1. A Stack Walks into a Bar and Says "It's Hard to Maintain Discipline While Getting Smashed" (20 points): Consider the following C datatype, intended to provide some protection against buffer overflow:

```
struct safe buffer {
  int size;
                           12 buffer, then 4 from
  char * buffer;
                                                    higher address to lower
                           size
} mybuf;
And the following function that creates a safe buffer:
                                                                                arg 1
void createbuf(struct safe buffer *buf, int size) {
  int i;
                                                             createbuf
  char tempbuf[12];
                                                              24 bytes
  (*buf).size=size;
  (*buf).buffer=calloc((*buf).size, sizeof(char));
  gets(tempbuf);
  for (i=0; i<size; i++)
    (*buf).buffer[i]=tempbuf[i];
  return;
}
```

Your friend argues that this function takes a size parameter and allocates that much buffer space – then only writes that much space to the buffer, ensuring that the safe buffer cannot overflow. Your friend is completely wrong. Prove that the code can be exploited – give us a string (plain text) that could be used to form an exploit string as you did in the buffer lab. You exploit string should maintain the correct value of the saved ebp on the stack, but should change the saved return address to 0x080485e8 so that the return from createbuf will take us to that address. Don't worry about the call to hex2string – just give us plain text. Here's some useful data from execution on IA32 Linux – the disassembled call to createbuf:

```
8048512: e8 e2 fe ff ff call 80483f9 <createbuf>
```

And the values of %esp and %ebp, and a gdb dump of some of the stack, after the call to gets in createbuf has completed and we are inside the for loop in createbuf:

```
0xffffdb40
esp
               0xffffdb58
ebp
                                              this is where the
(gdb) x/32x 0xffffdb40
                                             buffer starts
               0xffffdb48
0xffffdb40:
                              0x0000001
                                             0x74617257
                                                            0x0000068
               0x00000000
                                            0xffffdbc8
0xffffdb50:
                              0x0000000
                                                            0x08048517
0xffffdb60:
               0x08049720
                              0x0000000a
                                             0x00000000
                                                            0x0000000
0xffffdb70:
               0xf63d4e2e
                              0x0000000
                                             0x00000000
                                                            0x0000000
0xffffdb80:
               0x00000000
                              0x0000000
                                             0x00000000
                                                           0x08048340
0xffffdb90:
               0x00000000
                              0x080496f4
                                             0xffffdba8
                                                           0x080482a1
0xffffdba0:
               0x00299ff4
                              0x00298204
                                             0xffffdbd8
                                                           0x08048569
0xffffdbb0:
               0x00183e25
                              0xffffdc6c
                                             0xffffdbd8
                                                           0x00299ff4
```

2. Cache Me If You Can! (30 points): Consider the following C function.

```
void shrink(int *old, int *new, int dim new, int shrink factor) {
     int i, j;
     int u, v;
i = 0,1 for (i=0; i<dim new; i++) {
  j=0,1 for (j=0; j<dim new; j++) {
    0.1.2.3 \text{ new}[i*dim new+j]=0;
         for (u=0; u<shrink factor; u++) {</pre>
           for (v=0; v<shrink factor; v++) {</pre>
              new[i*dim_new+j]\overline{+}= 0,2,1,3 4 = 0,8,4,12
                   old[(i*shrink factor+u)*dim new*shrink factor
                         +(j*shrink factor+v)];
                                                    old[0,2,1,3,8,10,9,11,4,6,5,7,12,14,1
                                                    3.151 =
         }
         new[i*dim new+j]/=shrink factor*shrink factor;
                              4
     }
  }
```

This function effectively takes a 2D matrix (int *old) and outputs a new 2D matrix based on this called (int *new). The parameter dim_new defines the size of the new 2D matrix – it is effectively a matrix of (dim_new *dim_new) integers. The last parameter, shrink_factor, is how many times smaller one dimension of the new matrix is relative to the old matrix. So if we went from a 400x400 matrix to a 100x100 matrix, dim_new would be 100 and shrink_factor would be 4. This could be used for something like image scaling. The technique to shrink the matrix will basically just use a simple, non-overlapping average – probably not good enough for high quality image scaling, but we'll do the best we can with it.

This problem is intended to be the most challenging one on this exam – so before continuing be sure you understand the original code first – it may be useful to run through an example of the shrinking on a small matrix – like a 4x4 matrix shrinking to a 2x2 matrix (scaling factor is 2).

We want to optimize this code by using strength reduction and common subexpression elimination on as many multiplies as possible, by eliminating unneeded memory references, and by using blocking to improve locality in the loop structure. There are lots of ways to attack this, but we are going to force you to finish the one we have started on the next page (this one cuts the runtime of shrink in half). The author of this code segment has followed a *horrible* coding practice of naming some of their variable names in a completely irrelevant way to the code function – so you cannot rely on the variable names to help you discern their functionality.

Your job is to fill in the blanks to make this code work correctly. The blanks we have inserted will look like this: $\underline{\mathbf{A}}$ where the letter at the center of the blank is the label for the space on the answer key. So you should have 5 labels (\mathbf{A} - \mathbf{E}) to fill in for this problem. MIN(X,Y) is a macro that returns the minimum of values X and Y.

```
void shrink fast(int *old, int *new, int dim new, int shrink factor)
  int i, j;
  int u, v;
  int iidim, jj, ii;
  // HINT - all labels should be filled with one of these names
  int platypus, kangaroo, echidna, cassowary, koala, dingo, wallaby,
      wombat;
  int dimshrink, sf2, sf2dim, bdim;
  dimshrink=dim new*shrink factor;
  sf2=shrink factor*shrink factor;
8 sf2dim=sf2*dim new;
20 bdim=BSIZE*dim new;
  iidim=0;
  for (ii=0; ii<dim new; ii+=BSIZE)<sub>4</sub> {
    for (jj=0; jj<dim new; jj+=BSIZE) {
      wallaby: MIN(ii+BSIZE, dim new);
      wombat = iidim;
   4x pracypus=iidim*sf2;
  0.2 cassowary=jj*shrink factor;
      for (i = ii; i < wallaby A; i++) {
         kangaroo: cassowary +sf2dim;
       aingo=MIN(jj+BSIZE,dim new);
    0.2 echidna=cassowary;
         for (j = jj; j < dingo C ; j++) {
          koala=0;
           for (u=platypus; u<kangaroo; u+=dimshrink) {</pre>
             for (v=echidna; v<echidna+shrink factor; v++) {</pre>
                    D +=old[u+v];
           new[wombatE____+j]=koala/sf2;
           echidna+=shrink factor;
        wombat+=dim new;
        platypus+=sf2dim;
    iidim+=bdim;
}
```