

Machine-Level Programming: Introduction

Chapter 3 of B&O

ISA - operators and stuff machine needs
Micro arch

Language -> IR -> Assem -> Machine Code -> [ISA][Micro arch]

Some notes adopted from Bryant and O'Hallaron

IA32 Processors

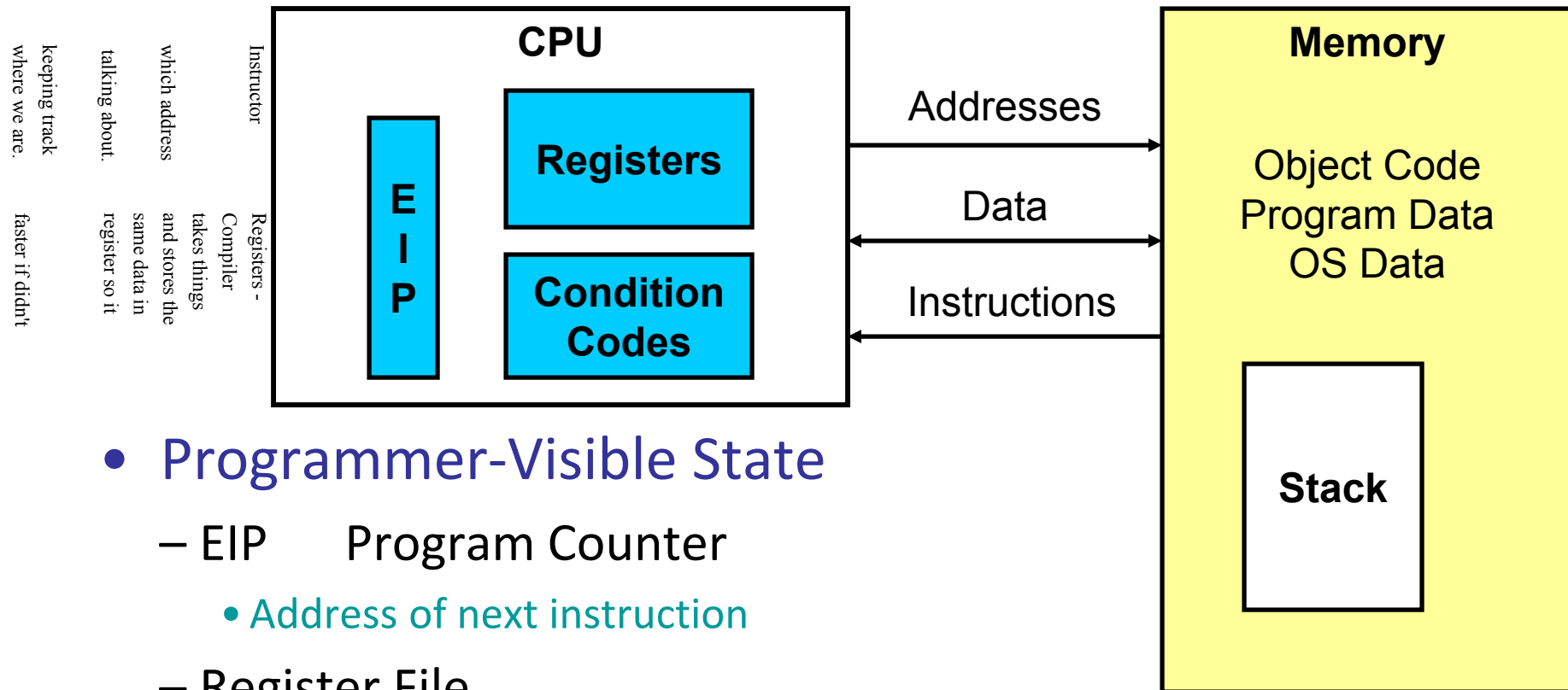
- Evolutionary Design

- Starting in 1978 with 8086
- Added more features as time goes on
- Still support old features, although obsolete

- Complex Instruction Set Computer (CISC)

- Many different instructions with many different formats
 - But, only small subset encountered with Linux programs
- Hard to match performance of Reduced Instruction Set Computers (RISC)
Low ops, Simple, Takes up more space (breaks up into pieces and can put in parallel)
- But, Intel has done just that!

Assembly Programmer's View



- **Programmer-Visible State**

- **EIP** Program Counter

- Address of next instruction

- **Register File**

- Heavily used program data

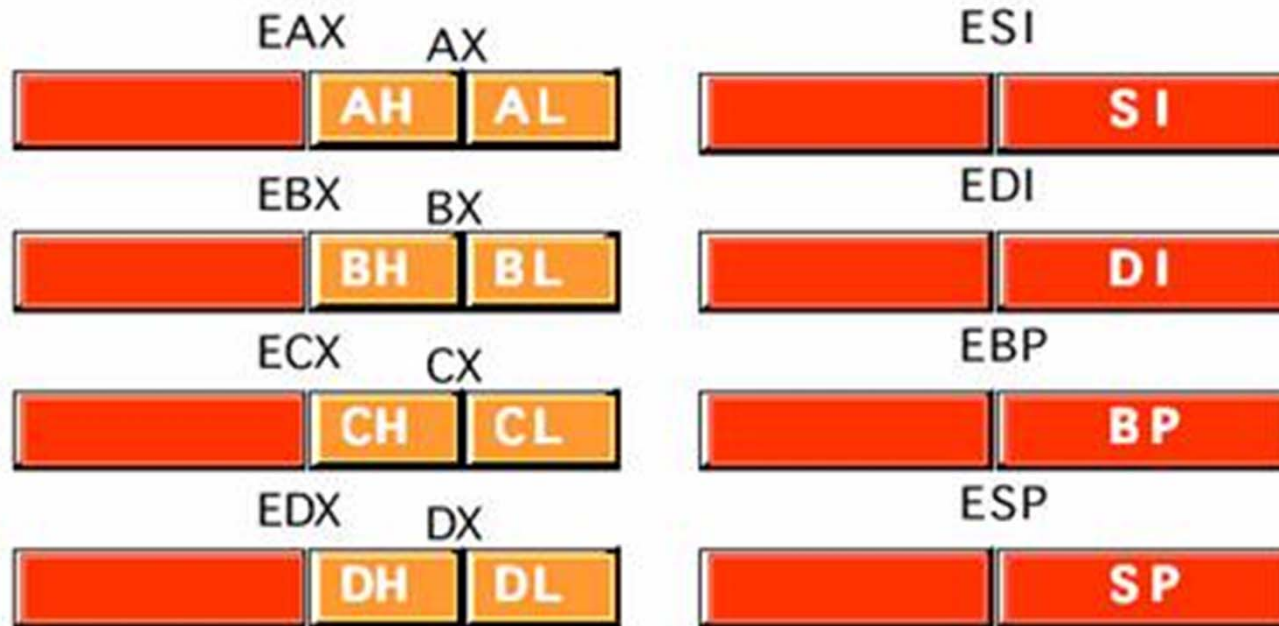
- **Condition Codes**

- Store status information about most recent arithmetic operation
 - Used for conditional branching

- **Memory**

- Byte addressable array
 - Code, user data, (some) OS data
 - Includes stack used to support procedures

IA32 Registers



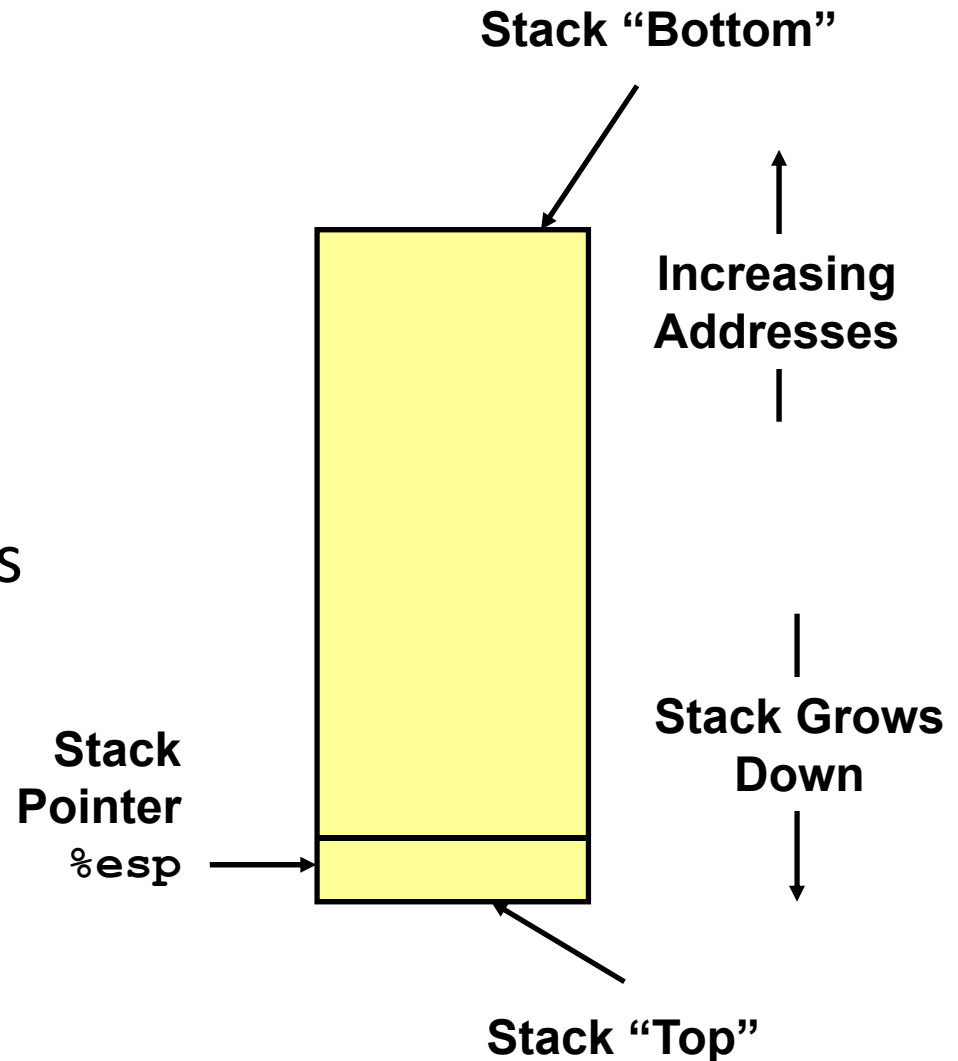
Integer registers

E P -- 32

IA32 Stack

- Region of memory managed with stack discipline
- Grows toward lower addresses
- Register `%esp` indicates lowest stack address
 - address of top element

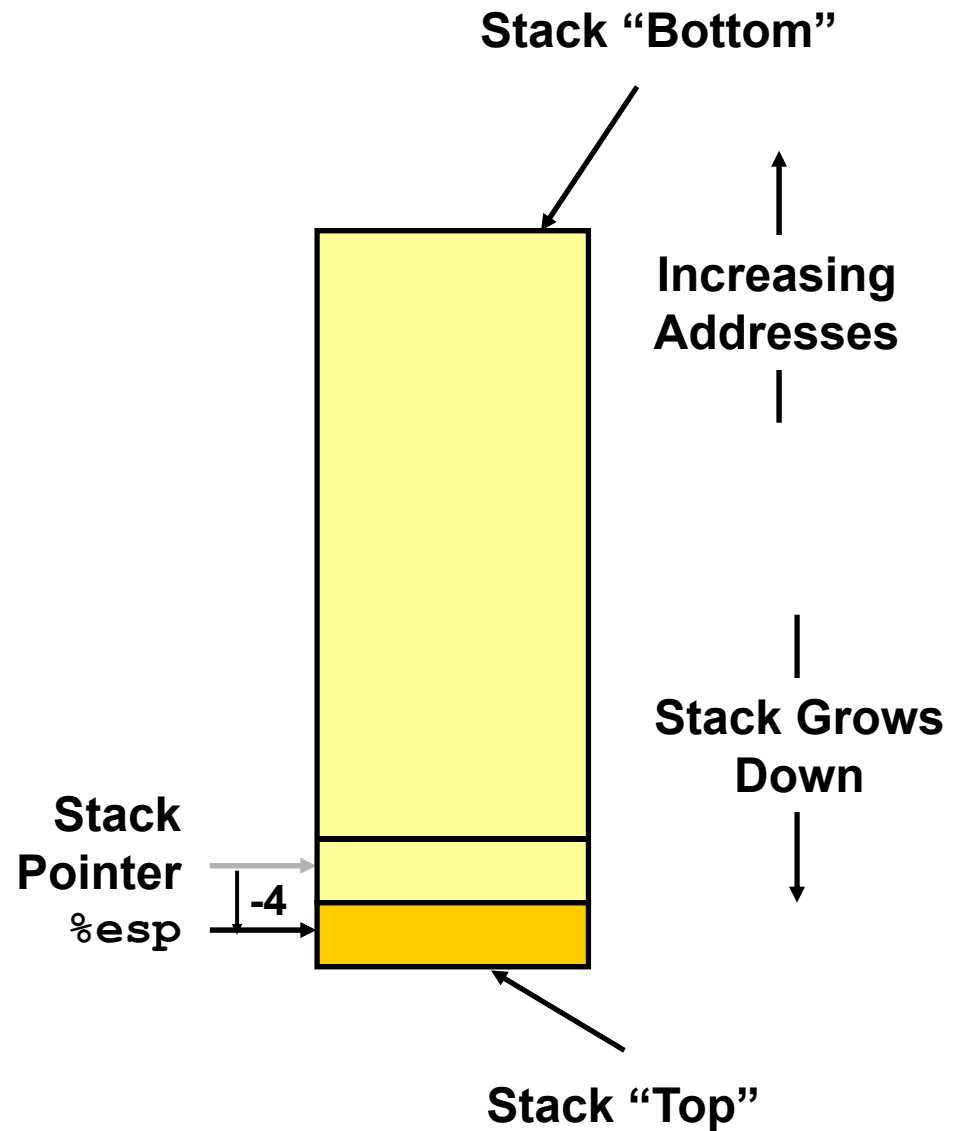
Register is not in memory, it is in a separate



IA32 Stack Pushing

- Pushing

- `pushl Src`
- Fetch operand at *Src*
- Decrement `%esp` by 4
- Write operand at address given by `%esp`

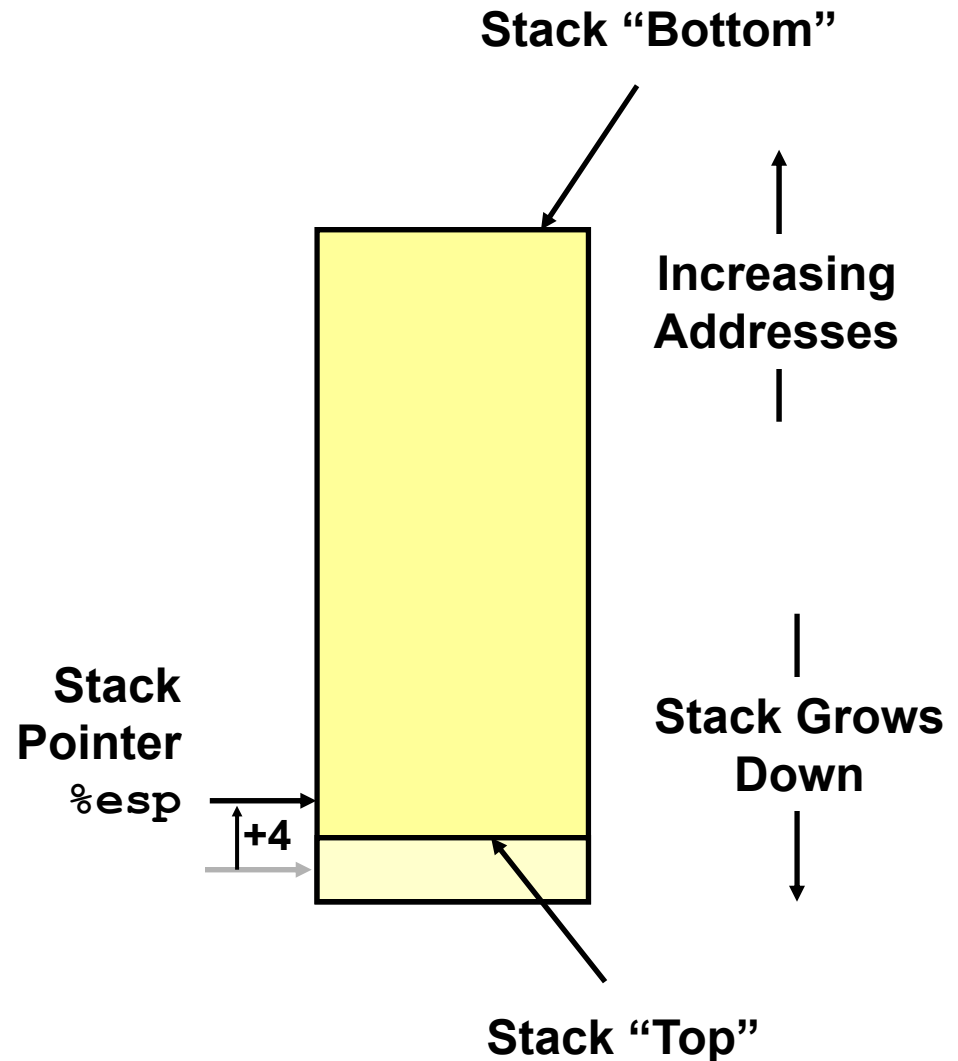


IA32 Stack Popping

- Popping

- `popl Dest`
- Read operand at address given by `%esp`
- Increment `%esp` by 4
- Write to `Dest`

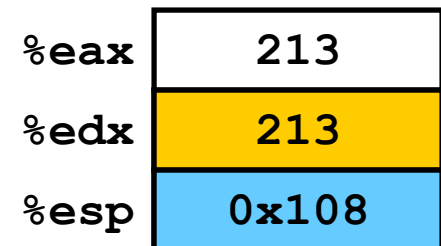
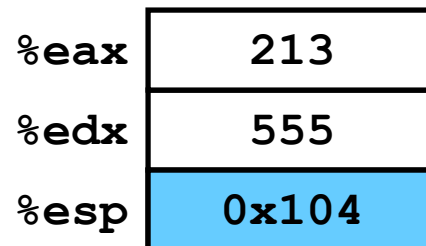
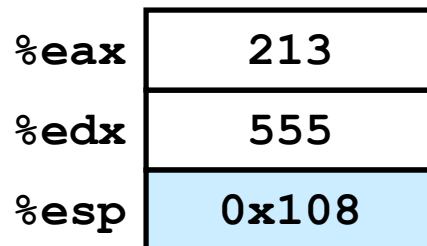
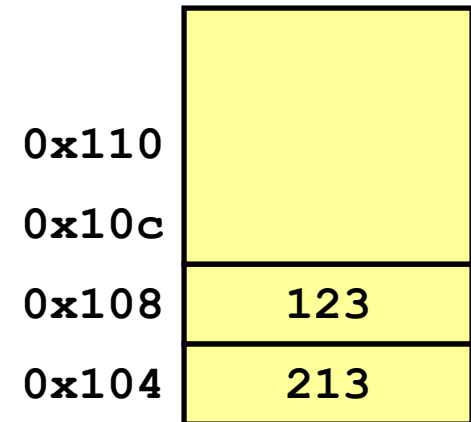
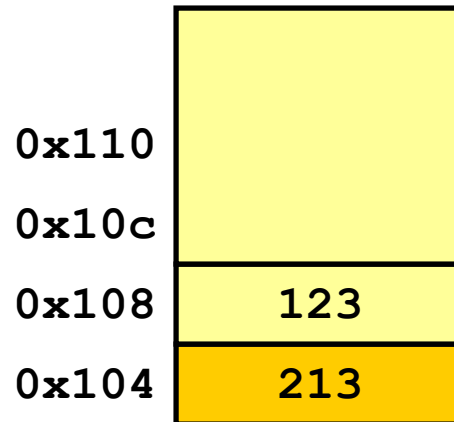
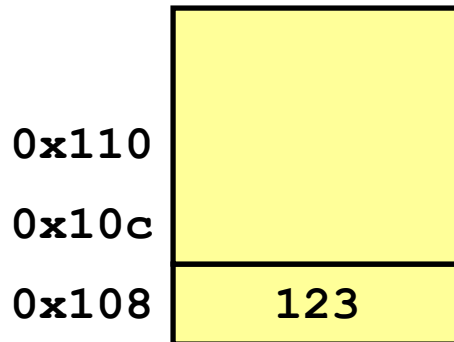
If you store 512 in the memory, and you pop, you don't erase that 512. You just change the esp pointer



Stack Operation Examples

`pushl %eax`

`popl %edx`



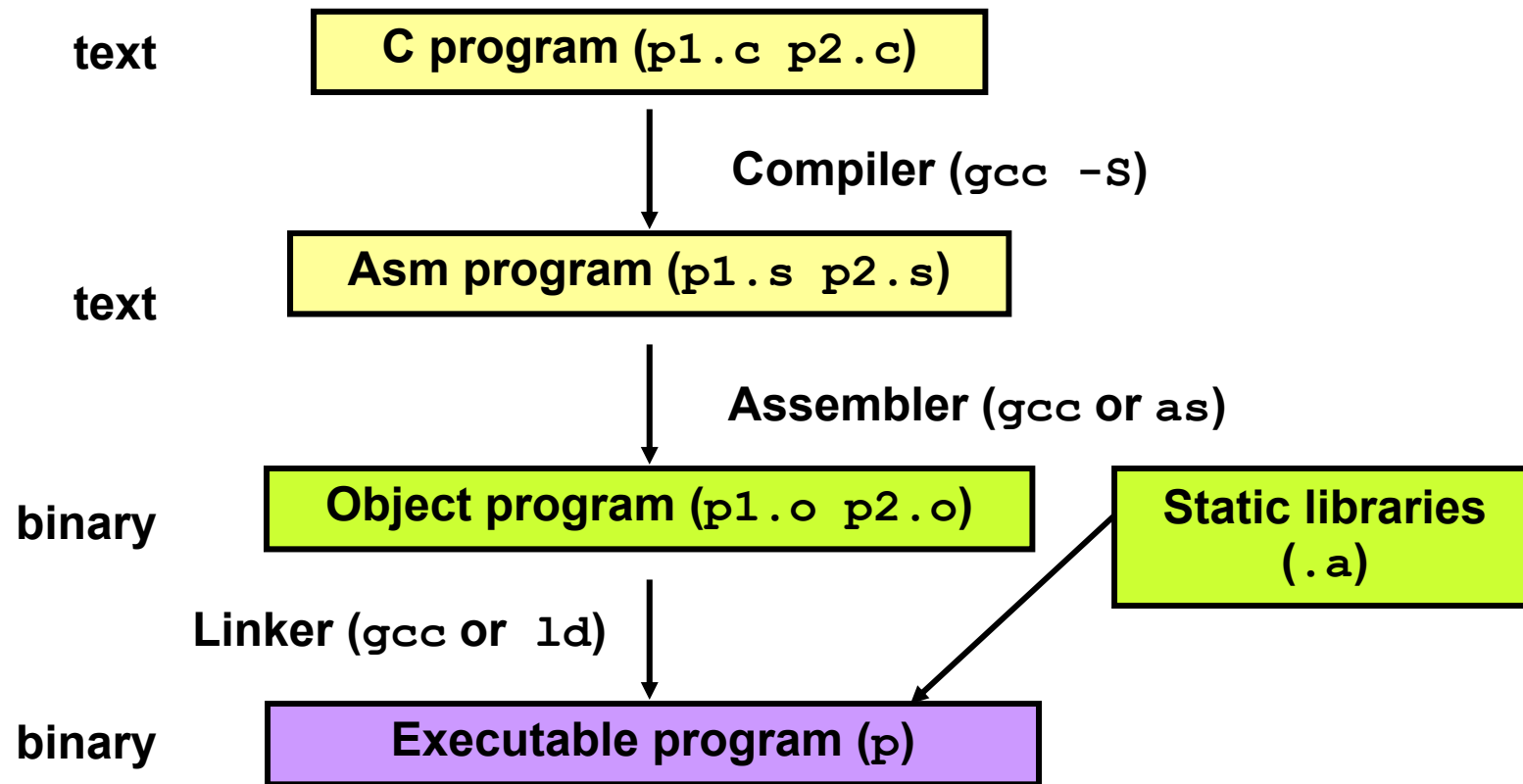
Turning C into Object Code

- Code in files `p1.c p2.c`

- Compile with command:

`gcc -O`

`p1.c p2.c -o p`



Compiling Into Assembly

C Code

```
int sum(int x, int y)
{
    int t = x+y;
    return t;
}
```

Generated Assembly

```
_sum:
    pushl %ebp
    movl %esp,%ebp
    movl 12(%ebp),%eax
    addl 8(%ebp),%eax
    movl %ebp,%esp
    popl %ebp
    ret
```

Obtain with command

```
gcc -O -S code.c
```

Produces file `code.s`

Assembly Characteristics

- Minimal Data Types

- “Integer” data of 1, 2, or 4 bytes
 - Data values
 - Addresses (untyped pointers)
- Floating point data of 4, 8, or 10 bytes
- No aggregate types such as arrays or structures
 - Just contiguously allocated bytes in memory

- Primitive Operations

- Perform arithmetic function on register or memory data
- Transfer data between memory and register
 - Load data from memory into register
 - Store register data into memory
- Transfer control
 - Unconditional jumps to/from procedures
 - Conditional branches

Object Code

Code for sum

0x401040 <sum>:

0x55

0x89

0xe5

0x8b

0x45

0x0c

0x03

0x45

0x08

0x89

0xec

0x5d

0xc3

- Total of 13 bytes
- Each instruction 1, 2, or 3 bytes
- Starts at address 0x401040

- **Assembler**

- Translates .s into .o
- Binary encoding of each instruction
- Nearly-complete image of executable code
- Missing linkages between code in different files

- **Linker**

- Resolves references between files
- Combines with static run-time libraries
 - E.g., code for `malloc`, `printf`
- Some libraries are *dynamically linked*
 - Linking occurs when program begins execution

Machine Instruction Example

```
int t = x+y;
```

```
addl 8(%ebp), %eax
```

Similar to
expression

x += y

x =
eax

```
0x401046:    03 45 08
```

- C Code

- Add two signed integers

- Assembly

- Add two 4-byte integers

- “Long” words in GCC parlance
- Same instruction whether signed or unsigned

- Operands:

x: Register %eax

y: Memory M[%ebp+8]

t: Register %eax

- Return function value in %eax

- Object Code

- 3-byte instruction
- Stored at address 0x401046

Disassembling Object Code

Disassembled

```
00401040 <_sum>:
  0:      55          push    %ebp
  1:      89 e5      mov     %esp, %ebp
  3:      8b 45 0c    mov     0xc(%ebp), %eax
  6:      03 45 08    add     0x8(%ebp), %eax
  9:      89 ec      mov     %ebp, %esp
  b:      5d        pop     %ebp
  c:      c3        ret
  d:      8d 76 00   lea     0x0(%esi), %esi
```

- Disassembler

`objdump -d p`

- Useful tool for examining object code
- Analyzes bit pattern of series of instructions
- Produces approximate rendition of assembly code
- Can be run on either a `.out` (complete executable) or `.o` file

Alternate Disassembly

Object

0x401040:

0x55

0x89

0xe5

0x8b

0x45

0x0c

0x03

0x45

0x08

0x89

0xec

0x5d

0xc3

Disassembled

0x401040	<sum>:	push	%ebp
0x401041	<sum+1>:	mov	%esp, %ebp
0x401043	<sum+3>:	mov	0xc(%ebp), %eax
0x401046	<sum+6>:	add	0x8(%ebp), %eax
0x401049	<sum+9>:	mov	%ebp, %esp
0x40104b	<sum+11>:	pop	%ebp
0x40104c	<sum+12>:	ret	
0x40104d	<sum+13>:	lea	0x0(%esi), %esi

- Within gdb Debugger

`gdb p`

`disassemble sum`

– Disassemble procedure

`x/13b sum`

– Examine the 13 bytes starting at `sum`

What Can be Disassembled?

```
% objdump -d WINWORD.EXE
```

```
WINWORD.EXE:      file format pei-i386
```

```
No symbols in "WINWORD.EXE".
```

```
Disassembly of section .text:
```

```
30001000 <.text>:
```

30001000:	55	push	%ebp
30001001:	8b ec	mov	%esp, %ebp
30001003:	6a ff	push	\$0xffffffff
30001005:	68 90 10 00 30	push	\$0x30001090
3000100a:	68 91 dc 4c 30	push	\$0x304cdc91

- Anything that can be interpreted as executable code
- Disassembler examines bytes and reconstructs assembly source

Moving Data

- Moving Data

`movl Source, Dest:`

- Move 4-byte (“long”) word
- Lots of these in typical code

- Operand Types

- Immediate: Constant integer data
 - Like C constant, but prefixed with ‘\$’
 - E.g., `$0x400`, `$-533`
 - Encoded with 1, 2, or 4 bytes
- Register: One of 8 integer registers
 - But `%esp` and `%ebp` reserved for special use
 - Others have special uses for particular instructions
- Memory: 4 consecutive bytes of memory
 - Various “address modes”

<code>%eax</code>
<code>%edx</code>
<code>%ecx</code>
<code>%ebx</code>
<code>%esi</code>
<code>%edi</code>
<code>%esp</code>
<code>%ebp</code>

movl Operand Combinations

	Source	Destination		C Analog
movl	Imm	Reg	movl \$0x4,%eax	temp = 0x4;
		Mem	movl \$-147, (%eax)	*p = -147;
	Reg	Reg	movl %eax,%edx	temp2 = temp1;
		Mem	movl %eax, (%edx)	*p = temp;
	Mem	Reg	movl (%eax), %edx	temp = *p;

- Cannot do memory-memory transfers with single instruction

Simple Addressing Modes

- Normal (R) Mem[Reg[R]]

- Register R specifies memory address

```
movl (%ecx), %eax
```

- Displacement D(R)
Mem[Reg[R]+D]

- Register R specifies start of memory region

- Constant displacement D specifies offset

```
movl 8(%ebp), %edx
```

additio
n add
the
outside
to the
inside
this
would
be
useful
for an
array

Using Simple Addressing Modes

```
void swap(int *xp, int *yp)
{
    int t0 = *xp;
    int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

swap:

```
    pushl %ebp
    movl  %esp,%ebp
    pushl %ebx
```

} Set Up

```
    movl 12(%ebp),%ecx
    movl 8(%ebp),%edx
    movl (%ecx),%eax
    movl (%edx),%ebx
    movl %eax, (%edx)
    movl %ebx, (%ecx)
```

} Body

```
    movl -4(%ebp),%ebx
    movl %ebp,%esp
    popl %ebp
    ret
```

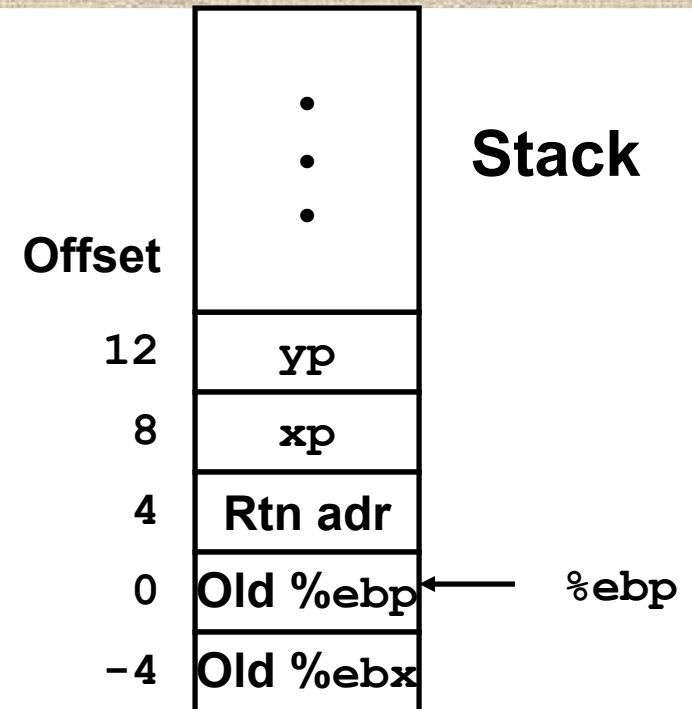
} Finish

Understanding Swap

```
void swap(int *xp, int *yp)
{
    int t0 = *xp;
    int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

Register	Variable
%ecx	yp
%edx	xp
%eax	t1
%ebx	t0

```
movl 12(%ebp), %ecx    # ecx = yp
movl 8(%ebp), %edx     # edx = xp
movl (%ecx), %eax      # eax = *yp (t1)
movl (%edx), %ebx      # ebx = *xp (t0)
movl %eax, (%edx)      # *xp = eax
movl %ebx, (%ecx)      # *yp = ebx
```



%eax	
%edx	
%ecx	
%ebx	
%esi	
%edi	
%esp	
%ebp	0x104

		Address	
		123	0x124
		456	0x120
			0x11c
			0x118
	Offset		0x114
yp	12	0x120	0x110
xp	8	0x124	0x10c
	4	Rtn adr	0x108
%ebp	→ 0		0x104
	-4		0x100

```
movl 12(%ebp), %ecx    # ecx = yp
movl 8(%ebp), %edx     # edx = xp
movl (%ecx), %eax      # eax = *yp (t1)
movl (%edx), %ebx      # ebx = *xp (t0)
movl %eax, (%edx)      # *xp = eax
movl %ebx, (%ecx)      # *yp = ebx
```

%eax	
%edx	
%ecx	0x120
%ebx	
%esi	
%edi	
%esp	
%ebp	0x104

		Address	
		123	0x124
		456	0x120
			0x11c
			0x118
	Offset		0x114
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	4	Rtn adr	0x108
%ebp	→ 0		0x104
	-4		0x100

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movl 12(%ebp), %ecx    # ecx = yp
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movl (%edx), %ebx      # ebx = *xp (t0)
movl %eax, (%edx)      # *xp = eax
movl %ebx, (%ecx)      # *yp = ebx
```

%eax	
%edx	0x124
%ecx	0x120
%ebx	
%esi	
%edi	
%esp	
%ebp	0x104

		Address
		123 0x124
		456 0x120
		0x11c
		0x118
Offset		0x114
yp	12	0x120 0x110
xp	8	0x124 0x10c
	4	Rtn adr 0x108
%ebp	0	0x104
	-4	0x100

```

movl 12(%ebp), %ecx    # ecx = yp
movl 8(%ebp), %edx    # edx = xp
movl (%ecx), %eax      # eax = *yp (t1)
movl (%edx), %ebx      # ebx = *xp (t0)
movl %eax, (%edx)      # *xp = eax
movl %ebx, (%ecx)      # *yp = ebx

```


%eax	456
%edx	0x124
%ecx	0x120
%ebx	
%esi	
%edi	
%esp	
%ebp	0x104

		Address
		123 0x124
		456 0x120
		0x11c
		0x118
Offset		0x114
yp	12	0x120 0x110
xp	8	0x124 0x10c
	4	Rtn adr 0x108
%ebp	→ 0	0x104
	-4	0x100

```

movl 12(%ebp), %ecx    # ecx = yp
movl 8(%ebp), %edx     # edx = xp
movl (%ecx), %eax      # eax = *yp (t1)
movl (%edx), %ebx      # ebx = *xp (t0)
movl %eax, (%edx)      # *xp = eax
movl %ebx, (%ecx)      # *yp = ebx

```

%eax	456
%edx	0x124
%ecx	0x120
%ebx	123
%esi	
%edi	
%esp	
%ebp	0x104

		Address	
		123	0x124
		456	0x120
			0x11c
			0x118
	Offset		0x114
yp	12	0x120	0x110
xp	8	0x124	0x10c
	4	Rtn adr	0x108
%ebp	→ 0		0x104
	-4		0x100

```
movl 12(%ebp), %ecx    # ecx = yp
movl 8(%ebp), %edx     # edx = xp
movl (%ecx), %eax      # eax = *yp (t1)
movl (%edx), %ebx      # ebx = *xp (t0)
movl %eax, (%edx)      # *xp = eax
movl %ebx, (%ecx)      # *yp = ebx
```

%eax	456
%edx	0x124
%ecx	0x120
%ebx	123
%esi	
%edi	
%esp	
%ebp	0x104

		Address
		456 0x124
		456 0x120
		0x11c
		0x118
Offset		0x114
yp	12	0x120 0x110
xp	8	0x124 0x10c
	4	Rtn adr 0x108
%ebp	0	0x104
	-4	0x100

```

movl 12(%ebp), %ecx    # ecx = yp
movl 8(%ebp), %edx     # edx = xp
movl (%ecx), %eax      # eax = *yp (t1)
movl (%edx), %ebx      # ebx = *xp (t0)
movl %eax, (%edx)    # *xp = eax
movl %ebx, (%ecx)      # *yp = ebx

```

%eax	456
%edx	0x124
%ecx	0x120
%ebx	123
%esi	
%edi	
%esp	
%ebp	0x104

		Address
		456 0x124
		123 0x120
		0x11c
		0x118
Offset		0x114
yp	12	0x120 0x110
xp	8	0x124 0x10c
	4	Rtn adr 0x108
%ebp	0	0x104
	-4	0x100

```

movl 12(%ebp), %ecx    # ecx = yp
movl 8(%ebp), %edx     # edx = xp
movl (%ecx), %eax      # eax = *yp (t1)
movl (%edx), %ebx      # ebx = *xp (t0)
movl %eax, (%edx)      # *xp = eax
movl %ebx, (%ecx)      # *yp = ebx

```

Indexed Addressing Modes

- Most General Form

D(Rb,Ri,S) Mem[Reg[Rb]+S*Reg[Ri]+ D]

- D: Constant “displacement” 1, 2, or 4 bytes
- Rb: Base register: Any of 8 integer registers
- Ri: Index register: Any, except for `%esp`
 - Unlikely you’d use `%ebp`, either
- S: Scale: 1, 2, 4, or 8

- Special Cases

(Rb,Ri) Mem[Reg[Rb]+Reg[Ri]]

D(Rb,Ri) Mem[Reg[Rb]+Reg[Ri]+D]

(Rb,Ri,S) Mem[Reg[Rb]+S*Reg[Ri]]

Address Computation Examples

<code>%edx</code>	<code>0xf000</code>
-------------------	---------------------

<code>%ecx</code>	<code>0x100</code>
-------------------	--------------------

Expression	Computation	Address
<code>0x8 (%edx)</code>	<code>0xf000 + 0x8</code>	<code>0xf008</code>
<code>(%edx, %ecx)</code>	<code>0xf000 + 0x100</code>	<code>0xf100</code>
<code>(%edx, %ecx, 4)</code>	<code>0xf000 + 4*0x100</code>	<code>0xf400</code>
<code>0x80(, %edx, 2)</code>	<code>2*0xf000 + 0x80</code>	<code>0x1e080</code>

Address Computation Instruction

- `leal Src, Dest`
 - *Src* is address mode expression
 - Set *Dest* to address denoted by expression
- Uses
 - Computing address without doing memory reference
 - E.g., translation of `p = &x[i];`
 - Computing arithmetic expressions of the form $x + k * y$
 - $k = 1, 2, 4, \text{ or } 8$.

18(%e ax) , %e bx = it is for ad ss cul ati on. It get s the me mo ry wit ho

Some Arithmetic Operations

Format

Computation

- Two Operand Instructions

<code>addl Src, Dest</code>	$Dest = Dest + Src$	
<code>subl Src, Dest</code>	$Dest = Dest - Src$	
<code>imull Src, Dest</code>	$Dest = Dest * Src$	
<code>sall Src, Dest</code>	$Dest = Dest \ll Src$	Also called <code>shll</code>
<code>sarl Src, Dest</code>	$Dest = Dest \gg Src$	Arithmetic
<code>shrl Src, Dest</code>	$Dest = Dest \gg Src$	Logical
<code>xorl Src, Dest</code>	$Dest = Dest \wedge Src$	
<code>andl Src, Dest</code>	$Dest = Dest \& Src$	
<code>orl Src, Dest</code>	$Dest = Dest Src$	

Some Arithmetic Operations

Format

Computation

- One Operand Instructions

`incl Dest` $Dest = Dest + 1$

`decl Dest` $Dest = Dest - 1$

`negl Dest` $Dest = - Dest$

`notl Dest` $Dest = \sim Dest$

More Arithmetic Operations

Format

Computation

`imull Source` $R[\%edx]:R[\%eax] = \text{Source} \times R[\%eax]$
(signed)

`mul Source` $R[\%edx]:R[\%eax] = \text{Source} \times R[\%eax]$
(unsigned)

`cld` $R[\%edx]:R[\%eax] = \text{SignExtend}(R[\%eax])$

take
es
a 3
bit
value
and
sign
it

Using leal for Arithmetic Expressions

```
int arith
(int x, int y, int z)
{
    int t1 = x+y;
    int t2 = z+t1;
    int t3 = x+4;
    int t4 = y * 48;
    int t5 = t3 + t4;
    int rval = t2 * t5;
    return rval;
}
```

arith:

```
pushl %ebp
movl %esp,%ebp
```

} Set
Up

```
movl 8(%ebp),%eax
movl 12(%ebp),%edx
leal (%edx,%eax),%ecx
leal (%edx,%edx,2),%edx
sall $4,%edx
addl 16(%ebp),%ecx
leal 4(%edx,%eax),%eax
imull %ecx,%eax
```

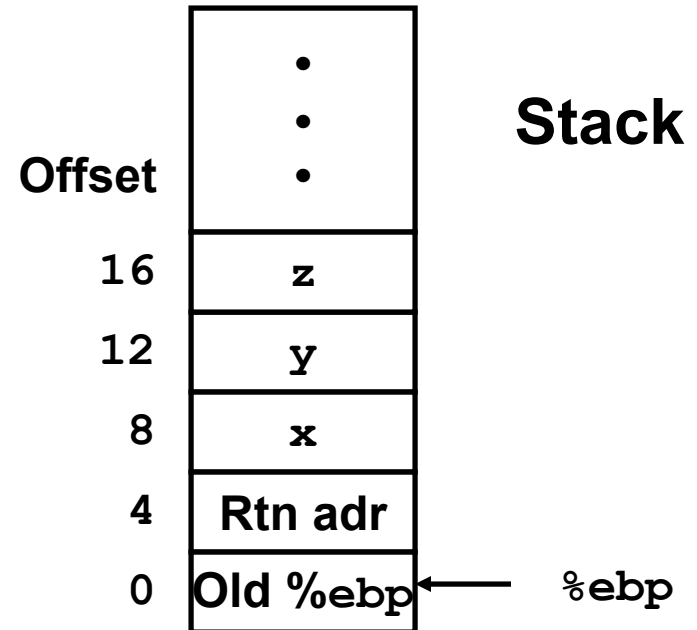
} Body

```
movl %ebp,%esp
popl %ebp
ret
```

} Finish

Understanding arith

```
int arith
(int x, int y, int z)
{
    int t1 = x+y;
    int t2 = z+t1;
    int t3 = x+4;
    int t4 = y * 48;
    int t5 = t3 + t4;
    int rval = t2 * t5;
    return rval;
}
```

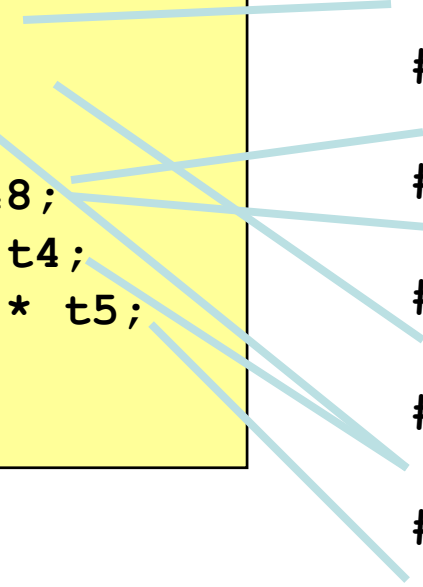


movl 8(%ebp), %eax	# eax = x
movl 12(%ebp), %edx	# edx = y
leal (%edx, %eax), %ecx	# ecx = x+y (t1)
leal (%edx, %edx, 2), %edx	# edx = 3*y
sall \$4, %edx	# edx = 48*y (t4)
addl 16(%ebp), %ecx	# ecx = z+t1 (t2)
leal 4(%edx, %eax), %eax	# eax = 4+t4+x (t5)
imull %ecx, %eax	# eax = t5*t2 (rval)

Understanding arith

```
int arith
(int x, int y, int z)
{
    int t1 = x+y;
    int t2 = z+t1;
    int t3 = x+4;
    int t4 = y * 48;
    int t5 = t3 + t4;
    int rval = t2 * t5;
    return rval;
}
```

```
# eax = x
movl 8(%ebp),%eax
# edx = y
movl 12(%ebp),%edx
# ecx = x+y (t1)
leal (%edx,%eax),%ecx
# edx = 3*y
leal (%edx,%edx,2),%edx
# edx = 48*y (t4)
sall $4,%edx
# ecx = z+t1 (t2)
addl 16(%ebp),%ecx
# eax = 4+t4+x (t5)
leal 4(%edx,%eax),%eax
# eax = t5*t2 (rval)
imull %ecx,%eax
```



Another Example

```
int logical(int x, int y)
{
    int t1 = x^y;
    int t2 = t1 >> 17;
    int mask = (1<<13) - 7;
    int rval = t2 & mask;
    return rval;
}
```

$2^{13} = 8192, 2^{13} - 7 = 8185$

logical:

```
pushl %ebp
movl %esp,%ebp
```

} Set Up

```
movl 8(%ebp),%eax
xorl 12(%ebp),%eax
sarl $17,%eax
andl $8185,%eax
```

} Body

```
movl %ebp,%esp
popl %ebp
ret
```

} Finish

```
movl 8(%ebp),%eax
xorl 12(%ebp),%eax
sarl $17,%eax
andl $8185,%eax
```

```
eax = x
eax = x^y      (t1)
eax = t1>>17  (t2)
eax = t2 & 8185
```

CISC Properties

- Instruction can reference different operand types
 - Immediate, register, memory
- Arithmetic operations can read/write memory
- Memory reference can involve complex computation
 - $Rb + S * Ri + D$
 - Useful for arithmetic expressions, too
- Instructions can have varying lengths
 - IA32 instructions can range from 1 to 15 bytes

Whose Assembler?

Intel/Microsoft Format

```
lea    eax, [ecx+ecx*2]
sub     esp, 8
cmp     dword ptr [ebp-8], 0
mov     eax, dword ptr [eax*4+100h]
```

GAS/Gnu Format

```
leal    (%ecx,%ecx,2), %eax
subl    $8, %esp
cmpl    $0, -8(%ebp)
movl    $0x100(, %eax, 4), %eax
```

- Intel/Microsoft Differs from GAS

- Operands listed in opposite order

`mov Dest, Src`

`movl Src, Dest`

- Constants not preceded by '\$', Denote hex with 'h' at end

`100h`

`$0x100`

- Operand size indicated by operands rather than operator suffix

`sub`

`subl`

- Addressing format shows effective address computation

`[eax*4+100h]`

`$0x100(, %eax, 4)`

Data Types in x86-64

C declaration	Intel data type	GAS suffix	x86-64 Size (Bytes)
char	Byte	b	1
short	Word	w	2
int	Double word	l	4
unsigned	Double word	l	4
long int	Quad word	q	8
unsigned long	Quad word	q	8
char *	Quad word	q	8
float	Single precision	s	4
double	Double precision	d	8
long double	Extended precision	t	16



x86-64 Registers

63	31	15	8	7	0	
%rax	%eax	%ax	%ah	%al		Return value
%rbx	%ebx	%bx	%bh	%bl		Callee saved
%rcx	%ecx	%cx	%ch	%cl		4th argument
%rdx	%edx	%dx	%dh	%dl		3rd argument
%rsi	%esi	%si		%sil		2nd argument
%rdi	%edi	%di		%dil		1st argument
%rbp	%ebp	%bp		%bpl		Callee saved
%rsp	%esp	%sp		%spl		Stack pointer
%r8	%r8d	%r8w		%r8b		5th argument
%r9	%r9d	%r9w		%r9b		6th argument
%r10	%r10d	%r10w		%r10b		Caller Saved
%r11	%r11d	%r11w		%r11b		Caller Saved
%r12	%r12d	%r12w		%r12b		Callee Saved
%r13	%r13d	%r13w		%r13b		Callee saved
%r14	%r14d	%r14w		%r14b		Callee saved
%r15	%r15d	%r15w		%r15b		Callee saved

rax rbp,
the r in
front
says it
represents
needed as
a 64 bit.



Swap in 32-bit Mode

```
void swap(int *xp, int *yp)
{
    int t0 = *xp;
    int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

swap:

pushl %ebp	}	Setup
movl %esp,%ebp		
pushl %ebx		
movl 12(%ebp),%ecx	}	Body
movl 8(%ebp),%edx		
movl (%ecx),%eax		
movl (%edx),%ebx		
movl %eax, (%edx)		
movl %ebx, (%ecx)		
movl -4(%ebp),%ebx	}	Finish
movl %ebp,%esp		
popl %ebp		
ret		



Swap in 64-bit Mode

```
void swap(int *xp, int *yp)
{
    int t0 = *xp;
    int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

```
swap:
    movl    (%rdi), %edx
    movl    (%rsi), %eax
    movl    %eax, (%rdi)
    movl    %edx, (%rsi)
    retq
```

- Operands passed in registers (why useful?)
 - First (xp) in %rdi, second (yp) in %rsi
 - 64-bit pointers
- No stack operations required
- 32-bit data
 - Data held in registers %eax and %edx
 - movl operation



Swap Long Ints in 64-bit Mode

```
void swap_l
(long int *xp, long int *yp)
{
    long int t0 = *xp;
    long int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

```
swap_l:
    movq    (%rdi), %rdx
    movq    (%rsi), %rax
    movq    %rax, (%rdi)
    movq    %rdx, (%rsi)
    retq
```

- 64-bit data
 - Data held in registers **%rax** and **%rdx**
 - **movq** operation
 - “q” stands for quad-word



gdb

```
40046b: c9          leaveq
40046c: c3          retq
40046d: 0f 1f 00    nopl    (%rax)

00000000400470 <frame_dummy>:
400470: 55          push    %rbp
400471: 48 83 3d 37 02 20 00    cmpl    $0x0,20977719(%rip)    # 6006b
<_JCR_END__>
400478: 00          mov     %rsp,%rbp
400479: 48 89 e5    je      400494 <frame_dummy+0x24>
40047c: 74 16      mov     $0x0,%eax
40047e: b8 00 00 00 00    test    %rax,%rax
400483: 48 85 c0    je      400494 <frame_dummy+0x24>
400486: 74 0c      mov     $0x6006b0,%edi
400488: bf b0 06 60 00    mov     %rax,%r11
40048d: 49 89 c3    leaveq  %r11
400490: c9          jmpq    *%r11
400491: 41 ff e3    leaveq
400494: c9          retq
400495: c3          nopl
400496: 90          nopl
400497: 90          nopl
400498: 90          nopl
400499: 90          nopl
40049a: 90          nopl
40049b: 90          nopl
40049c: 90          nopl
40049d: 90          nopl
40049e: 90          nopl
40049f: 90          nopl

000000004004a0 <main>:
4004a0: 48 89 5c 24 f0    mov     %rbx,0xffffffffffffff0(%rsp)
4004a5: 4c 89 64 24 f8    mov     %r12,0xffffffffffffff8(%rsp)
4004aa: 48 83 ec 18      sub     $0x18,%rsp
4004ae: e8 f5 fe ff ff    callq   4003a8 <rand@plt>
4004b3: 89 c3          mov     %eax,%ebx
4004b5: e8 ee fe ff ff    callq   4003a8 <rand@plt>
4004ba: 41 89 c4          mov     %eax,%r12d
4004bd: 41 21 dc          and     %ebx,%r12d
4004c0: e8 e3 fe ff ff    callq   4003a8 <rand@plt>
4004c5: 41 31 c4          xor     %eax,%r12d
4004c8: 48 8b 5c 24 08    mov     0x8(%rsp),%rbx
4004cd: 44 89 e0          mov     %r12d,%eax
4004d0: 4c 8b 64 24 10    mov     0x10(%rsp),%r12
4004d5: 48 83 c4 18      add     $0x18,%rsp
4004d9: c3          retq
4004da: 90          nopl
4004db: 90          nopl
4004dc: 90          nopl
4004dd: 90          nopl
4004de: 90          nopl
4004df: 90          nopl

000000004004e0 <__libc_csu_fini>:
4004e0: f3 c3      repz retq
4004e2: 0f 1f 80 00 00 00    nopl    0x0(%rax)
4004e7: 0f 1f 80 00 00 00    nopl    0x0(%rax)

000000004004f0 <__libc_csu_init>:
4004f0: 4c 89 64 24 e0    mov     %r12,0xffffffffffffe0(%rsp)
4004f5: 4c 89 6c 24 e8    mov     %r13,0xffffffffffffe8(%rsp)
4004fa: 4c 8d 25 8b 01 20 00    lea     2097547(%rip),%r12    # 60068
<_fini_array_end>
400501: 4c 89 74 24 f0    mov     %r14,0xfffffffffffff0(%rsp)
400506: 4c 89 7c 24 f8    mov     %r15,0xfffffffffffff8(%rsp)
40050b: 49 89 f6          mov     %rsi,%r14
40050e: 48 89 5c 24 d0    mov     %rbx,0xffffffffffffd0(%rsp)
400513: 48 89 6c 24 d8    mov     %rbp,0xffffffffffffd8(%rsp)
400518: 48 83 ec 38      sub     $0x38,%rsp
40051c: 41 89 ff          mov     %edi,%r15d
40051f: 49 89 d5          mov     %rdx,%r13
400522: e8 49 fe ff ff    callq   400370 <_init>
400527: 48 8d 05 5e 01 20 00    lea     2097502(%rip),%rax    # 60068
<_fini_array_end>
40052e: 49 29 c4          sub     %rax,%r12
400531: 49 c1 fc 03      sar     $0x3,%r12
400535: 4d 85 e4          test    %r12,%r12
400538: 74 1e          je      400558 <__libc_csu_init+0x68>
40053a: 31 ed          xor     %ebp,%ebp
40053c: 48 89 c3          mov     %rax,%rbx
40053f: 90          nopl
```

```
(gdb) break *0x4004c5
Breakpoint 3 at 0x4004c5
(gdb) run
Starting program: /w/fac.3/cs/reinman/code/a.out
```

Breakpoint 3, 0x000000004004c5 in main (< >)

```
(gdb) i r
rax      0x643c9869      1681692777
rbx      0x6b8b4567      1804289383
rcx      0x31ead5106c    214393229420
rdx      0x31ead5107c    214393229436
rsi      0x7fffffff9774  140737488349556
rdi      0x31ead514c0    214393230528
rbp      0x0            0
rsp      0x7fffffff990    0x7fffffff990
r8       0x31ead51064    214393229412
r9       0x31ead510e0    214393229536
r10      0x0            0
r11      0x31eaa340e0     214389965024
r12      0x220b0145     571146566
r13      0x7fffffff9a80  140737488349824
r14      0x0            0
r15      0x0            0
rip      0x4004c5 0x4004c5 <main+37>
eflags   0x206 [ PF IF ]
cs       0x33 51
ss       0x2b 43
ds       0x0 0
es       0x0 0
fs       0x0 0
gs       0x0 0
fctrl    0x37f 895
fstat    0x0 0
ftag     0xffff 65535
fiseg    0x0 0
fioff    0x0 0
foseg    0x0 0
fofff    0x0 0
fop      0x0 0
mxcsr    0x1f80 [ IM DM ZM OM UM PM ]
```

(gdb) si
0x000000004004c8 in main (< >)

```
(gdb) i r
rax      0x643c9869      1681692777
rbx      0x6b8b4567      1804289383
rcx      0x31ead5106c    214393229420
rdx      0x31ead5107c    214393229436
rsi      0x7fffffff9774  140737488349556
rdi      0x31ead514c0    214393230528
rbp      0x0            0
rsp      0x7fffffff990    0x7fffffff990
r8       0x31ead51064    214393229412
r9       0x31ead510e0    214393229536
r10      0x0            0
r11      0x31eaa340e0     214389965024
r12      0x4637992f     1178048815
r13      0x7fffffff9a80  140737488349824
r14      0x0            0
r15      0x0            0
rip      0x4004c8 0x4004c8 <main+40>
eflags   0x202 [ IF ]
cs       0x33 51
ss       0x2b 43
ds       0x0 0
es       0x0 0
fs       0x0 0
gs       0x0 0
fctrl    0x37f 895
fstat    0x0 0
ftag     0xffff 65535
fiseg    0x0 0
fioff    0x0 0
foseg    0x0 0
fofff    0x0 0
fop      0x0 0
mxcsr    0x1f80 [ IM DM ZM OM UM PM ]
```

```
(gdb) x/4x 0x4004c8
0x4004c8 <main+40>: 0x245c8b48 0xe0894408 0x24648b4c 0xc48348
(gdb) x/8x 0x4004c8
0x4004c8 <main+40>: 0x245c8b48 0xe0894408 0x24648b4c 0xc48348
0x4004d0 0x9090c318 0x90909090 0x1f0fc3f3 0x00000000
(gdb)
```