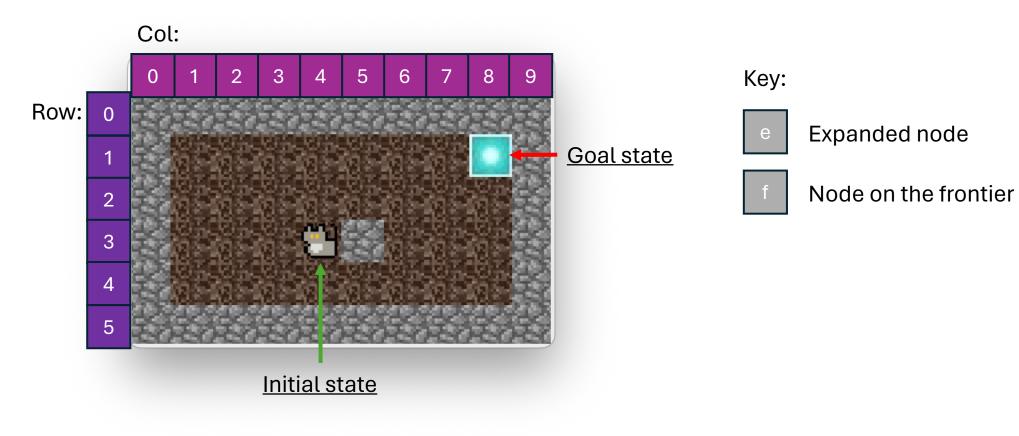
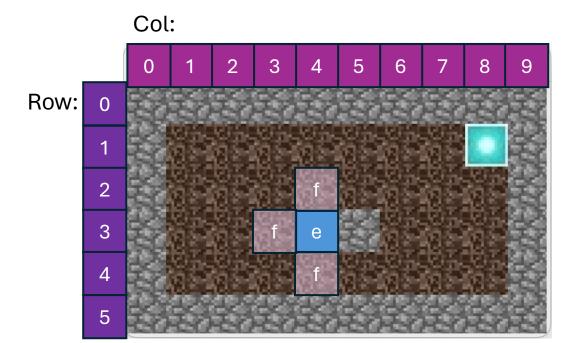
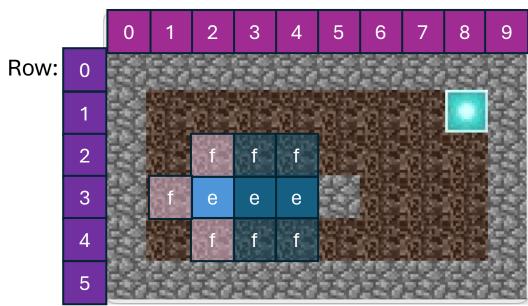
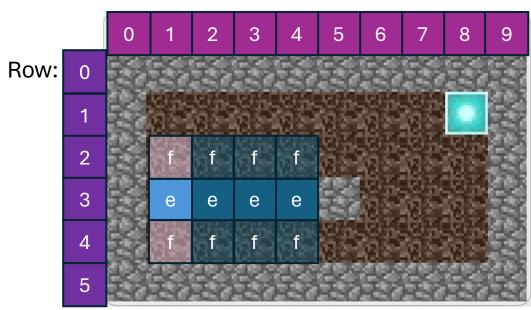
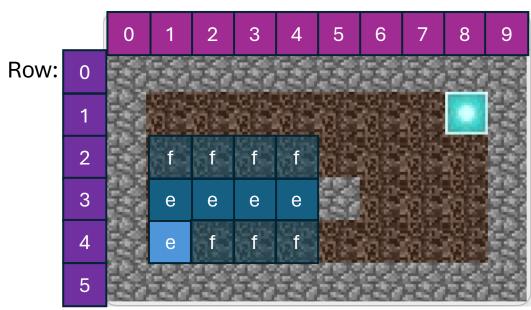
2-D Maze Example



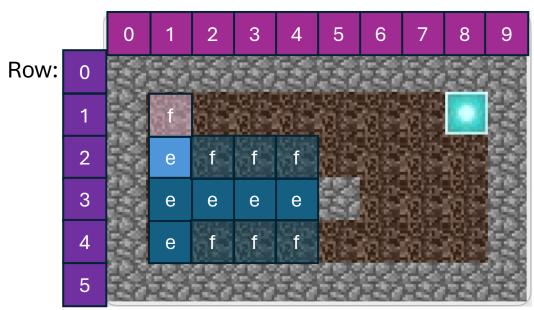




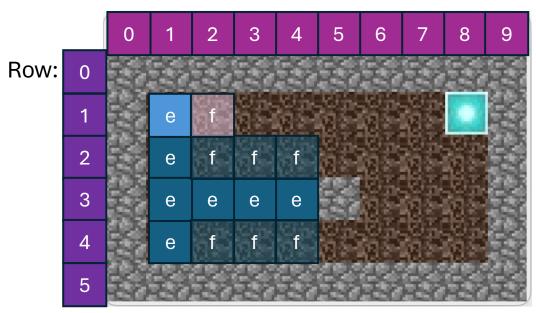




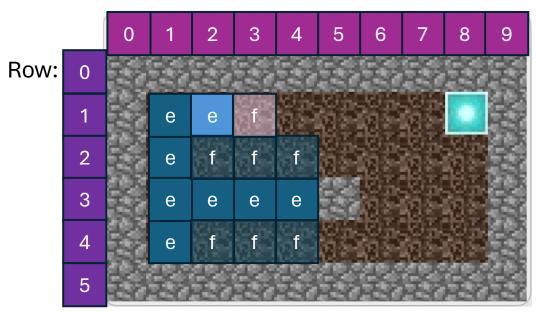
DFS



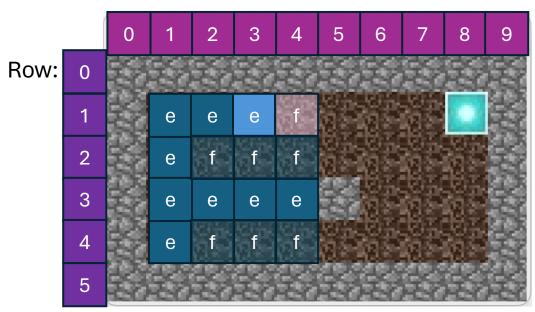
DFS



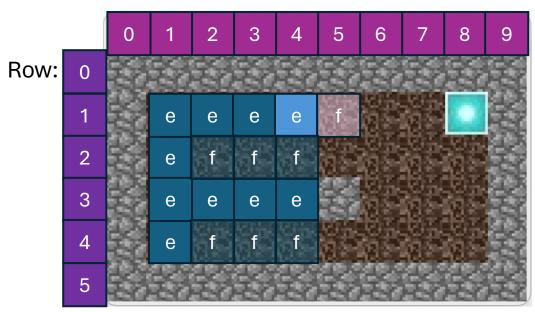
DFS



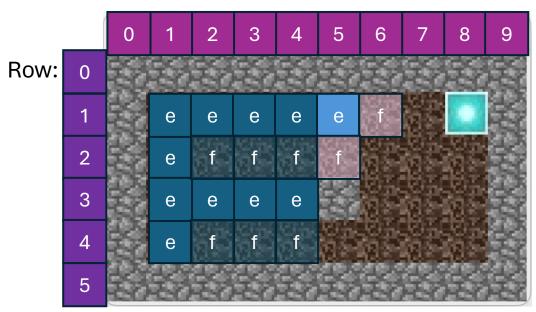
DFS



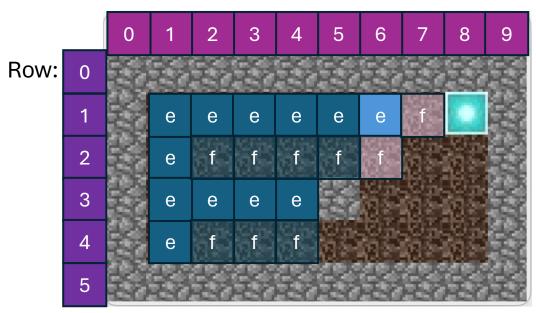
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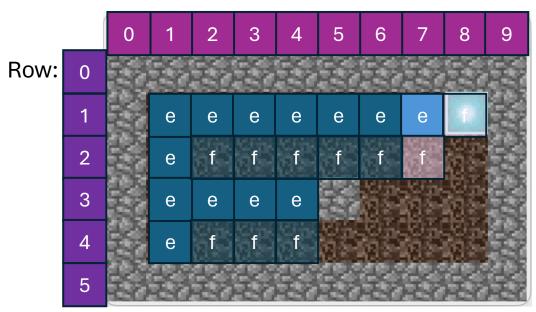
DFS



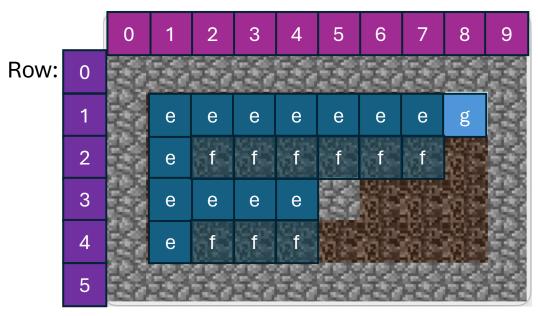
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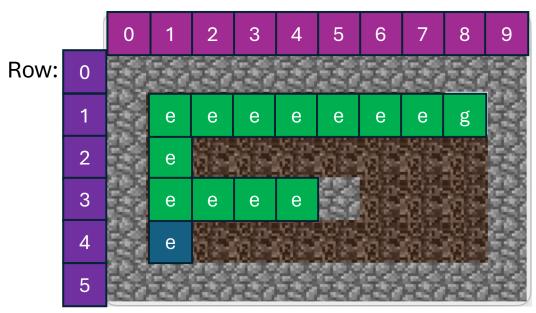


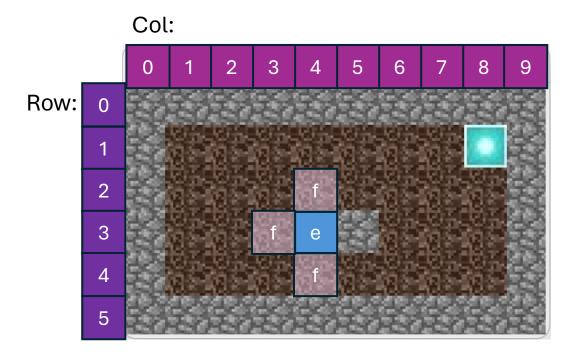
DFS

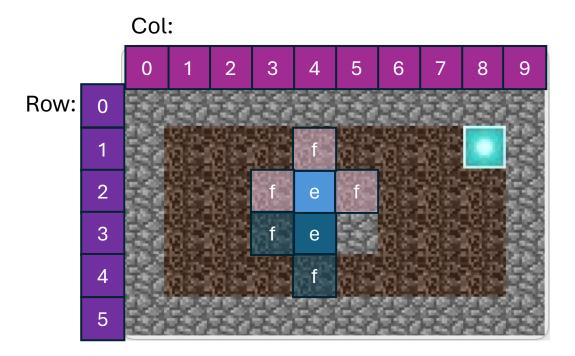


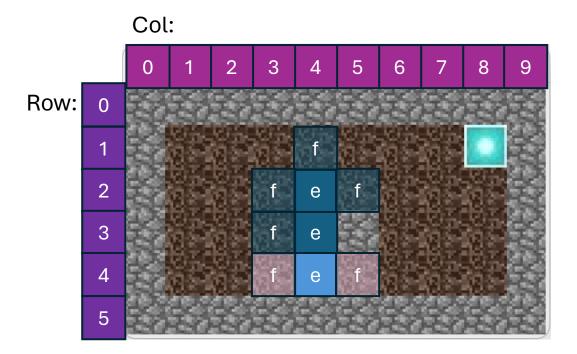
DFS

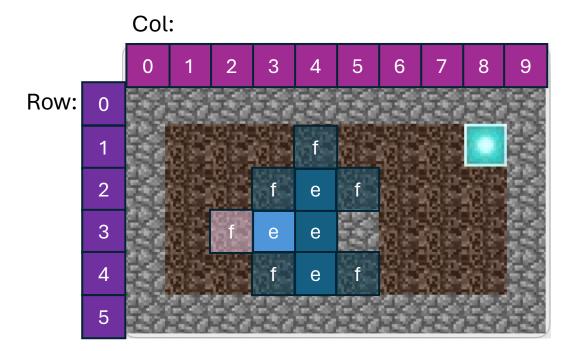


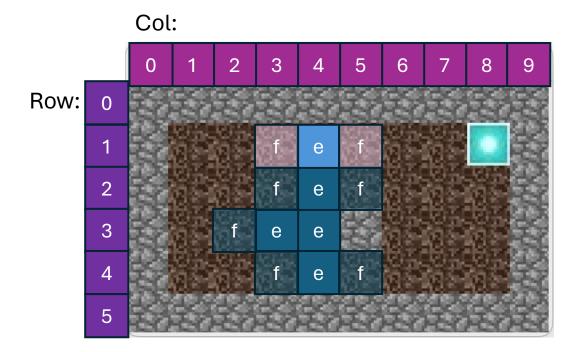








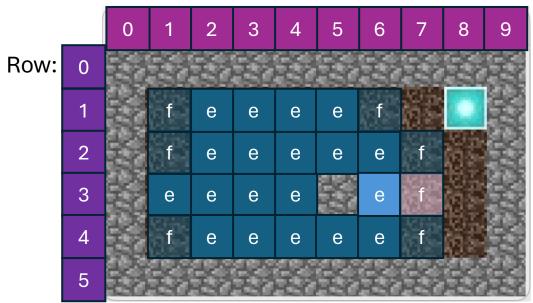


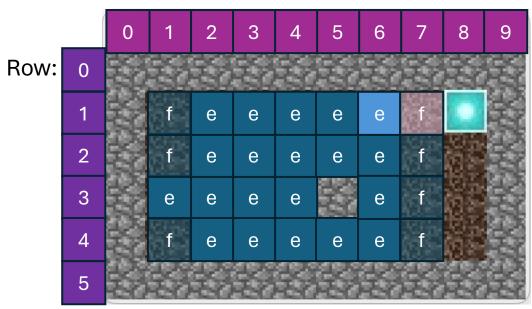


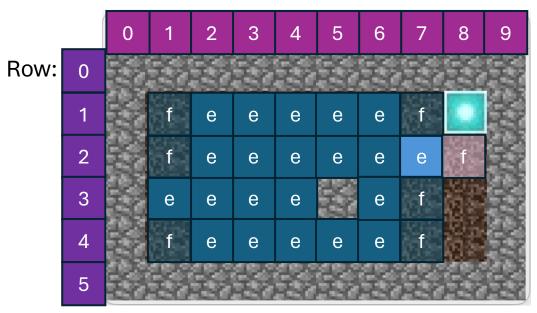
Col: O 1 2 3 4 5 6 7 8 9 Row: O f e e e 4 f e e f 5

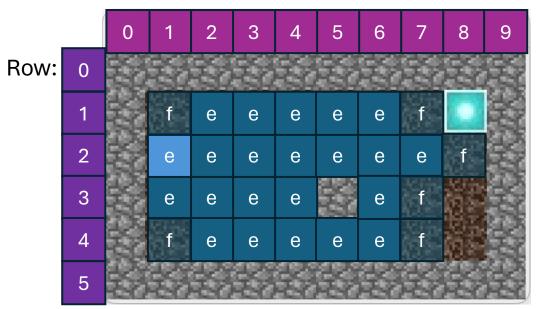
Col: O 1 2 3 4 5 6 7 8 9 Row: O f e e e e f 2 f e e e e f 4 f e e e e e f 5

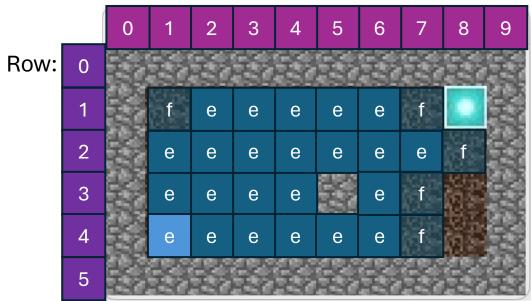
Col: O 1 2 3 4 5 6 7 8 9 Row: O f e e e e f 6 6 7 8 9 4 f e e e e e f 7 7 8 9

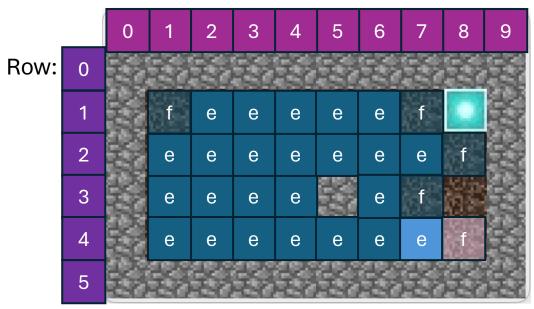


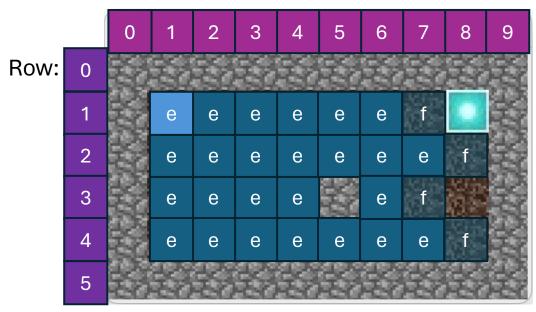


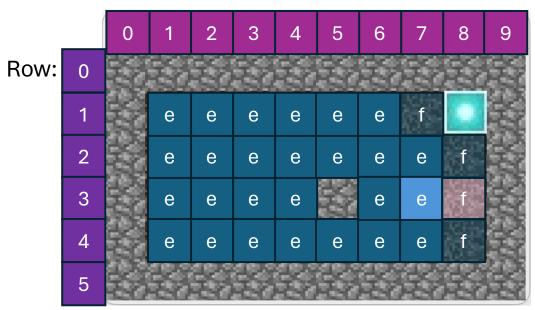


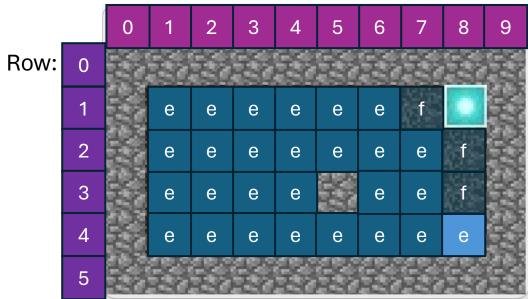




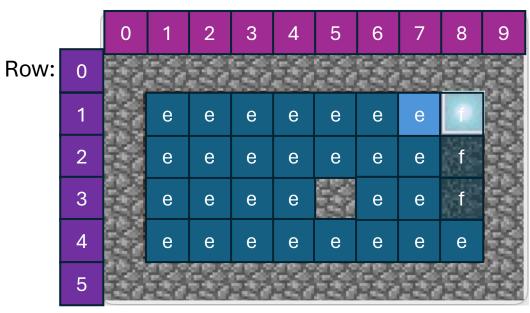








Col:



In **BFS**, the goal check is performed on node **generation**, so BFS will find the goal at this step. **UCS** performs the goal check on node **expansion** (i.e. when the node is selected from the frontier), so continues searching.

```
function Breadth-First-Search(problem) returns a solution node or failure

node ← Node(problem.initial)

if problem.Is-Goal(node.State) then return node

frontier ← a FIFO queue, with node as an element

reached ← {problem.initial}

while not Is-Empty(frontier) do

node ← Pop(frontier)

for each child in Expand(problem, node) do

s ← child State

if problem.Is-Goal(s) then return child

if s is not in reached then

add s to reached

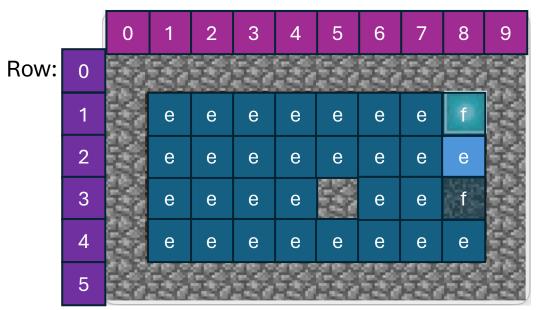
add child to frontier

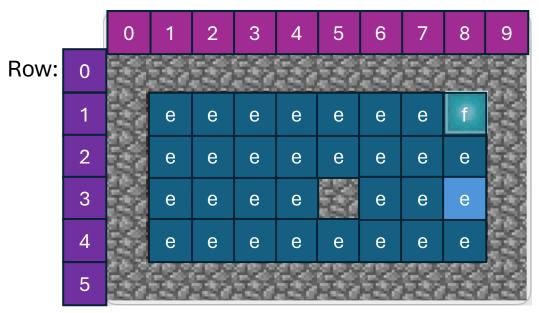
return failure
```

```
function UNIFORM-COST-SEARCH(problem) returns a solution node, or failure return BEST-FIRST-SEARCH(problem, PATH-COST)
```

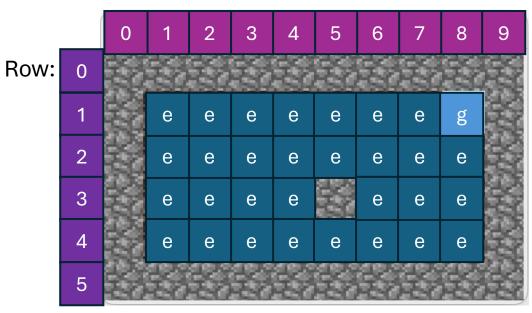
```
function BEST-FIRST-SEARCH(problem, f) returns a solution node or failure
  node \leftarrow Node(STATE=problem.INITIAL)
  frontier \leftarrow a priority queue ordered by f, with node as an element
  reached \leftarrow a lookup table, with one entry with key problem. INITIAL and value node
  while not IS-EMPTY(frontier) do
    node \leftarrow POP(frontier)
    if problem.IS-GOAL(node.STATE) then return node
    for each child in EXPAND(problem, node) do
       s \leftarrow child.STATE
       if s is not in reached or child.PATH-COST < reached[s].PATH-COST then
         reached[s] \leftarrow child
         add child to frontier
  return failure
function EXPAND(problem, node) vields nodes
  s \leftarrow node.STATE
  for each action in problem. ACTIONS(s) do
    s' \leftarrow problem.RESULT(s, action)
    cost \leftarrow node.PATH-COST + problem.ACTION-COST(s, action, s')
    yield NODE(STATE=s', PARENT=node, ACTION=action, PATH-COST=cost)
```

Figure 3.7 The best-first search algorithm, and the function for expanding a node. The data structures used here are described in Section 3.3.2. See Appendix B for **yield**.





Col:



Since **UCS** performs the goal check on node **expansion** (i.e. when the node is selected from the frontier), it only detects the goal now.

```
function BREADTH-FIRST-SEARCH(problem) returns a solution node or failure

node ← NODE(problem.INITIAL)

if problem.Is-GOAL(node.STATE) then return node

frontier ← a FIFO queue, with node as an element

reached ← {problem.INITIAL}

while not Is-EMPTY(frontier) do

node ← POP(frontier)

for each child in EXPAND(problem, node) do

s← child STATE

if problem.Is-GOAL(s) then return child

if s is not in reached then

add s to reached

add child to frontier

return failure
```

```
add child to frontier

return failure

function EXPAND(problem, node) yields nodes

s ← node.STATE

for each action in problem.ACTIONS(s) do

s' ← problem.RESULT(s, action)

cost ← node.PATH-COST + problem.ACTION-COST(s, action, s')

yield NODE(STATE=s', PARENT=node, ACTION=action, PATH-COST=cost)
```

function BEST-FIRST-SEARCH(problem, f) **returns** a solution node or failure

 $reached \leftarrow$ a lookup table, with one entry with key problem. INITIAL and value node

if *s* is not in *reached* **or** *child*.PATH-COST < *reached*[*s*].PATH-COST **then**

frontier \leftarrow a priority queue ordered by f, with node as an element

if problem.IS-GOAL(node.STATE) then return node for each child in EXPAND(problem, node) do

function UNIFORM-COST-SEARCH(problem) returns a solution node, or failure
return BEST-FIRST-SEARCH(problem, PATH-COST)

Figure 3.7 The best-first search algorithm, and the function for expanding a node. The data structures used here are described in Section 3.3.2. See Appendix B for yield.

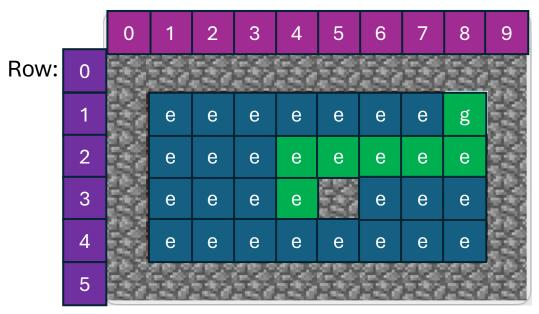
 $node \leftarrow Node(STATE=problem.INITIAL)$

while not IS-EMPTY(frontier) do

 $reached[s] \leftarrow child$

 $node \leftarrow POP(frontier)$

 $s \leftarrow child.STATE$



Heuristic \rightarrow Manhattan distance to the goal

		0	1	2	3	4	5	6	7	8	9
Row:	0										
	1		7	6	5	4	3	2	1	0	
	2		8	7	6	5	4	3	2	1	
	3		9	8	7	6		4	3	2	8
	4		10	9	8	7	6	5	4	3	
	5										

Col: O 1 2 3 4 5 6 7 8 9 Row: 1 7 6 5 4 3 2 1 0 2 8 7 6 5 4 3 2 1 3 9 8 7 6 4 3 2 4 10 9 8 7 6 5 4 3 5

Col: O 1 2 3 4 5 6 7 8 9 Row: O 7 6 5 4 3 2 1 0 2 8 7 6 5 4 3 2 1 3 9 8 7 6 5 4 3 2 4 10 9 8 7 6 5 4 3 5

Col: O 1 2 3 4 5 6 7 8 9 Row: O 7 6 5 4 3 2 1 0 2 8 7 6 e 4 3 2 1 3 9 8 7 e 4 3 2 4 10 9 8 7 6 5 4 3 5

Col: O 1 2 3 4 5 6 7 8 9 Row: O 7 6 5 e 3 2 1 O 2 8 7 6 e 4 3 2 1 3 9 8 7 e 4 3 2 4 10 9 8 7 6 5 4 3 5

Col: O 1 2 3 4 5 6 7 8 9 Row: O 7 6 5 e e 2 1 O 2 8 7 6 e 4 3 2 1 3 9 8 7 e 4 3 2 4 10 9 8 7 6 5 4 3 5

Col: O 1 2 3 4 5 6 7 8 9 Row: O 7 6 5 e e e 1 O 2 8 7 6 e 4 3 2 1 3 9 8 7 e 4 3 2 4 10 9 8 7 6 5 4 3 5

Col: 8 3 Row: е е е е 8 е 3 е 4 8 5 3

Col: 8 3 Row: е е е g е 8 е 3 е 4 8 5 3

Col: O 1 2 3 4 5 6 7 8 9 Row: O 7 6 5 e e e e e e 2 8 7 6 e 4 3 2 1 3 9 8 7 e 4 3 2 4 10 9 8 7 6 5 4 3