Lesson 110: Discrete Cosine Transform



Discrete Cosine Transform

8-point 1D DCT

$$X(u) = \frac{C(u)}{2} \sum_{i=0}^{7} x(i) \cos \frac{(2i+1)u\pi}{16}$$

where x(i) are the inputs, X(u) are the outputs, and $C(u)=1/\sqrt{2}$ for u=0, otherwise is 1

8 × 8-point 2D DCT:

$$X(u,v) = \frac{C(u)C(v)}{2} \sum_{i=0}^{7} \sum_{j=0}^{7} x(i,j) \cos \frac{(2i+1)u\pi}{16} \cos \frac{(2j+1)v\pi}{16}$$



product Using matrix notation, the 8×8 -point 2D DCT can be expressed as a matrix

$$[X] = [C_{8 \times 8}] * [x]$$

where $\left[x\right]$ and $\left[X\right]$ are the two-dimensional input and output sequences, respectively

A standard strategy to compute the 2D DCT is the row-column separation

$$[X] = [C] * [x] * [C]^t$$

where $\left[C\right]$ is the matrix for 1D DCT



The expanded matrix representation for the 8-point 1D DCT

$$\begin{bmatrix} X(0) \\ X(1) \\ X(1) \\ X(2) \\ X(3) \\ X(5) \\ X(6) \\ X(7) \end{bmatrix} = \begin{bmatrix} C4 & C4 & C4 & C4 & C4 & C4 & C4 \\ C1 & C3 & C5 & C7 & -C7 & -C5 & -C3 & -C1 \\ C2 & C6 & -C6 & -C2 & -C2 & -C6 & C6 & C2 \\ C2 & C6 & -C6 & -C2 & -C2 & -C6 & C6 & C2 \\ C3 & -C7 & -C1 & -C5 & C5 & C1 & C7 & -C3 \\ C4 & -C4 & -C4 & C4 & -C4 & -C4 & -C4 & -C4 \\ X(5) \\ X(6) \\ X(7) \end{bmatrix} = \begin{bmatrix} C4 & C4 & C7 & C3 & -C3 & -C7 & C1 & -C5 \\ C6 & -C2 & C2 & -C6 & -C6 & C2 & -C2 & C6 \\ C3 & -C1 & C1 & -C3 & -C7 & C1 & -C7 \end{bmatrix} \begin{bmatrix} x(0) \\ x(1) \\ x(1) \\ x(2) \\ x(2) \\ x(2) \\ x(3) \\ x(4) \\ x(4) \\ x(5) \\ x(6) \\ x(6) \\ x(7) \end{bmatrix}$$

where $Cn = \cos n\pi/16$

A direct implementation of would require 64 multiplications and 56 additions



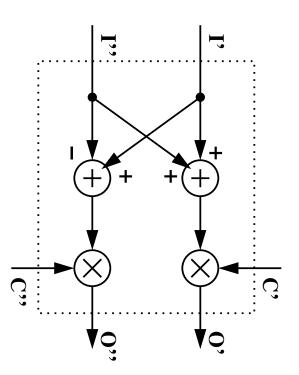
By exploiting the symmetry in the DCT coefficient matrix the number additions/subtractions of multiplications is reduced to 32, at the expense of additional 8

[X(7)]	X(5)	X(3)	X(1)	X(6)	X(4)	X(2)	[X(0)]
0	0	0	0	C6	C4	C2	C4
0	0	0	0	-C2	-C4	C6	C4
0	0	0	0	C2	-C4	-C6	C4
0	0	0	0	-C6	C4	-C2	C4
C7	C_5	C3	C1	0	0	0	0
-C5	-C1	-C7	C3	0	0	0	0
C3	C7	-C1	C5	0	0	0	0
-C1	C3	-C5	C7	0	0	0	0
*							
x(3) - x(4)	x(2) - x(5)			+	+	+	+



Butterfly computation is widely encountered:

$$O' = (I' + I'') \cdot C'$$
 $O'' = (I' - I'') \cdot C''$





- Butterfly vector-valued function
- Any restrictions due to ARM architecture?
- Additions and subtractions that's easy
- Multiplication by constant (seven constants)
- Software: use MUL operation
- Hardware: one/two full multipliers or seven multipliers-by-constant?
- How many new instructions?
- What syntax? where Rs1 = [I', C'], Rs2 = [I'', C''], Rt = [O1, O'']DCT Rs1, Rs2, Rt
- What latency?
- How to rewrite the code using these new instructions?



DCT - project requirements

- Build the testbench: the input is a 8-by-8 matrix of integers
- Write a DCT in software and highlight the Butterfly routine
- Implement Butterfly in:
- software (write C routines)
- horizontal firmware with two issue slots
- custom hardware (write VHDL/Verilog)
- Define a new instruction (or new instructions) that does Butterfly
- arguments and one result per instruction call) You must comply with the ARM architecture (you can have at most two



DCT - project requirements

- Rewrite the high-level code and instantiate the new instruction
- Use assembly inlining
- Estimate:
- the performance improvement of hardware-based solution versus softwarebased solution
- the performance improvement of a 2-issue slot firmware-based solution versus software-based solution
- Estimate the penalty in terms of number of gates for the hardware solution



Questions, feedbacks



