

# What can abstract mathematics tell us about programming climate models?

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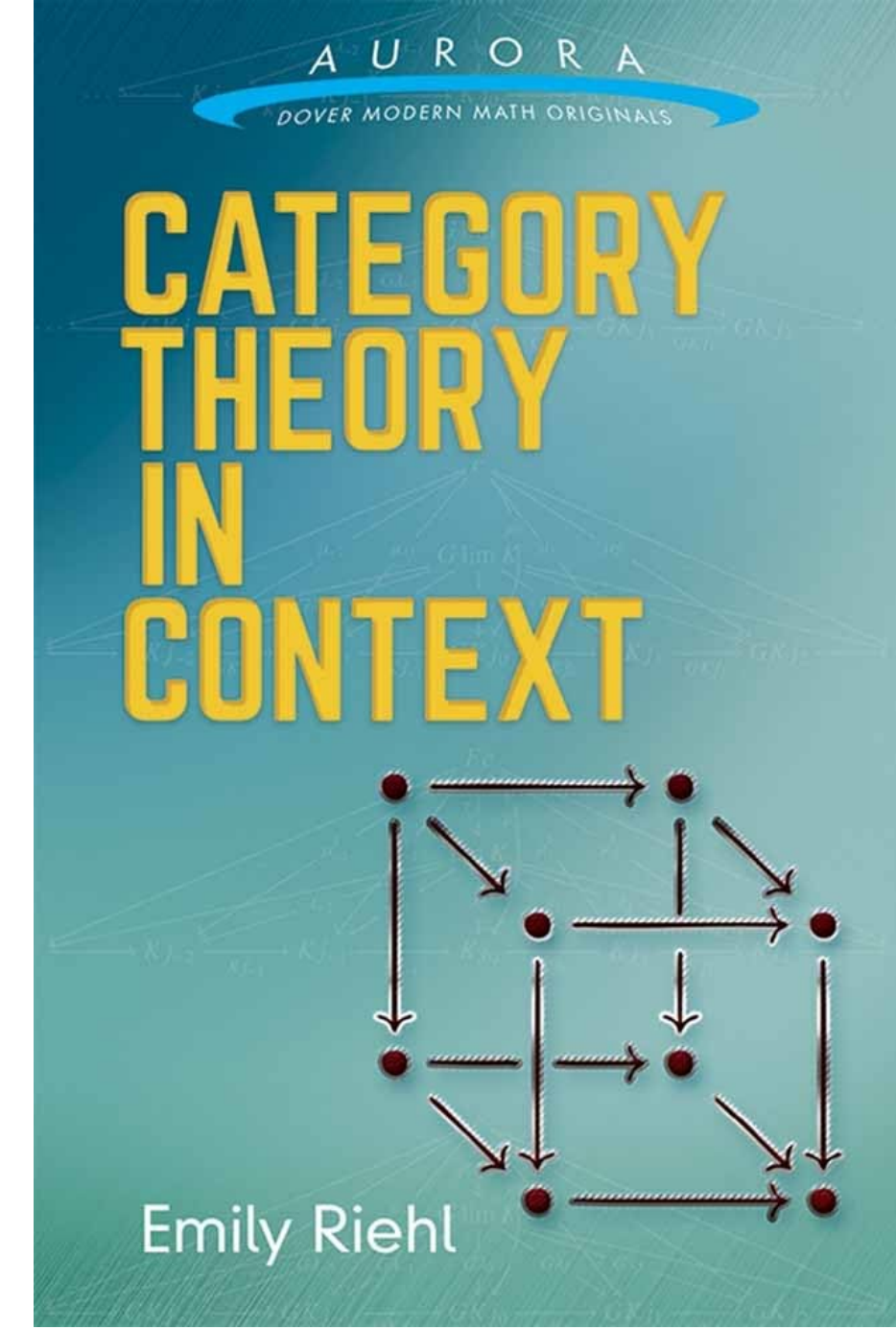
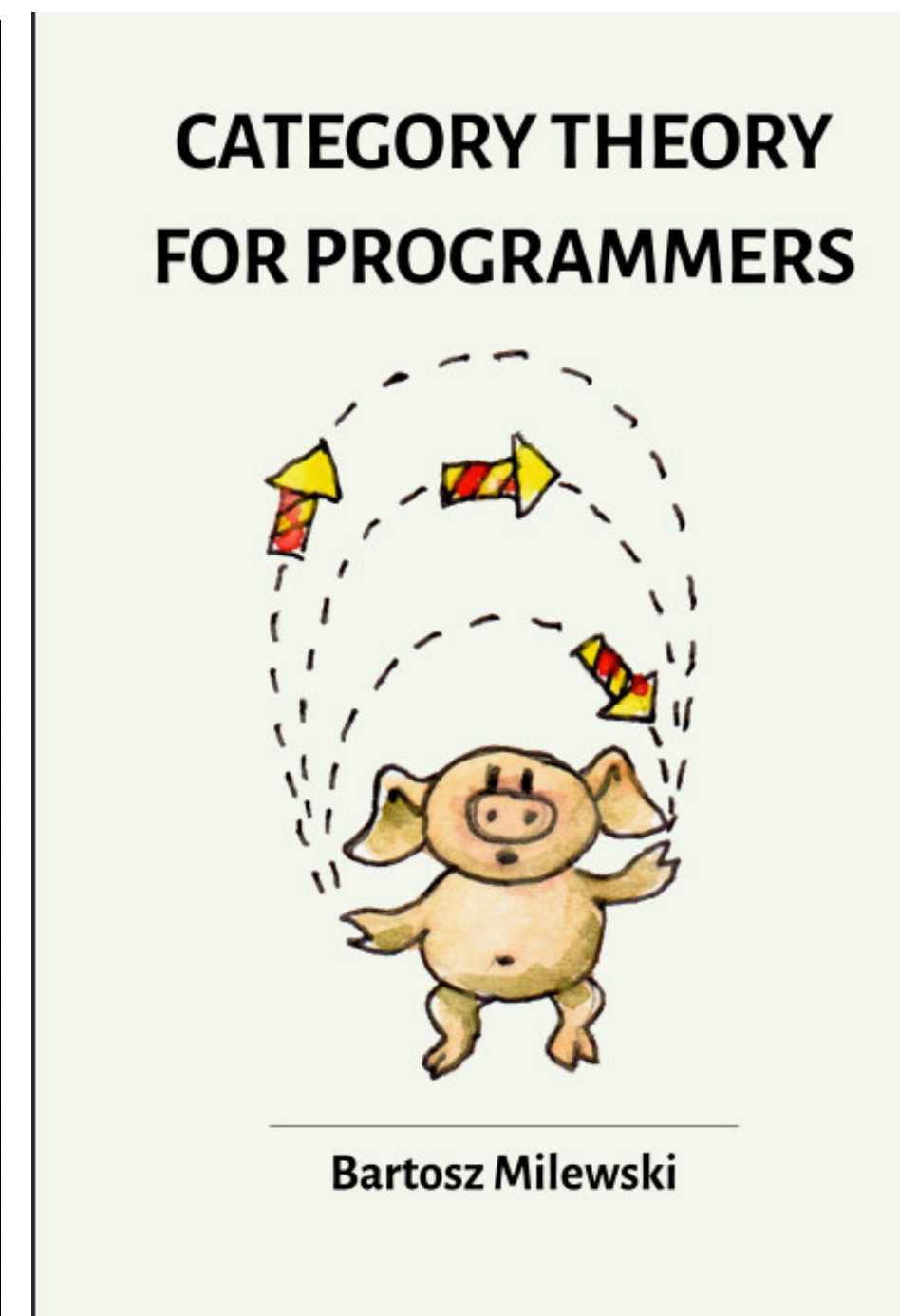
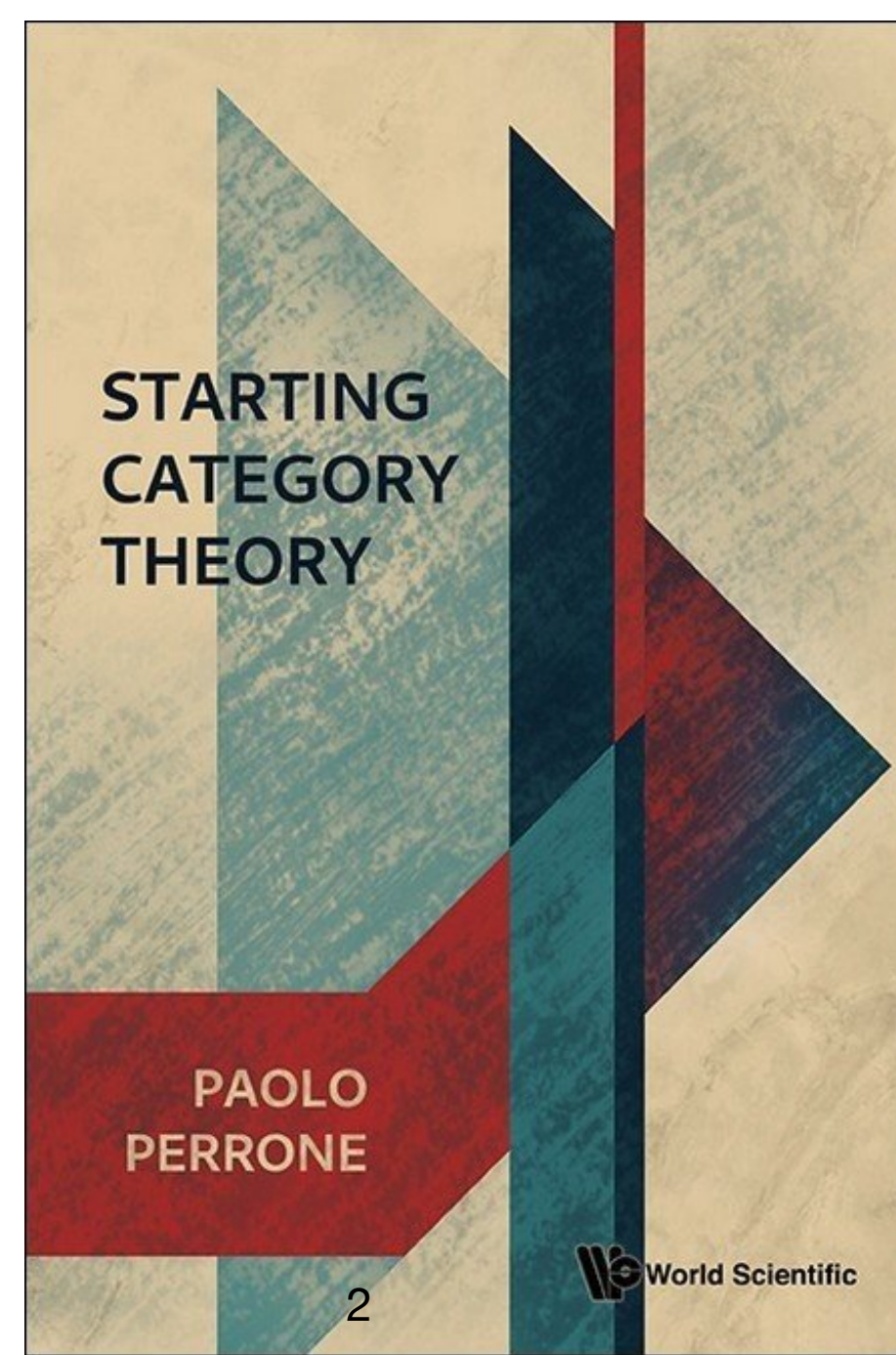
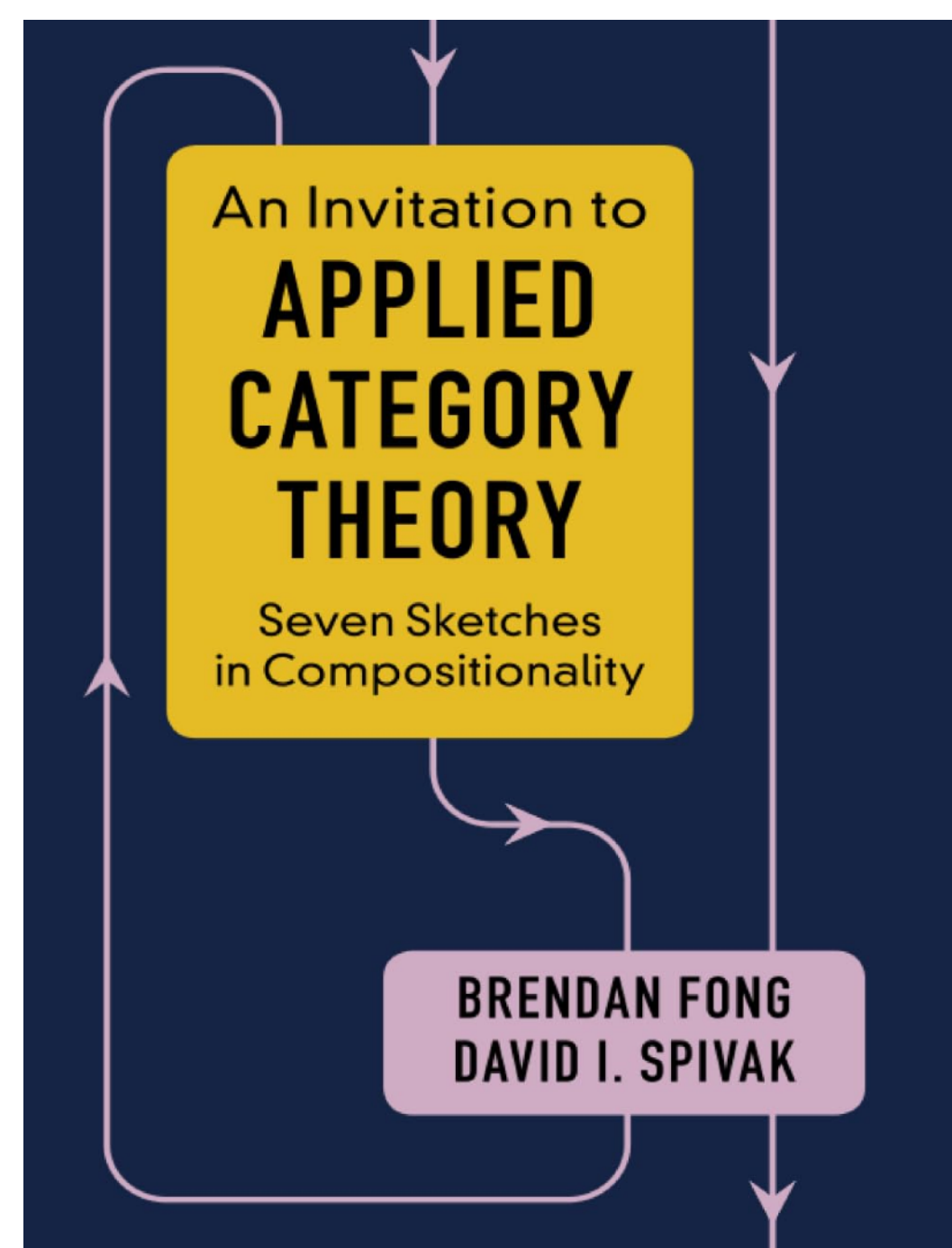
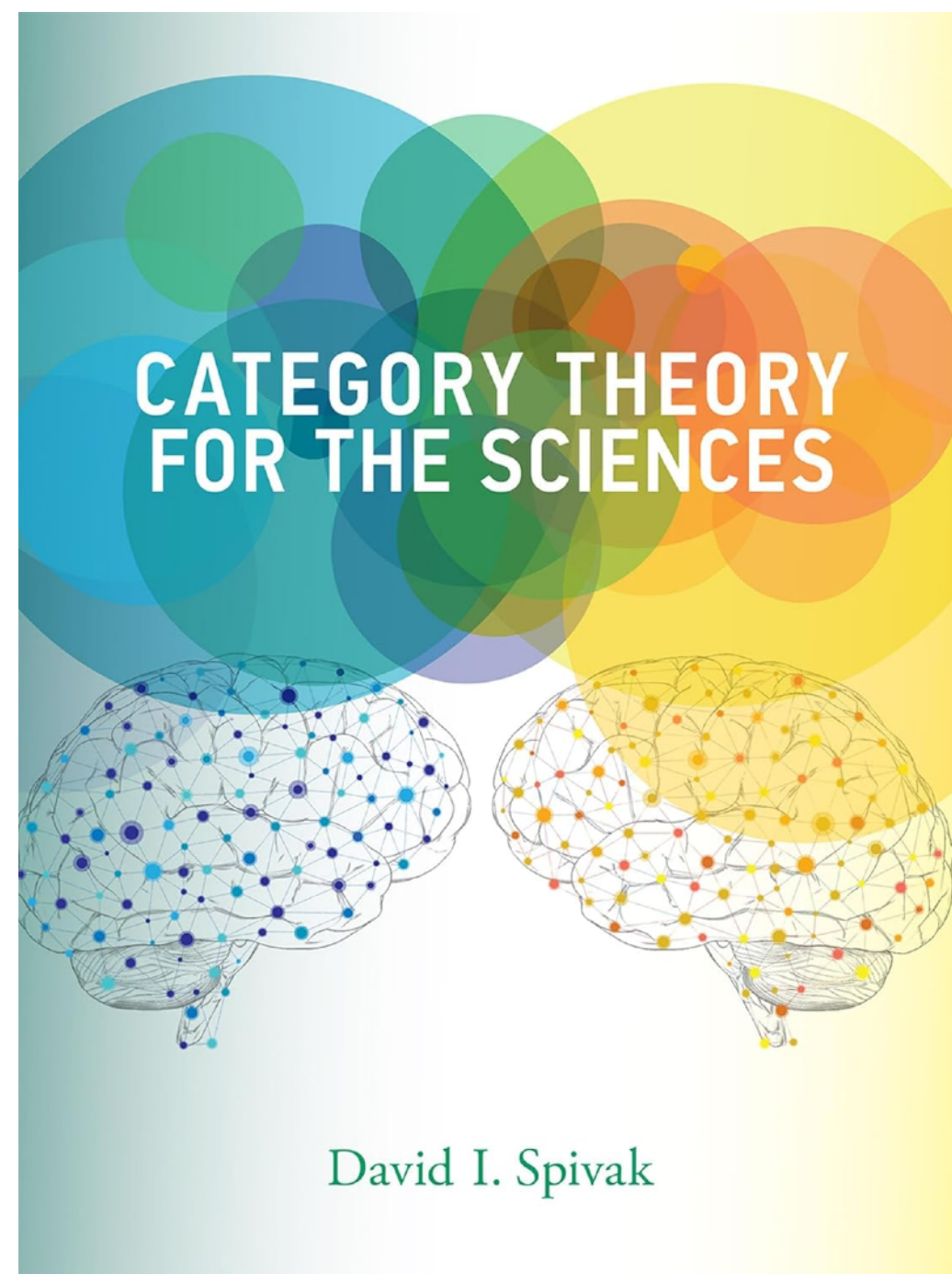




# Category theory

A branch of abstract mathematics

Provides organising principles exposing mathematical structure and structure within mathematics





# Today's objectives

*View some common ideas in numerical programming through the category theory lens*

- Understand what is a category, functor, and natural transformation
- Understand their relationship to particular programming patterns
- Use these basic categorical concepts to inform **testing** and **optimisation** strategies
- Understand these basics as a launch pad

# Categories

- Collection of objects  $\text{obj}(\mathcal{C})$ 
  - often denoted  $A, B, C, \dots \in \text{obj}(\mathcal{C})$
- Collection of “morphisms” (relationships) between pairs of objects  $\text{morph}(\mathcal{C})$ 
  - often denoted  $f : A \rightarrow B, g : B \rightarrow C, \dots \in \text{morph}(\mathcal{C})$
- Composition operation

$$\frac{f : A \rightarrow B \quad g : B \rightarrow C}{f; g : A \rightarrow C}$$

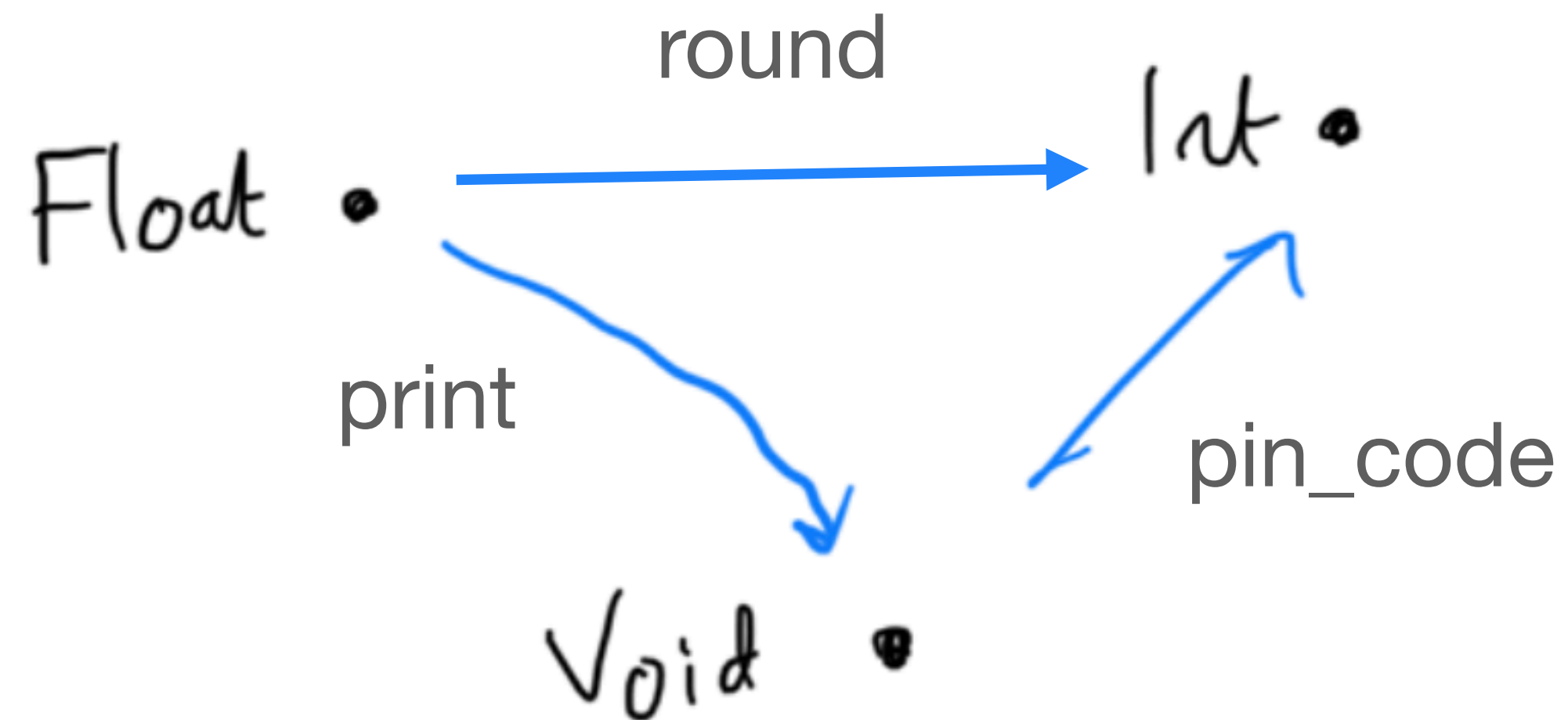
- Identity morphisms

$$\forall A \in \text{obj}(\mathcal{C}) . id_A : A \rightarrow A$$

- **Axioms:**  $f; id = id; f = f$   $f; (g; h) = (f; g); h$

# Categories in programming

## Types as objects, functions as morphisms



# Functors

A morphism between categories

$$F : \mathcal{C} \rightarrow \mathcal{D}$$

*Has two components:*

(1) mapping on objects:  $\forall A \in \text{obj}(\mathcal{C}) \implies F(A) \in \text{obj}(\mathcal{D})$

(2) mapping on morphisms:

$\forall (f : A \rightarrow B) \in \text{morph}(\mathcal{C}) \implies F(f) : F(A) \rightarrow F(B) \in \text{morph}(\mathcal{D})$

*Has two axioms:*

$$F(\text{id}_A) = \text{id}_{FA}$$

$$F(f; g) = F(f); F(g)$$

# Functors in programming

**Object mapping is a type constructor**

`list : Type -> Type`

**Morphism mapping lifts a function to apply over the type constructor**

$$\frac{A \rightarrow B}{\text{list}(A) \rightarrow \text{list}(B)}$$

# Natural transformations

A morphism between functors

$$F : \mathcal{C} \rightarrow \mathcal{D} \qquad G : \mathcal{C} \rightarrow \mathcal{D}$$

*A natural transformation  $\alpha : F \rightarrow G$  is a family of morphisms*

meaning  $\forall A \in \text{obj}(\mathcal{C}). \alpha_A : FA \rightarrow GA \in \text{morph}(\mathcal{D})$

along with the **naturality property**

$$\begin{array}{ccc} FA & \xrightarrow{\alpha_A} & GA \\ \downarrow Ff & & \downarrow Gf \\ FB & \xrightarrow{\alpha_B} & GB \end{array}$$



# Natural transformations in programming

# Connecting back to climate modelling...

- Functors represent pointwise traversals (common transformation)
- Structural transformations are natural transformations
- Stencil computations are a generalisation to functors....

# Comonads (*briefly and incompletely*)

An endofunctor  $D : \mathcal{C} \rightarrow \mathcal{C}$  with:

- Natural transformation  $\varepsilon_A : DA \rightarrow A$
- For every morphism, a lifting  $(-)^{\dagger}$

*Functorial mapping*

$$\frac{f : DA \rightarrow B}{f^{\dagger} : DA \rightarrow DB}$$

*cf.*

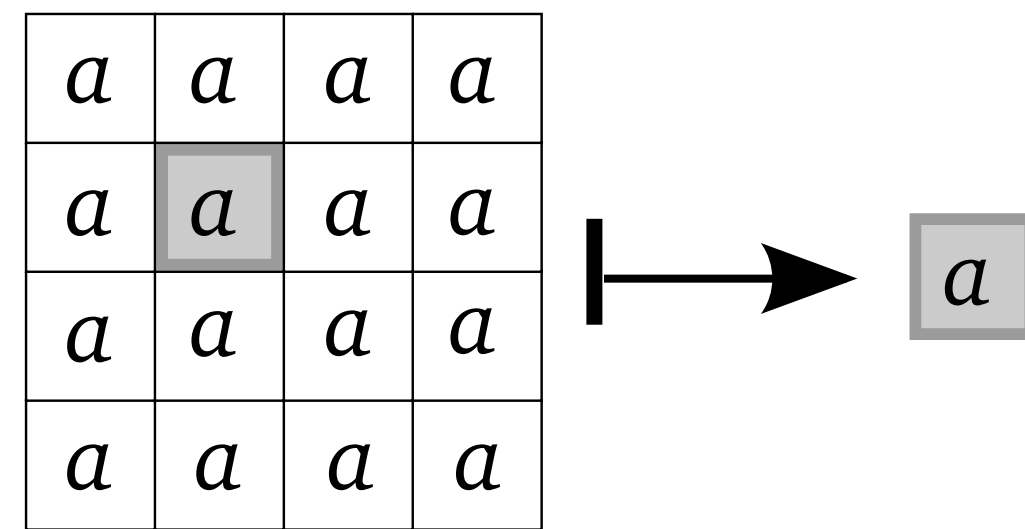
$$\frac{f : A \rightarrow B}{Ff : DA \rightarrow DB}$$

# Comonads in programming

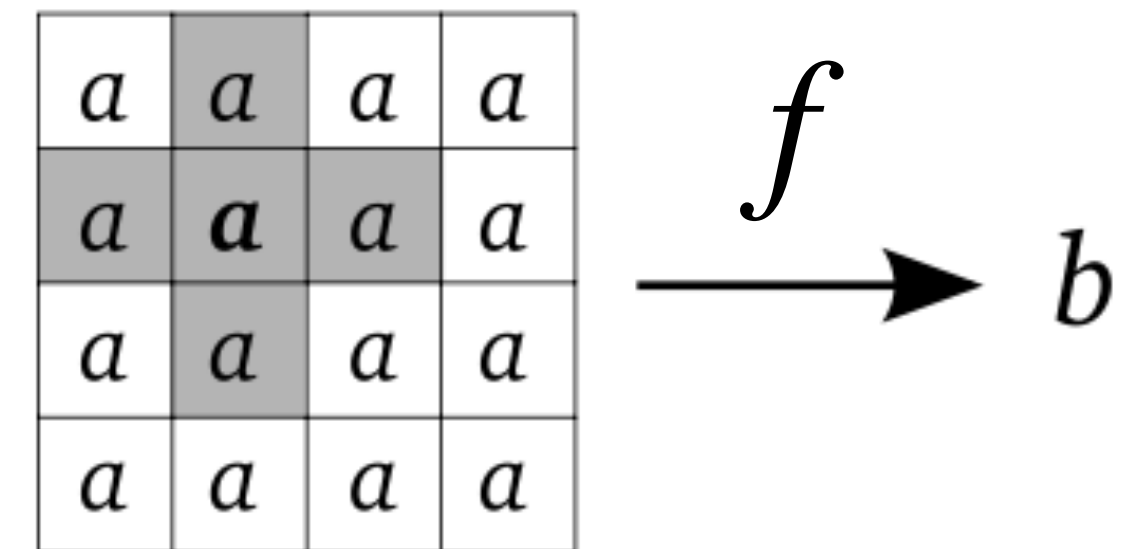
## Array comonad

$$DA = \text{Array } I \ A \times I$$

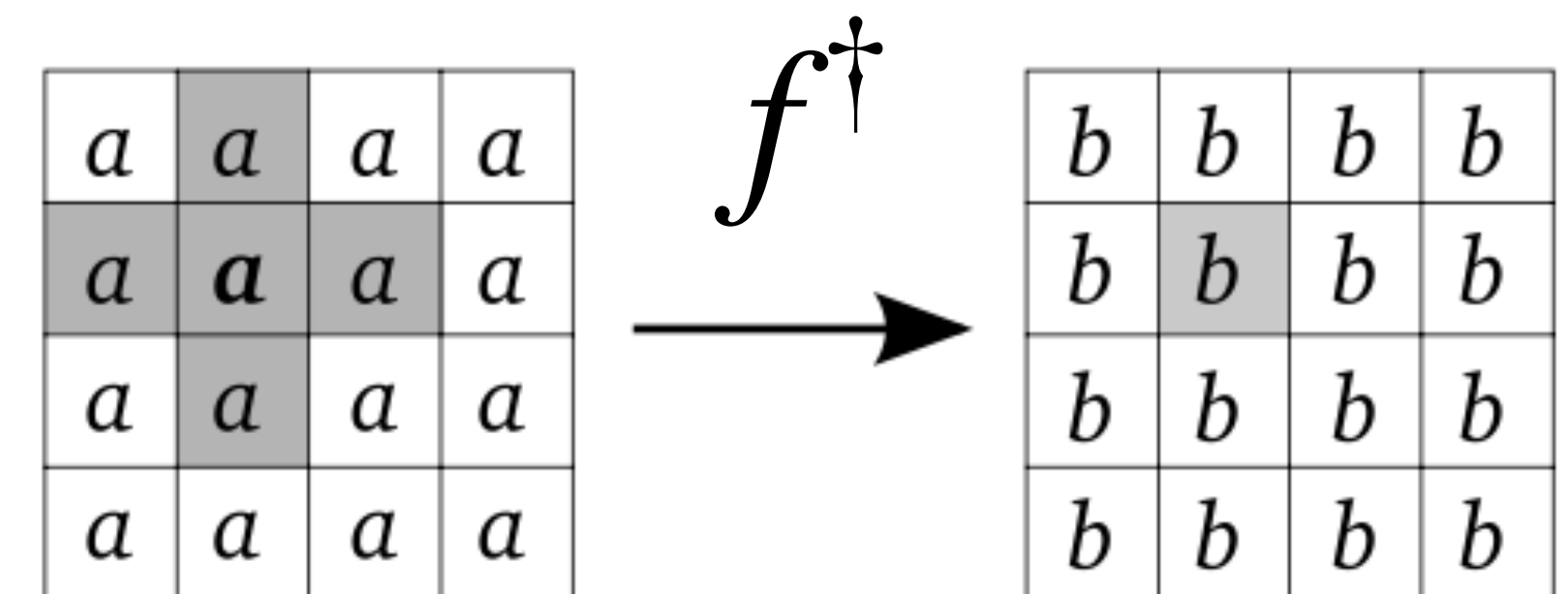
array-data  $\times$  cursor



$$\varepsilon_A : DA \rightarrow A$$



**Local computation  
(neighbourhood)**



**Global computation**



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# Thanks!

