# High Performance Graph and Data Analytics Project

Alessandro La Conca - Pablo Giaccaglia

#### What we have done?

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## Acceleration of Personalized PageRank using GPUs

briefly!

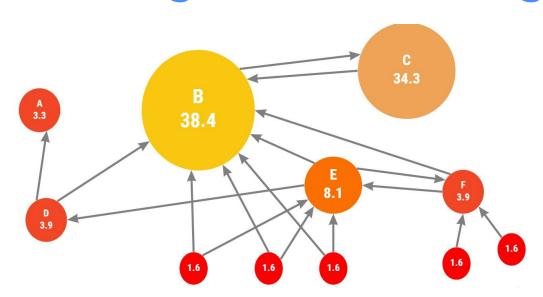
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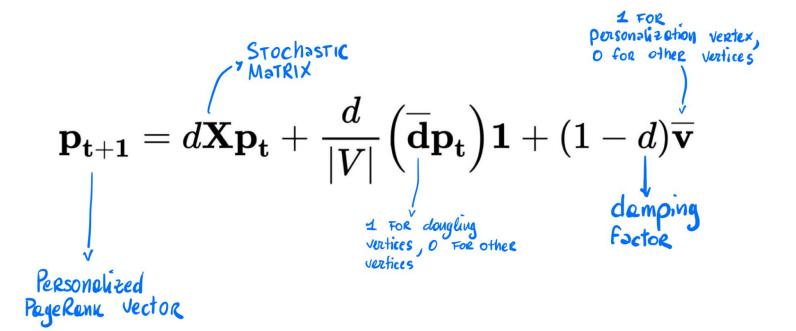
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- Given a web page, it gives a score to the others
- A ranking is generated, according to pages' relevance

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$$\mathbf{p_{t+1}} = d\mathbf{X}\mathbf{p_t} + rac{d}{|V|}\left(\overline{\mathbf{d}}\mathbf{p_t}
ight)\mathbf{1} + (1-d)\overline{\mathbf{v}}$$



#### Now, the fun part

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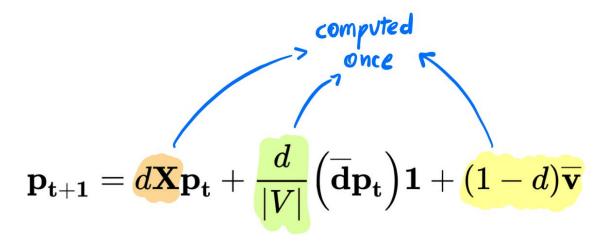
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- Using C++ and Cuda
- At least 80% precision on top-20 results

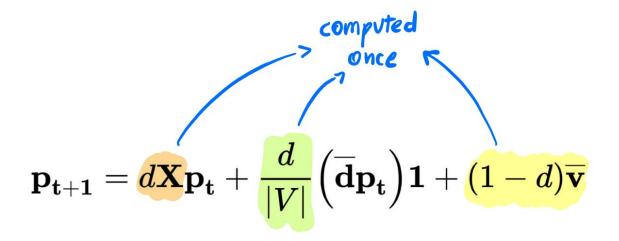
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- Tests run on NVIDIA Laptop GeForce RTX 3080 GPU

### 01 Naive Implementation

#### **PPR with Power method**





Graph is stored in COO format

Just the meaningful ones!

$$\left(\overline{\mathbf{d}}\mathbf{p_{t}}\right)$$

compute\_dangling\_factor\_gpu(dangling\_gpu, pr\_old, pDanglingFact\_gpu, \*pV\_gpu)

$$\left(\overline{\mathbf{d}}\mathbf{p_{t}}\right)$$

Ξ

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template<typename T1, typename T2> __global__ void compute_dangling_factor_gpu( T1 *dangling,  T2* pr, T2 *result, int V){
    for(size_t i = blockIdx.x * blockDim.x + threadIdx.x; i < V; i += blockDim.x * gridDim.x){
        T2 val = (T2) dangling[i] * pr[i];
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- Naive dot product kernel
- Linear global memory access pattern
- atomicAdd synchronized for the whole GPU
- Many threads are stalled!

cooSPMV(x\_gpu,y\_gpu,val\_gpu,pPpr->E,pr\_old,pr\_temp)

 $d\mathbf{X}\mathbf{p_t}$ 

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#### cooSPMV(x\_gpu,y\_gpu,val\_gpu,pPpr->E,pr\_old,pr\_temp)

```
template <typename T1, typename T2> _global__ void cooSPMV_math(
    const T1 * __restrict__ x_gpu,
    const T1 * __restrict__ y_gpu,
    const T2 * __restrict__ val_gpu,
    const T1 E,
    const T2 * __restrict__ pr_old,
    T2 * __restrict__ pr_temp)
{
    size_t index = blockIdx.x * blockDim.x + threadIdx.x;

    if (index < E){
        atomicAdd(&pr_temp[x_gpu[index]], val_gpu[index]*pr_old[y_gpu[index]]);
    }
}</pre>
```

Same inefficiencies of the previous kernel

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- Same inefficiencies of the previous kernel
- Works only with fixed number of threads
- Each thread does one load, computation and store
- Thread setup cost not properly "hided"

compute\_square\_error\_gpu(pr\_old, pr\_gpu, pSquareError\_gpu, pPpr->V)

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template<typename T> __global__ void compute_square_error_gpu( T *old, T *newVector, T *result, int V){
    for (size_t i = blockIdx.x * blockDim.x + threadIdx.x; i < V; i += blockDim.x * gridDim.x) {
        atomicAdd(result, (T)((old[i] - newVector[i]) * (old[i] - newVector[i])));
    }
}</pre>
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- Computes only sum of square of the difference, square root on CPU
- Same inefficiencies of the first kernel

Instances	Average	Minimum	Maximum	Kernel
900	15,5 ms	15 ms	16 ms	cooSPMV
900	6,7 ms	6,5 ms	7,6 ms	compute_square_error_gpu
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- Grid size k1 = (E + Block size 1) / (Block size)

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- Grid size k1 = (E + Block size 1) / (Block size)
- Grid size k2 = (V + Block size 1) / (Block size)
- Grid size k3 = (V + Block size 1) / (Block size)
- Mean computation time (100 iterations, threshold 1e-6) = 1010 ms
- Mean accuracy (100 iterations, threshold 1e-6) = 100%

# After many optimizations...

# Final coo Implementation

# Some improvements

- Floats instead of doubles
- Usage of shared memory
- Some heuristics introduced
- Usage of <u>single precision intrinsics</u>
  - Device only
  - Compiler inline-able set of highly optimized instructions
  - Fast but has low numerical accuracy
  - Allows do things with a small instruction count

Just the meaningful ones!

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template<typename T1, typename T2> __global__ void dangling_kernel(T1*x, T2 *y, T2 *dot, T2 dampingFract, unsigned int n){
    unsigned int index = threadIdx.x + blockDim.x*blockIdx.x;
    unsigned int stride = blockDim.x*gridDim.x;
    extern __shared__ T2 cache[];
    T2 temp = 0.0;
    while(index < n){</pre>
        temp = __fadd_rd(temp, y[x[index]]);
        index += stride:
    cache[threadIdx.x] = temp;
    __syncthreads();
    // reduction
    unsigned int i = blockDim.x/2;
    while(i != 0){
        if(threadIdx.x < i){
            cache[threadIdx.x] = fadd rd(cache[threadIdx.x], cache[threadIdx.x + i]);
        __syncthreads();
        i /= 2:
    if(threadIdx.x == 0){
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dangling\_kernel(pDanglingIndexes\_gpu, pr\_old, pDanglingFact\_gpu,dampingFract,danglingSize);

Pseudo dot product

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  - If a row extends over the 32 elements boundary, then the partial sum of a warp is carried over into the next warp, otherwise the result is written in global memory.
     The decision is taken by the first thread of the next warp.
  - Each warp conduct a segmented scan by looking if 2 values belong to the same row.

```
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 # thread_lane idx [ 0, 0, 0, 0, 0, 0, 1, 1, 1, 1, 2, 2, 2, 2, 2, 2] # row indices val [ 4, 6, 5, 0, 8, 3, 2, 8, 3, 1, 4, 9, 2, 5, 2, 4] # M(i,j) * x(j) \rightarrow result: val [ 4,10,15,15,23,26, 2,10,13,14, 4,13,15,20,22,26] # <math>A(i,j) * x(j)
```

\_\_spmv\_coo\_flat(x\_gpu, y\_gpu, val\_gpu, pr\_old, pr\_gpu, pPpr->E, num\_blocks, ...)

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- Each thread in a warp, except the last one, checks if it is the last of the row. The last thread writes the sum in the output vector
- The last thread in each warp writes its row index and partial sum in 2 global arrays (at an index corresponding to that warp).
- These steps are repeated until the warp reaches the end of its interval
- Finally, the values to carry between warps are added to the result through reduction.

```
template<typename T> _global__ void euclidean_kernel_math(const T * __restrict__ x, const T * __restrict__ y, T * __restrict__ result, unsigned int n)
   size_t index = threadIdx.x + blockDim.x * blockIdx.x;
   size_t stride = blockDim.x * gridDim.x;
   extern __shared__ T cache1[];
   T temp = 0.0;
   while (index < n) {
       temp = __fadd_rd(temp, (
                                 __fmul_rd(
                                         __fsub_rd(x[index],y[index]),
                                         fsub rd(x[index],y[index])
        );
        index += stride;
   cache1[threadIdx.x] = temp;
    __syncthreads();
   size t i = blockDim.x / 2:
   while (i != 0) {
           cache1[threadIdx.x] = __fadd_rd(cache1[threadIdx.x], cache1[threadIdx.x + i]);
        __syncthreads();
       i /= 2;
   if (threadIdx.x == 0) {
        atomicAdd(result, cache1[0]): //
```

euclidean\_kernel\_math(pr\_old, pr\_gpu, pSquareError\_gpu, pPpr->V)

Adaptation of dot product kernel

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Instances	Average	StdDev	Minimum	Maximum	Kernel
900	8,5 ms	0,18 ms	8 ms	9 ms	_spmv_coo_flat
900	0,1 ms	0,09 ms	0,1 ms	0,13 ms	euclidean_kernel_math
900	0,015 ms	0,0098 ms	0,013 ms	0,018 ms	dangling_kernel

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- Grid size k1 = (E + Block size 1) / (Block size)
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- Mean computation time (100 iterations, threshold 1e-6) = 750 ms
- Mean accuracy (100 iterations, threshold 1e-6) = 100%

# Now, the heuristics

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Easily adaptable with different graphs

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Mean execution time: 100 ms

Mean accuracy: 95%

Implementation	Mean total time (reset + exec)	Mean exec time	Accuracy
Naive	1023 ms	1010.62 ms	100%
Cublas coo, cusparse of dot	258 ms	255.081 ms	100%
Final implementation	103 ms	100 ms	>80%

Time (	%) Total Time (ns)	Instances	Avg (ns)	Med (ns)	Min (ns)	Max (ns)	StdDev (ns)	Name
		700	4/ 0/0 507 /	44 005 584 5	45 077 040	00 8/0 070	/01 000 0	ADW - 11 - 12 - 13 - 14 - 15 - 15 - 15 - 15 - 15 - 15 - 15
92 , T2 *)	.0 4,818,817,005	300	16,062,723.4	16,025,571.5	15,933,219	20,742,032	401,228.2	void cooSPMV_math <int, float="">(const T1 *, const T1 *, const T2 *, T2, const T2 *</int,>
6	.0 314,028,668	300	1,046,762.2	1,044,380.0	1,038,236	1,440,922	32,342.9	void euclidean_kernel <float>(T1 *, T1 *, T1 *, unsigned int)</float>
0		300	97,861.1	97,504.0	97,216	128,736	3,080.1	<pre>void dot_product_kernel_math<int, float="">(const T1 *, const T2 *, T2 *, T2, unsig</int,></pre>
ned int 0		300	82,705.6	82,703.5	77,344	85,312	737.6	void vectorScalarAddAndIncrement_math <float>(T1, T1 *, int, int, T1)</float>
0	.4 23,175,290	300	77,251.0	77,183.0	75,776	79,392	650.6	void copy_vector <float>(T1 *, T1 *, int)</float>
0	.3 13,469,514	10	1,346,951.4	1,332,683.0	1,322,139	1,487,226	49,449.1	void vectorScalarMul <float>(T1, T1 *, int)</float>
0	.2 12,004,502	300	40,015.0	39,968.0	39,296	41,344	324.3	<pre>void init_vector<float>(T1 *, int, T1)</float></pre>

For all the implementations seen until now the Time (%) were similar. Can we skip the SPMV multiplication?

Implementation	Mean total time (reset + exec)	Mean exec time	Accuracy
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Cublas coo, cusparse of dot	258 ms	255.081 ms	100%
Final implementation	103 ms	100 ms	>80%
Monte Carlo complete path	17 ms	14.2666 ms	>80%

# Monte Carlo Implementation

#### **Monte Carlo Complete path**

- Simulate m random walks starting from the personalized vertex.
- At each step the walker can terminate with probability 1-alpha.
- At each step the walker move to a random vertex connected to an outgoing edge with probability alpha.

Let  $N_j$  be the number of visits by all walkers to vertex j, then the pr of vertex j is defined as following:

$$pr_j(personalized\ vertex, alpha) = (1 - alpha)\ \frac{1}{m}N_j\ (personalized\ vertex)$$

#### **Monte Carlo Complete path**

Since we don't care about the values we can skip the multiplication and division

$$pr_j(personalized\ vertex, alpha) = (1 - alpha)\frac{1}{m}N_j\ (personalized\ vertex)$$



 $pr_j(personalized\ vertex, alpha) = N_j(personalized\ vertex)$ 

#### MC Complete path on CPU

- Uses CSC
   (Compressed sparse column matrix format)
- Walkers run sequentially.

```
int walkLen = 0;
for (int run = 0; run < nWalkers; run++) {</pre>
    int vIdx = s;
   while (walkLen < maxWalkLen){</pre>
        pPpr->pr[vIdx]+=1.0f;
        bool terminate = dist(mt) < 1.0f-pPpr->alpha;//Terminate with probability 1-c
        if (terminate) {
            break:
       } else {
            int neighStartIdx = csc.xPtr_cpu[vIdx];
            int neighEndIdx = csc.xPtr_cpu[vIdx+1];
            int neighSize = neighEndIdx - neighStartIdx;
            if (neighSize == 0) {
                break;}
            //select a random connected vertex
            int rIdx = vJump(mt)%neighSize;
            vIdx=csc.x_cpu[neighStartIdx+rIdx];
        walkLen++;
   //std::cout << "Walk length: " << walkLen << std::endl;
   walkLen = 0;
```

## MC Complete path on GPU naive

- 1thread = 1 walker
- If we try to run the walkers in parallel we need an atomic add and in the first step each walker thread will add to the same position

```
const unsigned int i = blockIdx.x * blockDim.x + threadIdx.x;
if (i < walkers) {
    int vIdx = s:
    while (walkLen < maxWalkLen){
        atomicAdd( address: &pr[vIdx], val: 1.0f);//Add 1 to each visited node
        bool terminate = dist( &: mt) < 1.0f-pPpr->alpha;//Terminate with probability 1-c
        if (terminate) {
            break;
            int neighStartIdx = csc.xPtr_cpu[vIdx];
            int neighSize = neighEndIdx - neighStartIdx;
            int rIdx = vJump( &: mt)%neighSize;
            vIdx=csc.x_cpu[neighStartIdx+rIdx];
        walkLen++;
```

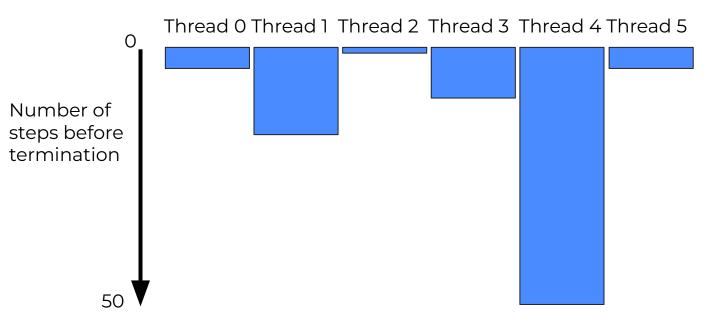
#### MC Complete path on GPU naive

Solution:

 Each thread starts
 from a random
 vertex between the outgoing edges.
 In this way the atomic adds act on different memory locations

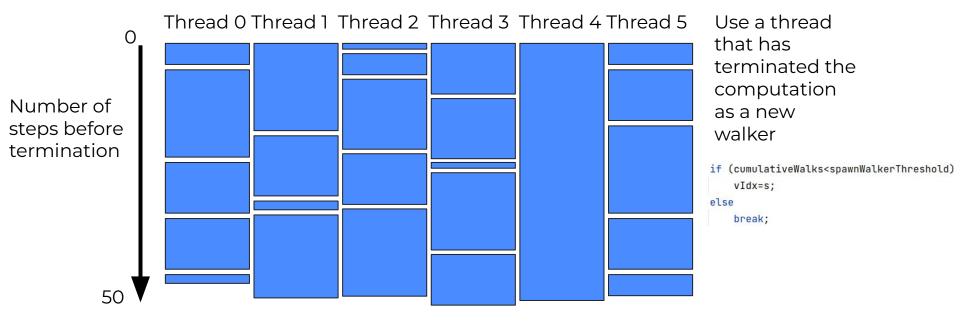
```
//Starting from neighbors of s
unsigned int rIdx0 = curand( state: &localState) % neighSizee[s];
int vIdx = x[neighStartPos + rIdx0];
int cumulativeWalks = 1;
```

#### Thread utilization



Most threads do only a few steps and then terminate the random walk

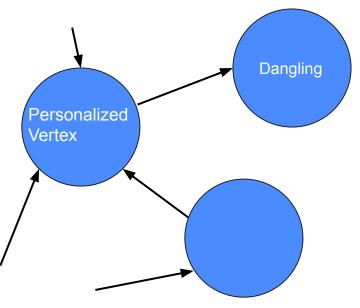
#### Thread utilization solution



Leads to a much higher accuracy due to the larger number of samples

#### Dangling nodes in MC

We have to consider an additional case



```
error in value! (0) correct=1288741 (val=0.15), found=1288741 (val=7.68e+05)
error in value! (1) correct=975307 (val=0.1275), found=975307 (val=6.53e+05)
error in rank! (2) correct=1702309 (val=0.001989), found=3566906 (val=2.804e-07)
error in rank! (3) correct=500469 (val=0.0009716), found=3566905 (val=2.804e-07)
error in rank! (4) correct=1518892 (val=0.0009589), found=3566904 (val=2.804e-07)
error in rank! (5) correct=24716 (val=0.0008311), found=3566903 (val=2.804e-07)
error in rank! (6) correct=689491 (val=0.0008219), found=3566902 (val=2.804e-07)
error in rank! (7) correct=195101 (val=0.0007462), found=3566901 (val=2.804e-07)
error in rank! (8) correct=932394 (val=0.00071), found=3566900 (val=2.804e-07)
error in rank! (9) correct=2984189 (val=0.0006991), found=3566899 (val=2.804e-07)
error in rank! (10) correct=1835017 (val=0.0006721), found=3566898 (val=2.804e-07)
error in rank! (11) correct=577659 (val=0.0006673), found=3566897 (val=2.804e-07)
error in rank! (12) correct=28196 (val=0.0006478),
                                                   found=3566896 (val=2.804e-07)
error in rank! (13) correct=7030 (val=0.0006427),
                                                  found=3566895 (val=2.804e-07)
error in rank! (14) correct=2257865 (val=0.0006064)
                                                     found=3566894 (val=2.804e-07)
error in rank! (15) correct=2532493 (val=0.0005939)
                                                     found=3566893 (val=2.804e-07)
error in rank!
               (16) correct=28020 (val=0.0005909),
                                                   found=3566892 (val=2.804e-07)
error in rank! (17) correct=2979297 (val=0.0005655)
                                                    found=3566891 (val=2.804e-07)
error in rank! (18) correct=9008 (val=0.0004983),
                                                  found=3566890 (val=2.804e-07)
error in rank! (19) correct=27566 (val=0.0004868),
                                                   found=3566889 (val=2.804e-07)
```

## Dangling nodes, MC

Correct first two places error in value (0) correct=1288741 (val=0.15), found=1288741 (val=7.68e+05) error in value correct=975307 (val=0.1275), found=975307 (val=6.53e+05) error in rank! (2) correct=1702309 (val=0.001989), error in rank! (3) correct=500469 (val=0.0009716), error in rank! (4) correct=1518892 (val=0.0009589) error in rank! (5) correct=24716 (val=0.0008311), error in rank! (6) correct=689491 (val=0.0008219), error in rank! (7) correct=195101 (val=0.0007462), error in rank! (8) correct=932394 (val=0.00071), for error in rank! (9) correct=2984189 (val=0.0006991) error in rank! (10) correct=1835017 (val=0.0006721) Same ranking as Initial pr values error in rank! (11) correct=577659 (val=0.0006673) dangling but shifted error in rank! (12) correct=28196 (val=0.0006478), by one towards the error in rank! (13) correct=7030 (val=0.0006427), bottom error in rank! (14) correct=2257865 (val=0.0006064 error in rank! (15) correct=2532493 (val=0.0005939) error in rank! (16) correct=28020 (val=0.0005909), error in rank! (17) correct=2979297 (val=0.0005655) error in rank! (18) correct=9008 (val=0.0004983), error in rank! (19) correct=27566 (val=0.0004868)

## Solution: Init pr with values for dangling

100% accuracy for vertices with outgoing edges towards dangling

100% accuracy and 2ms reset+execution time for danglings

# **Other MC optimizations**

- Copy curand state in local memory
- Use an array to store outgoing edges size

# Q&A Additional slides

# Performance of every implementation

Implementation	Mean total time (reset + exec)	Mean exec time	Accuracy
Naive	1023 ms	1010.62 ms	100%
Cublas bcr test	1304 ms	1215.45 ms	100%
Improved naive	814 ms	800.903 ms	100%
Cublas coo, cusparse of dot	258 ms	255.081 ms	100%
Final implementation	103 ms	100 ms	>80%
Monte Carlo complete path	17 ms	14.2666 ms	>80%

Tested on rtx 3080 laptop, 100 runs