

# Two-Echelon Multicommodity Location Model

## Sets

- $I$  = set of production plants; (*index*  $i$ )
- $J$  = set of potential Distribution Centers (DCs); (*index*  $j$ )
- $R$  = set of demand nodes; (*index*  $r$ )
- $K$  = set of homogeneous commodities; (*index*  $k$ )

## Parameters

- $p$  = maximum number of DCs that can be opened;
- $c_{ijr}^k$  = unit transportation cost of commodity  $k \in K$  from plant node  $i \in I$  to demand node  $r \in R$  across DC  $j \in J$ ;
- $d_r^k$  = demand of commodity  $k \in K$  from demand node  $r \in R$ ;
- $p_i^k$  = maximum quantity of commodity  $k \in K$  that can be manufactured by plant  $i \in I$ ;
- $q_j^-$  = minimum activity level of potential DC  $j \in J$ ;
- $q_j^+$  = maximum activity level of potential DC  $j \in J$ ;
- $f_j$  = fixed cost of potential DC  $j \in J$ ;
- $g_j$  = marginal cost of potential DC  $j \in J$ ;

## Variables

$$z_j = \begin{cases} 1 & \text{if DC } j \in J \text{ is opened} \\ 0 & \text{otherwise} \end{cases}$$

$$y_{jr} = \begin{cases} 1 & \text{if demand node } r \in R \text{ is assigned to DC } j \in J \\ 0 & \text{otherwise} \end{cases}$$

- $s_{ijr}^k$  = amount of commodity  $k \in K$  transported from plant node  $i \in I$  to demand node  $r \in R$  across DC  $j \in J$ ;

The following feasibility condition must hold:

$$\sum_{i \in I} p_i^k \geq \sum_{r \in R} d_r^k \quad k \in K$$

## TEMC Mathematical Formulation

$$\sum_{i \in I} \sum_{j \in J} \sum_{r \in R} \sum_{k \in K} c_{ijr}^k \cdot s_{ijr}^k + \sum_{j \in J} \left( f_j \cdot z_j + g_j \cdot \sum_{r \in R} \sum_{k \in K} d_r^k \cdot y_{jr} \right)$$

$$\sum_{j \in J} \sum_{r \in R} s_{ijr}^k \leq p_i^k \quad i \in I, k \in K \quad (1)$$

$$\sum_{i \in I} s_{ijr}^k = d_r^k \cdot y_{jr} \quad j \in J, r \in R, k \in K \quad (2)$$

$$\sum_{j \in J} y_{jr} = 1 \quad r \in R \quad (3)$$

$$q_j^- \cdot z_j \leq \sum_{r \in R} \sum_{k \in K} d_r^k \cdot y_{jr} \leq q_j^+ \cdot z_j \quad j \in J \quad (4)$$

$$\sum_{j \in J} z_j = p \quad (5)$$

$$z_j \in \{0, 1\} \quad j \in J \quad (6)$$

$$y_{jr} \in \{0, 1\} \quad j \in J, r \in R \quad (7)$$

$$s_{ijr}^k \geq 0 \quad i \in I, j \in J, r \in R, k \in K \quad (8)$$

## Demand Allocation Problem

If a set  $\bar{z}$   $j \in J$  and  $\bar{y}_{jr}$ ,  $j \in J, r \in R$  of feasible values is available, you just have to solve the following LP problem in order to determine the optimal demand allocation:

$$\sum_{i \in I} \sum_{j \in J} \sum_{r \in R} \sum_{k \in K} c_{ijr}^k \cdot s_{ijr}^k$$

$$\sum_{j \in J} \sum_{r \in R} s_{ijr}^k \leq p_i^k \quad i \in I, k \in K$$

$$\sum_{i \in I} s_{ijr}^k = d_r^k \cdot \bar{y}_{jr} \quad j \in J, r \in R, k \in K$$

$$s_{ijr}^k \geq 0 \quad i \in I, j \in J, r \in R, k \in K$$