

# OCG - Owner vs. Collector

## 1. Model Overview

This case study recreates the Owner–Collector Game (OCG) as described in Wu et al. (2020) using a Vickery–Clarke–Groves style auction, represented as a taxi hailing game. In the OCG setting, a data owner wants the benefit of a service (a taxi), and the data collector wants to extract truthful location reports while protecting user privacy. The model we implement follows the representative approach used in ride-matching systems:

- A VCG mechanism incentivizes honest reporting of user valuations.
- A differentially private exponential mechanism randomizes which user is selected, protecting location privacy.
- The system studies the tradeoff between efficiency (accurate allocation) and privacy (strength of noise).

This game fits the Owner vs. Collector category because the collector designs the mechanism (allocation + payment rule), and the owner chooses what information to reveal based on incentives and privacy concerns.

## 2. Players

- Owner (Passenger):
  - Reports a valuation for service and a location. Wants high utility while preserving privacy.
- Collector (Platform):
  - Wants accurate service allocation but must incorporate privacy protections that reduce utility.

## 3. Strategies

- Owner (Passenger):
  - Truthfully report valuation and location
  - Misreport valuation and/or location
- Collector (Platform):
  - Choose privacy level  $\epsilon$
  - Choose allocation rule (exponential mechanism)
  - Compute VCG-style payments

## 4. Payoff Functions

BLUF:

- If  $s_i$  is the best score and  $\epsilon$  is high, that person is almost always chosen.
- If  $\epsilon$  is low, exponentials compress and all choices look similar, having stronger privacy.

Let:

- $v_i$  = valuation of user i
- $c_i$  = distance/cost to serve user i
- $s_i = v_i - c_i$  (score)
- $\epsilon$  = DP privacy parameter

### Allocation Probability (Exponential Mechanism)

The probability of any one person being chosen is:

- $\exp(\epsilon s_i) / \sum_j \exp(\epsilon s_j)$

This formula ensures two things

1. probabilities are always positive (since exponentials are never negative).
2. Amplifies difference in scores when  $\epsilon$  is large

### Owner Utility

If the user is chosen their utility is:

- Utility of chosen owner = (valuation of user i) - (payment to user i)
- $U_i = v_i - \text{payment}_i$

If the user is not chosen, their utility is 0.

### Collector Payoff

Welfare of the collector is calculated by:

- $\sum_i \Pr(i \text{ chosen})(v_i - c_i)$
- $\Pr$  = probability
- Higher  $\epsilon$ , more accurate welfare. Lower  $\epsilon$ , more privacy

## 5. Equilibrium Concept

### Mechanism Design Optimum (Dominant-Strategy Truthfulness)

- The collector implements a VCG-style payment rule, ensuring that the owner's best response is truth-telling, regardless of other players' reports.

- Differential privacy introduces noise but does not break truthfulness; expected utilities preserve incentive compatibility.

Why this fits:

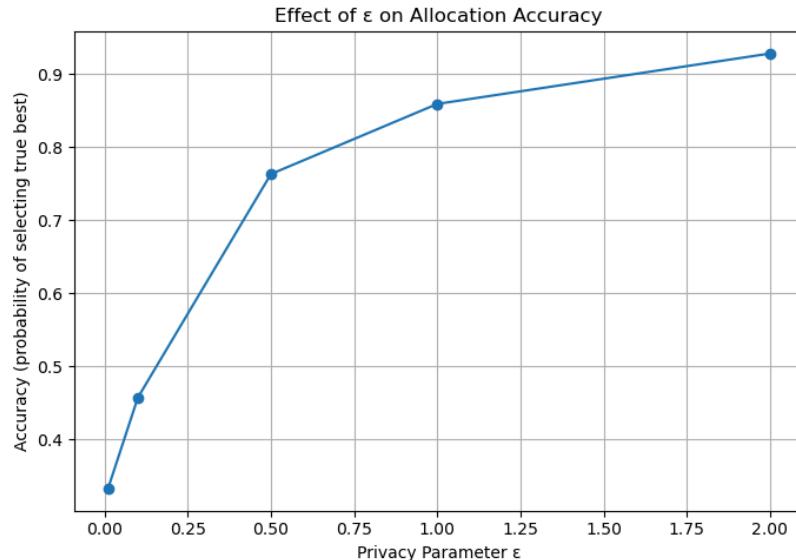
- OCG models are typically evaluated using truthfulness as the equilibrium. The owner optimally reveals truthful valuations, and the collector optimally chooses  $\epsilon$  to balance efficiency vs. privacy.

## 6. Implementation Details

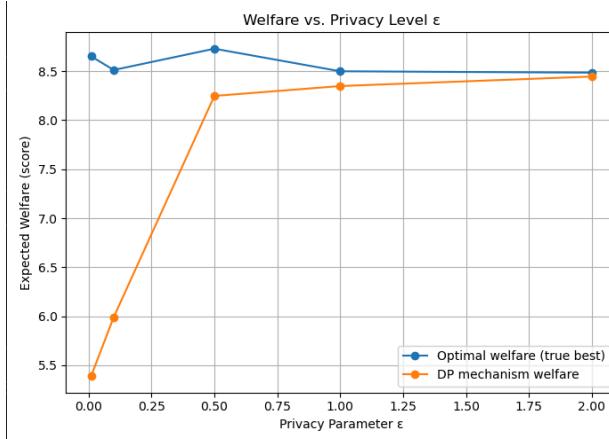
### Computational Approach

- Generate a set of passengers with valuations and costs.
- Compute each passenger's score  $s_i$ .
- For a chosen  $\epsilon$ , compute exponential-mechanism allocation probabilities.
- Randomly draw a selected passenger based on these probabilities.
- Compute VCG-style payments.
- Repeat over many runs to estimate expected utilities and welfare under different  $\epsilon$  values.

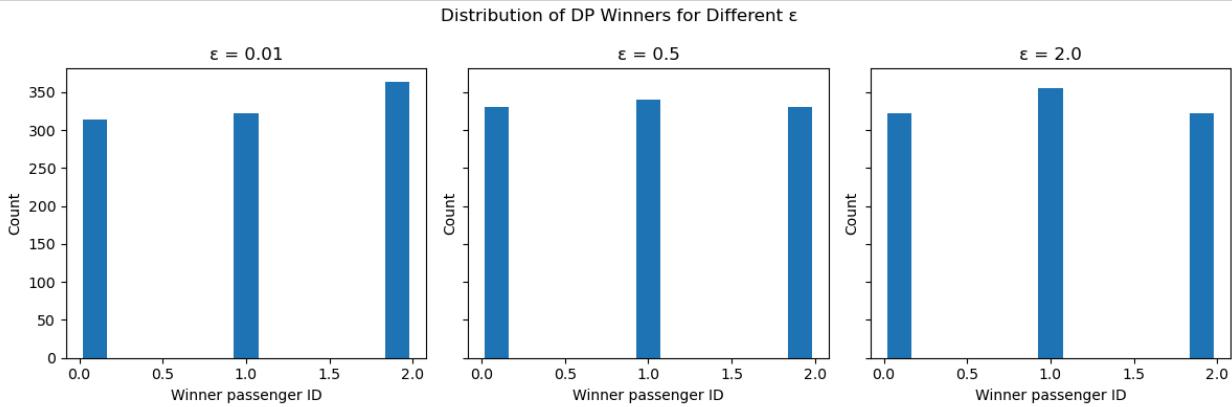
## 7. Results



- The probability of selecting the highest-score passenger increases sharply with  $\epsilon$ .
- For  $\epsilon = 0$ , allocation becomes almost uniform.



- Stronger privacy (small  $\epsilon$ ) hurts system performance, because choices become nearly random.
- Increasing  $\epsilon$  improves welfare, because the mechanism increasingly favors high-score passengers.
- At sufficiently large  $\epsilon$ , DP noise becomes negligible, so welfare converges to the optimal.



- Differential privacy randomizes outcomes, and low  $\epsilon$  makes the mechanism nearly uniform.
- As  $\epsilon$  increases, the distribution becomes more peaked toward the true best passenger.
- The figure visually captures the privacy–accuracy tradeoff, showing how noise drops and preference emerges as  $\epsilon$  grows.

## Observations

- Higher  $\epsilon$  = weaker privacy, higher welfare, more accuracy.
- Lower  $\epsilon$  = stronger privacy, lower welfare.
- The tradeoff is smooth and matches the pattern described in Wu et al. (2020): owners benefit from privacy protection but face increased randomness.

## 8. Interpretation

This OCG simulation shows how a platform can use differential privacy to protect user location data while still extracting truthful inputs through a VCG-based mechanism. Owners are incentivized to report truthfully, but their probability of being chosen depends on the privacy parameter  $\epsilon$ .

As Wu et al. emphasize, the key insight is the efficiency–privacy tradeoff:

- Stronger privacy requires more randomization (low  $\epsilon$ ),
- But randomization reduces allocation accuracy and welfare.

The collector must therefore choose  $\epsilon$  to balance user trust and system performance. Our results match the general trend found in the literature.