

A Brief Introduction to LFSR and NLFSR

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Chapter 1

The periodicity of LFSR sequences

1.1 Concepts

Definition 1.1.1. In \mathbb{F}_q , a infinite sequence $a = (a_0, a_1, a_2, \dots)$ is periodic if $\exists l, s.t. a_{l+k} = a_k, \forall k \geq 0$. The period of a is denoted by $p(a)$.

Lemma 1.1.2. $a = (a_0, a_1, a_2, \dots)$. If $\exists l, s.t. a_{l+k} = a_k$, then $p(a)|l$.

Lemma 1.1.3. $f = c_0 + c_1x + c_2x^2 + \dots + c_nx^n$ ($c_0c_n \neq 0$), $a \in G(f)$. $a = (a_0, a_1, \dots)$.

1. $S_0T^{p(a)} = S_0$

2. $S_0, S_0T, S_0T^2, \dots, S_0T^{p(a)-1}$ are distinct.

3. l is the minimal positive integer s.t. $S_0T^l = S_0$, then $p(n) = l$.

Theorem 1.1.4. $f = c_0 + c_1x + c_2x^2 + \dots + c_nx^n$ ($c_0c_n \neq 0$). $a \in G(f)$. a is periodic, and $p(a) \leq q^n - 1$.

Proof. Suppose $a = (a_0, a_1, a_2, \dots)$. The states are

$$S_0, S_1 = S_0T, S_2 = S_0T^2, \dots$$

If $\exists S_i = (0, 0, \dots, 0)$, then all the states of a are zero. $a = (0, 0, \dots)$, thus $p(a) = 1$.

If $a \neq (0, 0, \dots)$, then $\forall S_i \neq (0, 0, \dots, 0)$. In \mathbb{F}_q , there are $q^n - 1$ non-zero n -dimensional row vectors. Thus in the first q^n states, there must be at least two are the same. i.e. $\exists i, j, 0 \leq i < j \leq q^n - 1, S_0T^i = S_0T^j$. $\Rightarrow S_0T^{j-i} = S_0$. Let $l = j - i$, we have $a_{l+k} = a_k, \forall k \geq 0 \Rightarrow a$ is periodic and $p(a) \leq l \leq q^n - 1$ \square

Note:

1. If $p(a) = q^n - 1$, a is called **q-ary m-sequence**.

2. $a \in G(f)$ is periodic, $\exists f$, s.t. $a \in G(f)$?. The answer is yes, because we have $a_k - a_{k-l} = 0, k \geq l$. Thus $f(x) = 1 - x^l$.

3. No matter what the initial state S_0 is, f generates periodic sequence.

1.2 Minimal Polynomials

Theorem 1.2.1. a is a periodic sequence over \mathbb{F}_q , $\exists! f(x) \in \mathbb{F}_q$, s.t. $a \in G(h)$ if and only if $f(x)|h(x)$

($I = \{h(x) \in \mathbb{F}_q | a \in G(h)\}$). Only need to prove $I \triangleleft \mathbb{F}_q$, because $a \in G(h) \Leftrightarrow h \in I \Leftrightarrow f|h$)

Proof. 1. $p(a) = l \Rightarrow a_k - a_{k-l} = 0 \Rightarrow 1 - x^l \in I, I$ is non-zero.

2. $g(x), h(x) \in I$. Suppose $a \in G(g), a \in G(h)$. It's easy to conclude that $a \in G(g - h)$. Thus we have $g - h \in I$. Hence, I is a additive subgroup of \mathbb{F}_q

3. $\forall h(x) \in \mathbb{F}_q$, if $g(x) \in I, a \in G(g) \Rightarrow a \in G(hg) \Rightarrow hg \in I$. I is closed under multiplication.

Therefore, $I \triangleleft \mathbb{F}_q \Rightarrow \exists f(x)$ s.t. $I = (f(x))$. Naturally, we can suppose the coefficient of zero-order term of $f(x)$ is 1. If $f_1(x)$ satisfies the same property as $f(x)$, then $f(x)|f_1(x), f_1(x)|f(x) \Rightarrow f(x) = cf_1(x) \Rightarrow c = 1$. Therefore, the polynomial is unique. \square

Definition 1.2.2. f is the minimal polynomial of a .

Theorem 1.2.3. $f(x) \in \mathbb{F}_q, c_0 = 1, \exists$ a periodic sequence whose minimal polynomial is $f(x)$.

Proof. $\deg f(x) = n, (1)n = 0$, zero sequence have $f(x) = 1$ as minimal polynomial.

(2) $n > 0$, suppose $f(x) = 1 + c_1x + c_2x^2 + \cdots + c_nx^n$. Let $S_0 = (0, 0, \dots, 0, 1)$. Let $a \in G(f)$ whose initial state is S_0 satisfying

$$a_k + c_1a_{k-1} + c_2a_{k-2} + \cdots + c_na_{k-n} = 0, k \geq n$$

$a = (a_0, a_1, a_2, \dots)$. $a_0 = a_1 = a_2 = \cdots = a_{n-2} = 0, a_{n-1} = 1$. If $h(x)$ is a minimal polynomial of a , then $h|f$.

If $\deg h(x) = \deg f(x)$, $f(x)$ is the minimal polynomial of (a) .

If $f(x)$ is not the minimal polynomial of a , then $\deg f(x) > \deg h(x) = m$. $h(x) = 1 + d_1x + \cdots + d_mx^m \Rightarrow a_k + d_1a_{k-1} + \cdots + d_ma_{k-m} = 0, k \geq 1$. i.e. $(a_0, a_1, \dots, a_{m-1})$ is the initial state of a . Then a is zero sequence. This leads a contradiction. \square

Note: $\forall f(x) \in \mathbb{F}_q, \exists$ a periodic sequence whose minimal polynomial is $f(x)$.

Definition 1.2.4. $f(x) \in \mathbb{F}_q[x], \deg f \geq 1$. The period of $f(x)$ is $p(f) = \min\{l \mid f(x)|x^l - 1\}$.

Definition 1.2.5. $p(f) = \text{ord } x$, in $\mathbb{F}_q[x]_{f(x)}^*$.

Lemma 1.2.6. $f(x)|(x^l - 1) \Rightarrow p(f)|l$.

Lemma 1.2.7. $p(f) = p(\tilde{f})$.

Proof. $p(f) = \min\{l : f(x)|x^l - 1\}, \deg(f) = n$

Let $p(\tilde{f}) = l$.

$$x^l - 1 = \tilde{f}(x)g(x)$$

Substitute x with x^{-1} ,

$$(x^{-1})^l - 1 = \tilde{f}(x^{-1})g(x^{-1})$$

Let $h(x) = x^{l-n}g(x^{-1})$,

$$1 - x^l = x^n \tilde{f}(x^{-1})h(x)$$

$$f(x) = x^n \tilde{f}(x^{-1})|(x^l - 1)$$

Thus we have

$$p(f)|p(\tilde{f})$$

For the same reason,

$$p(\tilde{f})|p(f)$$

Hence,

$$p(f) = p(\tilde{f})$$

\square

Lemma 1.2.8. $f(x) \in \mathbb{F}_q[x]$ is irreducible, $\deg f(x) = n$. $p(f) \mid q^n - 1$.

Definition 1.2.9. $A \in GL(\mathbb{F}_q)$, if $\exists l$, s.t. $A^l = I$, $p(A) \triangleq \min\{l \mid A^l = I\}$.

Lemma 1.2.10. $A \in GL(\mathbb{F}_q)$, $\exists l$, s.t. $A^l = I$, and $p(A)|l$.

Lemma 1.2.11. $p(f) = p(T)$.

Proof. $f(x) = 1 + c_1x + \cdots + c_nx^n \in \mathbb{F}_q[x]$. T is a matrix determined by $f(x)$.

$$\tilde{f}(x)|x^{p(\tilde{f})} - 1, \tilde{f}(T) = 0 \Rightarrow T^{p(\tilde{f})} = I \Rightarrow p(T)|p(\tilde{f}).$$

$$T^{p(T)} = I \Rightarrow T \text{ satisfies } x^{p(T)} - 1 = 0.$$

$$\tilde{f}(x) \text{ is the minimal polynomial of } T. \Rightarrow \tilde{f}|x^{p(T)} - 1. \Rightarrow p(\tilde{f})|p(T)$$

$$\Rightarrow p(\tilde{f}) = p(T).$$

$$p(f) = p(\tilde{f}) \Rightarrow p(f) = p(T). \quad \square$$

Theorem 1.2.12. a is a LFSR sequence whose minimal polynomial is $f(x)$, then $p(a) = p(f)$.

Proof. a satisfies $a_k - a_{k-p(a)} = 0, k \geq p(a)$.

$$h(x) = 1 - x^{p(a)}, a \in G(h) \Rightarrow f|h = 1 - x^{p(a)} \Rightarrow p(f)|p(a)$$

$$T^{p(T)} = I \Rightarrow S_k T^{p(T)} = S_k \Rightarrow S_{k+p(T)} = S_k \Rightarrow a_{k+p(T)} = a_k, \forall k \geq 0 \Rightarrow p(T)|p(a)$$

$$p(f) = p(T) \Rightarrow p(f)|p(a)$$

Thus we have

$$p(a) = p(f) \quad \square$$

Corollary 1.2.13. $f(x) \in \mathbb{F}_q[x]$, irreducible. $\forall a \in G(f), p(a) = p(f)$

Corollary 1.2.14. $a \in G(f), S_0 = (0, 0, \dots, 0, 1)$ is the initial state of a . $p(a) = p(f)$

Corollary 1.2.15. $\forall a \in G(f), p(a)|p(f), f(x) \in \mathbb{F}_q[x], \deg f(x) \geq 1$

Chapter 2

Shift Equivalent Class

2.1 Shift Equivalent Class

Definition 2.1.1. $a = (a_0, a_1, a_2, \dots)$. Left shifting transform L is an operator over $G(f)$:

$$L(a) = (a_1, a_2, a_3, \dots)$$

$$L^0(a) = a$$

$$L^t(a) = L(L^{t-1}(a)), t \geq 1$$

Definition 2.1.2. If a and b are shift equivalent, $a \sim b$ if $\exists t \geq 0, b = L^t(a)$.

Lemma 2.1.3. $a \sim b$ is an equivalence relation.

Proof. (1) $a \sim a$

$$(2) a \sim b \Rightarrow b \sim a$$

$$(3) a \sim b, b \sim c \Rightarrow a \sim c$$

□

Lemma 2.1.4. (1) $a \sim b \Rightarrow p(a) = p(b)$

$$(2) L^{p(a)}(a) = L^0(a) = a$$

$$(3) \{b | b \sim a\} = a, L(a), L^2(a), \dots, L^{p(a)-1}(a)$$

Definition 2.1.5. A set of q -ary periodic sequences is shift equivalent class(C) if $\forall a, b \in C \Rightarrow a \sim b$, and $d \sim a \Rightarrow d \in C$

Theorem 2.1.6. a is q -ary periodic sequence. The shift equivalent class C that contains a :

$$C = a, L(a), L^2(a), \dots, L^{p(a)-1}(a)$$

C_1 and C_2 are any two shift equivalent classes. $C_1 = C_2$ or $C_1 \cap C_2 = \emptyset$

Corollary 2.1.7. $|C| = p(a)$

Corollary 2.1.8. $G(f)$ is a disjoint sum of several shift equivalent classes.

Proof. $a \in G(f) \Rightarrow L(a) \in G(f) \Rightarrow C \subset G(f)$

□

Theorem 2.1.9. $f(x) = f_1(x)f_2(x) \cdots f_r(x)$. $\deg f_i = n_i$. $\gcd(f_i, f_j) = 1, i \neq j$. Then

$$G(f) = G(f_1) \oplus G(f_2) \oplus \cdots \oplus G(f_r)$$

Proof. If $a \in G(f_i) \Rightarrow \exists g_i$, s.t. $g_i|f_i \Rightarrow G(f_i) \subset G(f)$. i.e. $G(f_i)$ is a subspace of $G(f)$.

$\deg f_i = n_i \Rightarrow$ suppose $(a_{i_1}, a_{i_2}, \dots, a_{i_{n_i}})$ is a basis of $G(f_i)$, $1 \leq i \leq r$

First, we need to prove

$$(a_{11}, \dots, a_{1_{n_1}}, a_{21}, \dots, a_{2_{n_2}}, \dots, a_{r1}, \dots, a_{r_{n_r}}) \quad (2.1)$$

is a basis of $G(f_1) + G(f_2) + \dots + G(f_r)$.

Assume $\exists c_i (i = 1, 2, \dots, r, \quad j = 1, 2, \dots, n_i)$ s.t.

$$\sum_{j=1}^{n_i} c_{ij} a_{ij} = 0$$

Let

$$a_i = \sum_{j=1}^{n_i} c_{ij} a_{ij} \in G(f_i)$$

$$a_1 = -(a_2 + \dots + a_r) \in G(f_2) + \dots + G(f_r)$$

Notice that

$$f_i | f_2 f_3 \dots f_r, i = 2, 3, \dots, r$$

Thus we have

$$G(f_2) + \dots + G(f_r) \subset G(f_2 f_3 \dots f_r)$$

Then

$$a_1 \in G(f_2 f_3 \dots f_r)$$

Suppose the minimal polynomial of a_1 is $m(x)$.

$$a_1 \in G(f_1) \Rightarrow m(x) | f_1(x)$$

$$a_1 \in G(f_2 f_3 \dots f_r) \Rightarrow m(x) | f_2 \dots f_r$$

Notice that $\gcd(f_1, f_2 \dots f_r) = 1$, then $m(x) = 1$.

Hence,

$$a_1 = (0, 0, \dots)$$

For the same reason

$$a_2 = \dots = a_r = 0$$

Thus

$$\sum_{j=1}^{n_i} c_{ij} a_{ij} = 0, i = 1, \dots, r$$

Then

$$c_{ij} = 0, i = 1, \dots, r, j = 1, \dots, n_i$$

Therefore (3) is linearly independent.

$$f = f_1 \dots f_r \Rightarrow n = n_1 + n_2 + \dots + n_r$$

(3) is a basis of $\sum_{i=1}^r G(f_i)$. Then $\sum_{i=1}^r G(f_i)$ is a direct sum.

Notice that

$$\dim G(f) = n = \dim G(f_1) + \dots + \dim G(f_r)$$

Hence

$$G(f) = G(f_1) \oplus G(f_2) \oplus \dots \oplus G(f_r)$$

□

2.2 The Decompositions of $G(f)$

Lemma 2.2.1. *If $\gcd(f_1(x), f_2(x)) = 1$, $a \in G(f_1), b \in G(f_2)$, then $p(a+b) = \text{lcm}[p(a), p(b)]$*

Proof. Let $l = \text{lcm}[p(a), p(b)], l' = p(a+b)$,

$$a = (a_0, a_1, a_2, \dots)$$

$$b = (b_0, b_1, b_2, \dots)$$

Let s be the period of

$$(a_0, b_0), (a_1, b_1), (a_2, b_2), \dots$$

On the one hand $a_l = a_0, b_l = b_0 \Rightarrow (a_l, b_l) = (a_0, b_0)$, then $s|l$.

$$(a_s, b_s) = (a_0, b_0) \Rightarrow a_s + b_s = a_0 + b_0 \Rightarrow l'|s. \text{ Thus we have } \Rightarrow l'|l$$

On the other hand, $p(a+b) = l' \Rightarrow$

$$a_{l'+k} + b_{l'+k} = a_k + b_k, \forall k \geq 0$$

$$(a_{l'+k} - a_k) + (b_{l'+k} - b_k) = 0, \forall k \geq 0$$

$$(L^{l'}(a) - a) + (L^{l'}(b) - b) = 0$$

$$\gcd(f_1, f_2) = 1 \Rightarrow G(f_1) \oplus G(f_2)$$

$$L^{l'}(a) - a \in G(f_1) \Rightarrow L^{l'}(a) = a \Rightarrow p(a)|l'$$

$$L^{l'}(b) - b \in G(f_2) \Rightarrow L^{l'}(b) = b \Rightarrow p(b)|l'$$

$$\Rightarrow l = \text{lcm}[p(a), p(b)]|l'$$

In conclusion, $l = l'$

□

Corollary 2.2.2. *$\gcd(f_i, f_j) = 1, i \neq j, a_i \in G(f_i), 1 \leq i \leq r$, then $p(a_1+a_2+\dots+a_r) = \text{lcm}[p(a_1), p(a_2), \dots, p(a_r)]$.*

Theorem 2.2.3. *$f = f_1 f_2$. $\gcd(f_1, f_2) = 1, \deg(f_i) = n_i$.*

Suppose $G(f_1)$ consists of two shift equivalent classes: C_{11}, C_{12} .

Suppose $G(f_2)$ consists of two shift equivalent classes: C_{21}, C_{22} . $|C_{ij}| = p_{ij}$.

$G(f)$ consists of

$$\sum_{k_2=1}^2 \sum_{k_1=1}^2 p_{1k_1} p_{2k_2} / \text{lcm}[p_{1k_1} p_{2k_2}]$$

shift equivalent classes

Proof.

$$G(f) = G(f_1) \oplus G(f_2)$$

$$= (C_{11} \cup C_{12}) \oplus (C_{21} \cup C_{22})$$

$$= (C_{11} \oplus C_{21}) \cup (C_{11} \oplus C_{22}) \cup (C_{12} \oplus C_{21}) \cup (C_{12} \oplus C_{22})$$

$|C_{11} \oplus C_{21}| = p_{11} p_{21}$, then the period of every sequence in $C_{11} \oplus C_{21}$ is $\text{lcm}[p_{11}, p_{21}]$ according to

Lemma 3.2.3. Thus $C_{11} \oplus C_{21}$ consists

$$p_{11} p_{21} / \text{lcm}[p_{11}, p_{21}]$$

shift equivalent classes.

Hence $G(f)$ consists of

$$\sum_{k_2=1}^2 \sum_{k_1=1}^2 p_{1k_1} p_{2k_2} / \text{lcm}[p_{1k_1} p_{2k_2}]$$

shift equivalent classes.

□

Corollary 2.2.4. $f = f_1 \cdot f_2 \cdots f_r$, $\deg(f) = n_i \geq 1$, $\gcd(f_i, f_j) = 1$ ($i \neq j$)

$G(f_i)$ consists of m_i shift equivalent classes, ($i = 1, 2, 3, \dots, r$), i.e. $C_{i1}, C_{i2}, \dots, C_{im_i}$

The period of sequences in C_{ij} is p_{ij} ($i = 1, 2, \dots, r; j = 1, 2, \dots, m_i$). Then $G(f)$ consists of

$$\sum_{k_1=1}^{m_1} \sum_{k_2=1}^{m_2} \cdots \sum_{k_r=1}^{m_r} \frac{p_{1k_1} \cdot p_{2k_2} \cdots p_{rk_r}}{\text{lcm}[p_{1k_1}, p_{2k_2}, \dots, p_{rk_r}]}$$

shift equivalent classes.

Lemma 2.2.5. $f(x) \in F_q[x]$, $c_0 = 1$. irreducible. $m = \min\{i \in \mathbb{Z}_+ | p^i \geq e\}$, ($p^m \geq e > p^{m-1}$)

$$\implies p(f^e) = p^m \cdot p(f)$$

Proof. On the one hand,

$$f \mid x^{p(f)} - 1 \implies f^{p^m} \mid (x^{p(f)} - 1)^{p^m} = x^{p^m \cdot p(f)} - 1$$

Notice that $p^m \geq e \implies f^e \mid f^{p^m}$, thus we have

$$f^e \mid x^{p^m \cdot p(f)} - 1 \implies p(f^e) \mid p^m \cdot p(f)$$

On the other hand, suppose

$$p(f) = p^j \cdot t, p \nmid t$$

Then we have

$$f^e \mid (x^{p^j t} - 1) = (x^t - 1)^{p^j}$$

$p \nmid t \implies (x^t - 1)' = t \cdot x^{t-1} \neq 0 \implies \gcd(x^t - 1, (x^t - 1)') = 1$, then $x^t - 1$ is irreducible. f is irreducible, thus

$$f \mid x^t - 1, e \leq p^j$$

Notice that $m = \min\{i \in \mathbb{Z}_+ | p^i \geq e\}$, then $m \leq j$, Thus $p(f) \mid t \implies p^m p(f) \mid p^m t \implies p^m p(f) \mid p^j t = p(f^e)$
Hence, we can conclude that $p(f^e) = p^m p(f)$. \square

Note: $p^m \geq e > p^{m-1} \implies e = p^{m-1} + 1, p^{m-1} + 2, \dots, p^m$, then $p(f^e) = p^m \cdot p(f)$

$$p(f^2) = p(f^3) = \cdots = p(f^p) = p \cdot p(f)$$

$$p(f^{p+1}) = p(f^{p+2}) = \cdots = p(f^{p^2}) = p^2 \cdot p(f)$$

$$\vdots$$

$$p(f^{p^{i-1}+1}) = p(f^{p^{i-1}+2}) = \cdots = p(f^{p^i}) = p^i \cdot p(f)$$

Theorem 2.2.6. $f(x) \in F_q[x]$, $c_0 = 1$, irreducible, $\deg(f) = n \cdot q = p^r$, $m = \min\{i \in \mathbb{Z}_+ | p^i \geq e\}$

The periods of sequences in $G(f^e)$ are

$$1, p(f), p^j p(f) (j = 1, 2, \dots, m-1), p^m p(f)$$

Respectively, the numbers of the sequences which have the periods above are

$$1, q^n - 1, q^{np^i} - q^{np^{i-1}} (i = 1, 2, \dots, m-1), q^{ne} - q^{np^{m-1}}$$

Thus, the numbers of the shift equivalent class in $G(f^e)$ which have the periods above are.

$$1, \frac{q^n - 1}{p(f)}, \frac{q^{np^i} - q^{np^{i-1}}}{p^j p(f)}, (i = 1, 2, \dots, m-1), \frac{q^{ne} - q^{np^{m-1}}}{p^m p(f)}$$

Proof. f is irreducible, then $\forall a \in G(f^e)$ the minimal polynomial of a is $f(x)$. And we have

$$G(f^0) \subset G(f^1) \subset G(f^2) \subset \cdots G(f^e)$$

Let's prove the most simple cases first.

$$0. \quad G(f^0) = 0. \quad |G(f^0)| = 1 \quad p(0) = 1$$

1. $\forall a \in G(f^1) \setminus G(f^0)$, $p(a) = p(f)$ (f is irreducible), $|G(f^1) \setminus G(f^0)| = q^n - 1$. Then $G(f^1) \setminus G(f^0)$ has $\frac{q^n-1}{p(f)}$ shift equivalent classes.

2.

$$\forall a \in \left\{ \begin{array}{ll} G(f^2) \setminus G(f^1) & p(a) = p(f^2) \\ G(f^3) \setminus G(f^2) & p(a) = p(f^3) \\ \cdots & \cdots \\ G(f^{i+1}) \setminus G(f^i) & p(a) = p(f^{i+1}) \\ \cdots & \cdots \\ G(f^p) \setminus G(f^{p-1}) & p(a) = p(f^p) \end{array} \right\} \implies p(a) = p(f^2) = \cdots = p(f^p) = p \cdot f(p)$$

$|G(f^p) \setminus G(f^1)| = q^{np} - q^n$, $G(f^p) \setminus G(f^1)$ has $\frac{q^{np}-q^n}{p \cdot p(f)}$ shift equivalent classes.

And so on,

$$\forall a \in G(f^{p^i}) \setminus G(f^{p^{i-1}}), p(a) = p(f^{p^{i-1}+1}) = \cdots = p(f^{p^i}) = p^i \cdot p(f)$$

$|G(f^{p^i}) \setminus G(f^{p^{i-1}})| = q^{np^i} - q^{np^{i-1}}$, thus $G(f^{p^i}) \setminus G(f^{p^{i-1}})$ has $\frac{q^{np^i}-q^{np^{i-1}}}{p^i \cdot p(f)}$ shift equivalent classes.

($i = 1, 2, \dots, m-1$)

$$\forall a \in G(f^{p^m}) \setminus G(f^{p^{m-1}}), p(a) = p^m \cdot p(f)$$

$|G(f^{p^m}) \setminus G(f^{p^{m-1}})| = q^{np^m} - q^{np^{m-1}}$, thus $G(f^{p^m}) \setminus G(f^{p^{m-1}})$ has $\frac{q^{np^m}-q^{np^{m-1}}}{p^m \cdot p(f)}$ shift equivalent classes.

□

Definition 2.2.7. All the states of LFSR compose of a set. We denote it by $V_n(\mathbb{F}_q)$.

$$V_n(\mathbb{F}_q) = \{(a_1, a_2, \dots, a_n) | a_1, a_2, \dots, a_n \in \mathbb{F}_q\}$$

While running LFSR, there is a transform acting in such set. We call it State-transition transform, denoted by T_f .

$$T_f : V_n(\mathbb{F}_q) \longrightarrow V_n(\mathbb{F}_q)$$

$$S_k \longmapsto S_k T (= S_{k+1})$$

Chapter 3

m-sequence and its sampling

3.1 m-sequence

Definition 3.1.1. $a = (a_0, a_1, a_2, \dots)$ is a q -ary n -level LFSR sequence, satisfying

$$a_k + c_1 a_{k-1} + c_2 a_{k-2} + \dots + c_n a_{k-n} = 0 \quad (k \geq n, c_n \neq 0)$$

If $p(a) = q^n - 1$, then a is called **m-sequence**.

Theorem 3.1.2. $f(x) \in \mathbb{F}_q[x], f(x) = 1 + c_1 x + c_2 x^2 + \dots + c_n x^n \quad (c_n \neq 0)$, a is a nonzero sequence in $G(f)$. If a is a m-sequence, then

- (1) $L^k(a)$ is m-sequence, for $k = 0, 1, 2, \dots$
- (2) $a, L(a), L^2(a), \dots, L^{q^n-2}(a)$ are all nonzero sequence in $G(f)$. and $L^{q^n-1}(a) = a$.
- (3) $S_0, S_1, \dots, S_{q^n-2} \in V_n(\mathbb{F}_q)$ distinct nonzero $q^n - 1$ and $S_{q^n-1} = S_0$

Proof. (1)(2) can be concluded from Lemma 3.4

- (3) If $\exists S_i = S_j, (0 \leq i < j < q^n - 1)$, then $L^i(a) = L^j(a) \implies i = j$.

□

Corollary 3.1.3. $f(x) \in \mathbb{F}_q[x], f(x) = 1 + c_1 x + c_2 x^2 + \dots + c_n (c_n \neq 0)$, $a(\neq 0) \in G(f)$. Suppose $t_1 > t_2 > 0$,. If a is a m-sequence,

$$L^{t_1}(a) + L^{t_2}(a) \in G(f)$$

is a m-sequence when $(q^n - 1) \nmid (t_1 - t_2)$.

Theorem 3.1.4. In \mathbb{F}_2 , a is a periodic sequence. $\forall i, j (0 \leq i, j \leq p(a) - 1), L^i(a) + L^j(a) = 0$ or $L^k(a) (0 \leq k \leq p(a) - 1)$, then

$\exists n$, s.t. a is a m-sequence with $p(a) = 2^n - 1$.

Proof. Let $L^i(a)$ denote a consecutive sequence in a with $p(a)$ elements $(0 \leq i \leq p(a) - 1)$

$$L^i(a) = (a_i, a_{i+1}, \dots, a_{p(a)-1}, a_0, a_1, \dots, a_{i-1})$$

It's easy to verify that

$$V \triangleq 0 \cup \{L^i(a) | 0 \leq i \leq p(a) - 1\}$$

is a multiplicative group.

Define operations on V .

$$0(c_0, c_1, \dots, c_{p(a)-1}) = (0, 0, \dots, 0)$$

$$1(c_0, c_1, \dots, c_{p(a)-1}) = (c_0, c_1, \dots, c_{p(a)-1})$$

The V is a vector space over \mathbb{F}_2 . Suppose $p(a) = 2^n - 1$.

$\dim(V) = n$. Thus we can assume $L^0(a), L^1(a), L^2(a), \dots, L^{r-1}(a)$ is linearly independent over \mathbb{F}_2 while $L^0(a), L^1(a), L^2(a), \dots, L^{r-1}(a), L^r(a)$ is linearly dependent over \mathbb{F}_2 . Then we have for $r \leq n$

$$L^r(a) = c_1 L^{r-1}(a) + c_2 L^{r-2}(a) + \dots + c_{r-1} L^1(a) + c_r L^0(a), c_i \in \mathbb{F}_2$$

Let $L^{k-r}(k \geq r)$ act over two sides of the equation.

$$L^k(a) = c_1 L^{k-1}(a) + c_2 L^{k-2}(a) + \dots + c_1 L^{k-r+1}(a) + c_r L^{k-r}(a), c_i \in \mathbb{F}_2$$

\implies

$$a_k = c_1 a_{k-1} + c_2 a_{k-2} + \dots + c_{r-1} a_{k-r+1} + c_r a_{k-r}, k \geq r$$

then, $a \in G(f)$, and $f = 1 + c_1 x + c_2 x^2 + \dots + c_r x^r$

$$L^k(a) \in G(f), \forall k \geq 0, \text{ thus } V \subset G(f), |G(f)| = 2^r \geq 2^n \implies r \geq n$$

Hence $r=n$, i.e. $\deg(f) = n, p(a) = 2^n - 1$ so a is a m -sequence with $p(a) = 2^n - 1$. □

3.2 Polynomials of m -sequences

Theorem 3.2.1. $f(x) \in \mathbb{F}_q[x], f(x) = 1 + c_1 x + c_2 x^2 + \dots + c_n x^n (c_n \neq 0), a(\neq 0) \in G(f)$
 a is a m -sequence $\implies f(x)$ is irreducible.

In fact, we have more strong theorem.

Theorem 3.2.2. $f(x) \in \mathbb{F}_q[x], f(x) = 1 + c_1 x + c_2 x^2 + \dots + c_n x^n (c_n \neq 0), a(\neq 0) \in G(f)$
 a is a m -sequence $\iff f(x)$ is a primitive polynomial.

Definition 3.2.3. a is a periodic sequence \mathbb{F}_q . $a = (a_0, a_1, a_2, \dots)$. $s \in \mathbb{Z}^+$,

$$a^{(s)} \triangleq (a_0, a_s, a_{2s}, \dots)$$

is a sampling sequence of a with s as its period.

Lemma 3.2.4. a is a periodic sequence \mathbb{F}_q , $s \equiv s_1 \pmod{p(a)}$, then

$$a^{(s)} = a^{(s_1)}$$

Lemma 3.2.5. a is a periodic sequence \mathbb{F}_q , $s \in \mathbb{Z}^+$, $a^{(s)}$ is periodic, then

$$p(a^{(s)}) \mid \frac{p(a)}{\gcd(s, p(a))}$$

Lemma 3.2.6. a is a periodic sequence \mathbb{F}_q , $s \in \mathbb{Z}^+$, $\gcd(s, p(a)) = 1$, then $a^{(s)}$ is also periodic,

$$p(a^{(s)}) = p(a)$$

Definition 3.2.7. (the Trace Representation of m -sequence) $Tr_{\mathbb{F}_{q^n}/\mathbb{F}_q}(\xi) = \xi + \xi^q + \xi^{q^2} + \dots + \xi^{q^{n-1}}$ is the trace of ξ .

Lemma 3.2.8. $\forall \xi \in \mathbb{F}_{p^n}, Tr(\xi) \in \mathbb{F}_q$

3.3 Trace representations

Theorem 3.3.1. a is m -sequence over \mathbb{F}_q , $a = (a_0, a_1, a_2, \dots), p(a) = q^n - 1$. The minimal polynomial of a , $f(x) = 1 + c_1 x + c_2 x^2 + \dots + c_n x^n$, is a primitive polynomial. α is a root of $\tilde{f}(x)$, $\exists \beta \in \mathbb{F}_{q^n}^*$, s.t.

$$a = (Tr(\beta), Tr(\beta\alpha), Tr(\beta\alpha^2), \dots)$$

On the other hand, suppose $f(x) \in \mathbb{F}[x], f(x) = 1 + c_1 x + \dots + c_n x^n$ is primitive polynomial. Let α is a root of $\tilde{f}(x), \beta \in \mathbb{F}_{q^n}$,

$$(Tr(\beta), Tr(\beta\alpha), Tr(\beta\alpha^2), \dots) \in G(f)$$

is m -sequence.

Proof. $f(x) = 1 + c_1x + c_2x^2 + \cdots + c_nx^n$. $\tilde{f}(x) = x^n + c_1x^{n-1} + \cdots + c_n$.

$$\tilde{f}(\alpha) = \alpha^n + c_1\alpha^{n-1} + \cdots + c_n = 0 \implies \alpha^k + c_1\alpha^{k-1} + \cdots + c_n\alpha^{k-n}, k \geq n$$

Then $(1, \alpha, \alpha^2, \dots) \in G(f) \implies (\beta, \beta\alpha, \beta\alpha^2, \dots) \in G(f), \forall \beta \in \mathbb{F}_{q^n}$

For the same reason,

$$(\beta^{q^j}, (\beta\alpha)^{q^j}, \dots) \in G(f), \forall j, 0 \leq j \leq n-1$$

Add the corresponding items, then we have

$$(Tr(\beta), Tr(\beta\alpha), Tr(\beta\alpha^2), \dots) \in G(f)$$

It's not difficult to prove that, let $\beta = 0, 1, \dots, q^n - 1$, we can get q^n distinct sequences, which are all of the sequences in $G(f)$.

Hence, $\forall \alpha, \exists \beta, s.t. a = (Tr(\beta), Tr(\beta\alpha), Tr(\beta\alpha^2), \dots)$ \square

Theorem 3.3.2. a is m -sequence over $(F)_q$. $p(a) = q^n - 1, \gcd(s, q^n - 1) = 1$. $a^{(s)}$ is m -sequence, and $p(a^{(s)}) = q^n - 1$

Proof. Lemma 4.10 $\implies a^{(s)}$ is periodic, and $p(a) = q^n - 1$.

Assume the minimal polynomial of α is $f(x)$, $\forall \alpha$ is a root of $\tilde{f}(x)$, $\exists \beta \in \mathbb{F}_{q^n}^*, s.t.$

$$a = (Tr(\beta), Tr(\beta\alpha), Tr(\beta\alpha^2))$$

Thus,

$$a^{(s)} = (Tr(\beta), Tr(\beta\alpha^s), Tr(\beta\alpha^{2s}), \dots)$$

$$a^{(s)} = (Tr(\beta), Tr(\beta\alpha^s), Tr(\beta\alpha^{2s}), \dots)$$

$\tilde{f}(x)$ is a primitive polynomial, thus α is a primitive element. $\gcd(s, q^n - 1) = 1$, hence $\alpha^{(s)}$ is also a primitive element.

Assume the minimal polynomial of $\alpha^{(s)}$ is $\tilde{f}_s(x)$, which is a primitive polynomial. $f_x(s)$ is a primitive polynomial. According to **Theorem 4.5**, $\alpha^{(s)}$ is a m -sequence and $p(a) = q^n - 1$. \square

Corollary 3.3.3. a is a m -sequence in $\mathbb{F}_q, p(a) = q^n - 1, \forall m$ -sequence in \mathbb{F}_q is shifting equivalent to a sampling sequence of a .

Proof. Assume the min.poly. of a is $f(x)$, $\gamma \in \mathbb{F}_{q^n}$ is a root of $f(x)$, and b is a m -sequence in \mathbb{F} , $p(b) = q^n - 1$ and assume γ^s is a root of the minimal polynomial of b ,

$$\gcd(s, q^n - 1) = 1, \xrightarrow{Thm 4.14} a^{(s)} \text{ is a } m\text{-sequence}, p(a^{(s)}) = q^n - 1. \gamma^s \text{ is a root min.poly. of } a^{(s)}.$$

γ^s is a root of the minimal polynomial of b .

b and $a^{(s)}$ have the same minimal polynomial $\xrightarrow{Thm 4.2} b \sim a^{(s)}$. \square

Lemma 3.3.4. a is a m -sequence in $\mathbb{F}_q, p(a) = q^n - 1$. Any n successive states

$$S_m, S_{m+1}, S_{m+2}, \dots, S_{m+n-1}, m \geq 0$$

are linearly independent in \mathbb{F}_q

Chapter 4

The pseudo-randomness of m-sequence

4.1 The pseudo-randomness of m-sequence

In \mathbb{F}_2

Theorem 4.1.1. *a is a m-sequence in \mathbb{F}_2 , $p(a) = 2^n - 1$,*

Arrange a period of a in a circle successively. $0 < k \leq n$, $\forall k$ -tuple (b_1, b_2, \dots, b_k) in \mathbb{F}_2 appear

$$\begin{cases} 2^n - k & (b_1, b_2, \dots, b_k) \neq (0, 0, \dots, 0) \\ 2^{n-k} - 1 & (b_1, b_2, \dots, b_k) = (0, 0, \dots, 0) \end{cases}$$

time in the circle.

Corollary 4.1.2. *a is a m-sequence in \mathbb{F}_2 . $p(a) = 2^n - 1$, 1 appear 2^{n-1} times, 0 appear $2^{n-1} - 1$*

Definition 4.1.3. *η is a group isomorphism from additive group of \mathbb{F}_2 to $\{+1, -1\}$ which is a multiplicative group.*

$$\eta: \mathbb{F}_2 \longrightarrow \{+1, -1\}$$

$$0 \longmapsto +1$$

$$1 \longmapsto -1$$

a is a periodic sequence in \mathbb{F}_2 , the autocorrelation function of a:

$$C_a(t) \triangleq \sum_{i=0}^{p(a)-1} \eta(a_i) \eta(a_{i+t})$$

Note:

$$(1) C_a(t) = p(a), t | p(a).$$

$$(2) C_a(p(a) + k) = C_a(k), \forall k \geq 0$$

$$(3) C_a(t) = \sum_{i=0}^{p(a)+1} \eta(a_i + a_{i+t})$$

Theorem 4.1.4. *a is a m-sequence in \mathbb{F}_2 . $p(a) = 2^n - 1$, $C_a(t) = -1$ if $t \nmid 2^n - 1$*

Proof. $a = (a_0, a_1, a_2, \dots)$, assume the min.poly. of a is $f(x)$.

$$f(x) = c_0 + c_1x + c_2x^2 + \dots + c_nx^n, \quad c_0, c_n \neq 0$$

a satisfies with:

$$c_0a_k + c_1a_{k-1} + c_2a_{k-2} + \dots + c_na_{k-n}, k \geq 0(*)$$

$L^t(a) = (a_t, a_{t+1}, a_{t+2}, \dots)$, $a + L^t(a)$ also satisfies with (*)

$$t \nmid 2^n - 1, a + L^t(a) \neq 0$$

according to Thm4.2 $a + L^t(a)$ is m-sequence.

$$C_a(t) = \sum_{i=0}^{p(a)-1} \eta(a_i) \eta(a_{i+t}) = \sum_{i=0}^{p(a)-1} \eta(a_i + a_{i+t}) = 2^{n-1} \dots (-1) + (2^{n-1} - 1) = -1$$

□

Definition 4.1.5. a is a periodical sequence. If $C_a(t) = -1, t \nmid p(a)$, a is a pseudo-random-sequence.

Theorem 4.1.6. m -sequence is a pseudo-random-sequence.

Chapter 5

NLFSR

5.1 Basic Concepts

In the section, we are going to study a bit of NLFSR(Nonlinear feedback shift register).

Definition 5.1.1. *The Initial state of FSR(Feedback shift register) is $(a_0, a_1, \dots, a_{n-1})$. Its Feedback function is $f(x_1, x_2, \dots, x_n)$.*

*If $f(0, x_2, \dots, x_n) = f(1, x_2, \dots, x_n)$, then we say f is **degenerated**, otherwise it is **non-degenerated**. If f is degenerated, let $g(x_1, x_2, \dots, x_n) = f(0, x_1, x_2, \dots, x_{n-1})$.*

When it comes to FSR, we always assume it is non-degenerated.

$V(\mathbb{F}_n)$ is a set that contains all states generated by $f(x)$. $|V(\mathbb{F}_n)| = 2^n$.

*We often use **state diagram** G_f to study FSR sequence. G_f is a directed graph with 2^n vertices and 2^n arcs. It determines **state-transition transition** T_f from $V(\mathbb{F}_n)$ to $V(\mathbb{F}_n)$. T_f is defined as follows.*

$$\begin{aligned} T_f : V(\mathbb{F}_n) &\rightarrow V(\mathbb{F}_n) \\ (a_1, a_2, \dots, a_n) &\mapsto (a_2, a_3, \dots, a_n, f(a_1, a_2, \dots, a_n)) \end{aligned}$$

Theorem 5.1.2. *The state diagrams of n -level FSR must have cycles.*

Proof. Initial state $S_0 = (a_0, a_1, \dots, a_{n-1})$. Let $S_k = T_f S_{k-1} = (a_k, a_{k+1}, a_{k+n-1})$. $k = 1, 2, \dots$. Notice that $|V(\mathbb{F}_n)| = 2^n$. Thus in S_0, S_1, \dots, S_{2^n} , there must be at least two identical states.

Let $n_1 \leq 2^n$ be the minimal number satisfying that $\exists n_0 < n_1$, s.t. $S_{n_1} = S_{n_0}$. Thus S_{n_0} is unique. Then $S_0, S_1, \dots, S_{n_0}, S_{n_0+1}, \dots, S_{n_1-1}$ are distinct. Hence, $S_{n_0}, S_{n_0+1}, \dots, S_{n_1-1}$ form a cycle. \square

Note: The length of the cycle is often called the period of the cycle. $S_0, S_1, \dots, S_{n_0-1}$ are in a branch of the cycle.

Corollary 5.1.3. *\forall n -level FSR sequence, \exists a non-negative integer n_0 , s.t. $a_{n_0}, a_{n_0+1}, \dots$ is a periodic sequence. Its period $\leq 2^n$, which is independent of n_0 .*

5.2 State Diagrams of FSRs

Now let's study the state diagram of FSR. A diagram without branches has more elegant and simple structure that composes of only several cycles. As you probably have found out, a problem arises. In what circumstances it has no branches? The following are some sufficient and necessary conditions that ensure no branches in the diagram.

Theorem 5.2.1. *The state-diagram of n -level FSR has no branches if and only if the diagram is composed of cycles which has no common vertices.*

Theorem 5.2.2. *The state-diagram of n -level FSR whose feedback function is $f(x_1, x_2, \dots, x_n)$ has no branches if and only if T_f is bijective.*

Proof. If G_f has no branches, then it is composed of cycles which have no common vertices. Thus T_f is bijective.

Suppose T_f is bijective. Initial state $S_0, S_k = T_f S_{k-1}, k = 1, 2, \dots$.

Let n_0, n_1 satisfying the conditions in **Theorem 5.1.2**. Then $S_0, S_1, S_2, \dots, S_{n_0}, \dots, S_{n_1-1}$ are distinct.

If $n_0 \neq 0$, then $T_f(S_{n_0-1} = S_0, T_f(S_{n_1-1} = S_{n_1} = S_0$. Notice that $S_{n_0-1} \neq S_{n_1-1}$, which contradicts the fact that T_f is bijective.

Thus $n_0 = 0$, which means that G_f has no branches. \square

5.3 Main theorems

Theorem 5.3.1. *The state-diagram of n -level FSR whose feedback function is $f(x_1, x_2, \dots, x_n)$ has no branches if and only if f can be expressed as $f(x_1, x_2, \dots, x_n) = x_1 + f_0(x_2, x_3, \dots, x_n)$.*

Proof. \forall State $(a_1, a_2, \dots, a_n) \neq (b_1, b_2, \dots, b_n)$,

$$T_f(a_1, a_2, \dots, a_n) = (a_2, \dots, a_n, f(a_1, \dots, a_n))$$

$$T_f(b_1, b_2, \dots, b_n) = (b_2, \dots, b_n, f(b_1, \dots, b_n))$$

If $a_1 = b_1, (a_2, \dots, a_n) \neq (b_2, \dots, b_n)$, then

$$T_f(a_1, \dots, a_n) \neq T_f(b_1, \dots, b_n)$$

If $a_1 \neq b_1, i.e. b_1 = \bar{a}_1, (a_2, \dots, a_n) = (b_2, \dots, b_n)$, then

$$f(\bar{a}_1, a_2, \dots, a_n) = a_1 + 1 + f_0(a_2, \dots, a_n) = 1 + f(a_1, a_2, \dots, a_n)$$

Then

$$\begin{aligned} T_f(a_1, a_2, \dots, a_n) &= (a_2, \dots, a_n, f(a_1, a_2, \dots, a_n)) \\ &\neq (a_2, \dots, a_n, f(\bar{a}_1, a_2, \dots, a_n)) \\ &= T_f(\bar{a}_1, a_2, \dots, a_n) \end{aligned}$$

In conclusion, T_f is bijective $\rightarrow G_f$ has no branches.

On the other hand, suppose

$$f(x_1, x_2, \dots, x_n) = x_1 f_1(x_2, \dots, x_n) + f_0(x_2, \dots, x_n)$$

If $f_1(x_2, \dots, x_n) \neq 0$, then $\exists (a_2, \dots, a_n) \in \mathbb{F}_{2^n-1}$, s.t. $f(a_2, \dots, a_n) = 0$.

$\forall a_1 \in \mathbb{F}_2$,

$$f(a_1, a_2, \dots, a_n) = f_0(a_2, \dots, a_n) = f(\bar{a}_1, a_2, \dots, a_n)$$

$$T_f(a_1, a_2, \dots, a_n) = T_f(\bar{a}_1, a_2, \dots, a_n)$$

Then T_f is not bijective. G_f has no branches. \square

Note: we call feedback function $f(x_1, x_2, \dots, x_n)$ non-singular. If $f(x_1, x_2, \dots, x_n)$ can be expressed as $f(x_1, x_2, \dots, x_n) = x_1 + f_0(x_2, \dots, x_n)$

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