948367112155225808364994454077172753899719772317155328.

We see that both results have different last digit  $(6 \neq 8)$ .

(1c)

test 2

# Problem 2

In this exercise, we write an algorithm to determine the relative speeds of addition, multiplication, division, exponentiation, and the logarithmic function of our computer. The computer uses Intel(R) Core(TM) i5-1035G4 CPU @ 1.10GHz, 1.50 GHz with 64-bit operating system and 8.00 GB installed RAM. We proceed by generating two matrices A and B of size  $10^4 \times 10^4$  using the rand() function in Matlab. Then, we use them in element-wise operations of addition (A + B), multiplication  $(A \cdot *B)$ , division  $(A \cdot /B)$ , exponentiation  $(A \cdot ^B)$ , as well as the logarithmic function (log(A)). The statements tic and toc are used around them to measure the computation time. This procedure is iterated 100 times, and the average computation time for each operation is computed. The results are as follows:

- Average computation time for variable initiation is 1.54030 seconds
- Average computation time for addition is 0.12073 seconds
- Average computation time for multiplication is 0.12083 seconds
- Average computation time for division is 0.11823 seconds
- Average computation time for exponentiation is 3.70586 seconds
- Average computation time for log function is 0.91528 seconds

# Problem 3

In order to find our machine  $\varepsilon$ , we follow Ken Judd's definition<sup>†</sup> by writing a while loop to subtract (resp. add) progressively smaller numbers  $\epsilon$  from (resp. to) 1 until the condition  $1 + \epsilon > 1 > 1 - \epsilon$  is no longer satisfied. Repeat the exercise using 0.001 and 1000 instead of 1. We verify our results by comparing them with the epsilons delivered by the built-in Matlab function eps(x), and the exponents with log2(eps(x)). The results are shown on Table 1.

x	$\varepsilon(x)$ - decimal	$\varepsilon(x)$ - exponential
0.001	2.16840434497101e-19	$2^{-62}$
1	2.22044604925031e-16	$2^{-52}$
1000	1.13686837721616e-13	$2^{-43}$

Table 1: Machine  $\varepsilon$  for diverse values of x

Comment: As x increases, the machine  $\varepsilon(x)$  also does increase. This is to be expected, since  $\varepsilon(x) \approx x\varepsilon(1)$ . Thus,  $\varepsilon(1000) > \varepsilon(1) > \varepsilon(0.001)$ .

#### Problem 4

We wrote a loop to evaluate the convergence of the following sequences:

(4a) 
$$x_k = \sum_{k=1}^n \frac{1}{2^n}$$
, where  $\lim_{k \to \infty} x_k = 1$ 

(4b) 
$$y_k = \sum_{k=1}^n \frac{1}{n}$$
, where  $\lim_{k \to \infty} y_k = \infty$ 

 $<sup>^\</sup>dagger$ Kenneth L. Judd, 1998. "Numerical Methods in Economics," MIT Press Books, The MIT Press, p.30

We use absolute and relative convergence criteria, and tolerance distance  $\delta \in \{10^{-2}, 10^{-4}, 10^{-6}\}$  as stopping rule. We limit the maximum number of iterations to 100,000. The number of iterations before convergence, as well the final guess are reported on Table 2 and 3 for  $x_k$  and  $y_k$ , respectively.

δ	$10^{-2}$	$10^{-4}$	$10^{-6}$
no. of iteration - absolute criteria	7	14	20
final guess - absolute criteria	0.992187500000000	0.999938964843750	0.999999046325684
no. of iteration - relative criteria	7	14	20
final guess - relative criteria	0.992187500000000	0.999938964843750	0.999999046325684

Table 2: Convergence for  $x_k$ 

δ	$10^{-2}$	$10^{-4}$	$10^{-6}$
no. of iteration - absolute criteria	100	10,000	100,000
final guess - absolute criteria	5.18737751763962	9.78760603604435	12.0901461298633
no. of iteration - relative criteria	100	10,000	100,000
final guess - relative criteria	5.18737751763962	9.78760603604435	12.0901461298633

Table 3: Convergence for  $y_k$ 

Comment: Table 2 shows that both absolute and relative convergence criteria lead to the same number of iterations for the sequence  $x_k$ . As tolerance distance  $\delta$  lowers, the final guess for  $x_k$  approaches its true limit 1. Table 3 also shows that both absolute and relative criteria lead to the same number of iterations for the sequence  $y_k$ . Since  $y_k$  is a divergent sequence, the final guess will be higher (with no upper bound), the more we iterate. By lowering tolerance distance  $\delta$ , the algorithm needs to iterate longer in order to satisfy the stopping rule, which in turn results in higher final guess. Since there is no upper bound for  $y_k$ , we can always make the tolerance distance  $\delta$  even lower, and obtain even higher final guess for  $y_k$ .