



CSC258: Computer Organization

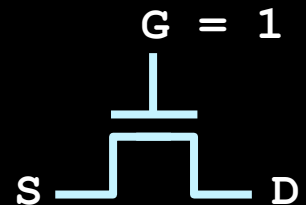


Circuit Creation

*Based on slides originally created by Prof. Steve Engels

Review - Transistors

- Transistors are switches whose behavior can be controlled electrically
- nMOS conducts electricity between source and drain when positive voltage (5V) is applied to the gate
- pMOS conducts electricity between source and drain when ground (0 V) is applied to the gate

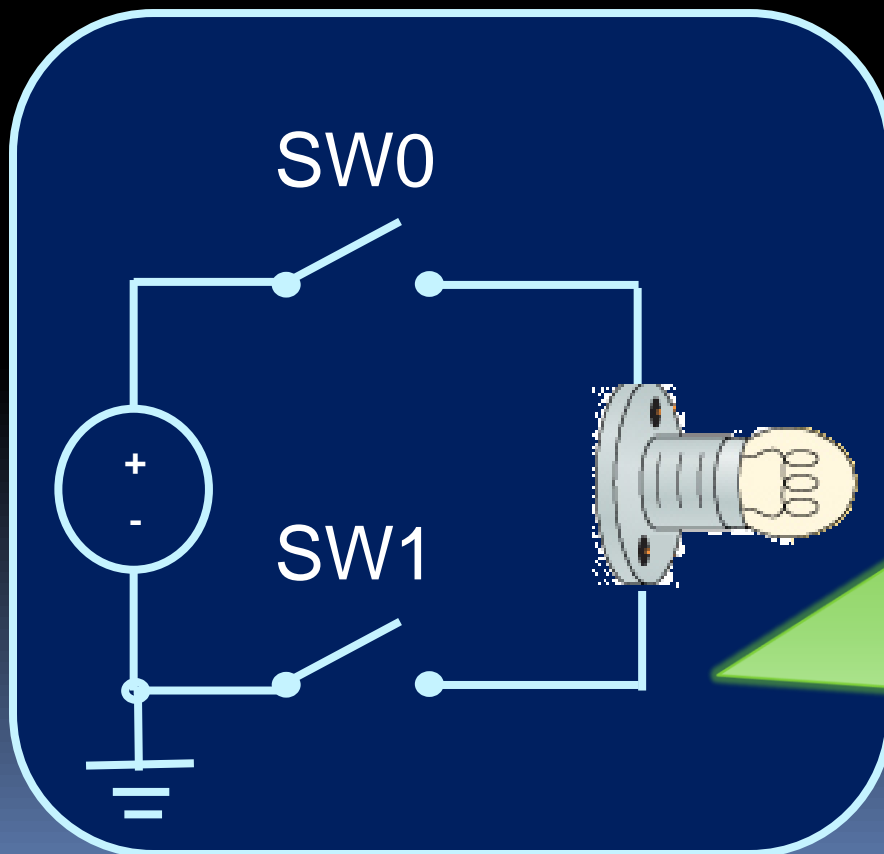


Review – Logic Values

- Logic values are about potential
 - When a point in circuit is at 5V potential, we call that high value or 1
 - The point is at 5V potential, if there is a direct path (no resistance) from **V_{cc}** to this point in the circuit
 - When a point in circuit is at 0V potential, we call that **low** value or 0
 - The point is at 0V potential, if there is a direct path (no resistance) from **GND** to this point in the circuit
- A point in the circuit which does not have a path to either VCC or GND is at an **undefined** potential

Review: Disconnected $\neq 0$

- Easy to demonstrate that disconnected points are **not** the same as 0

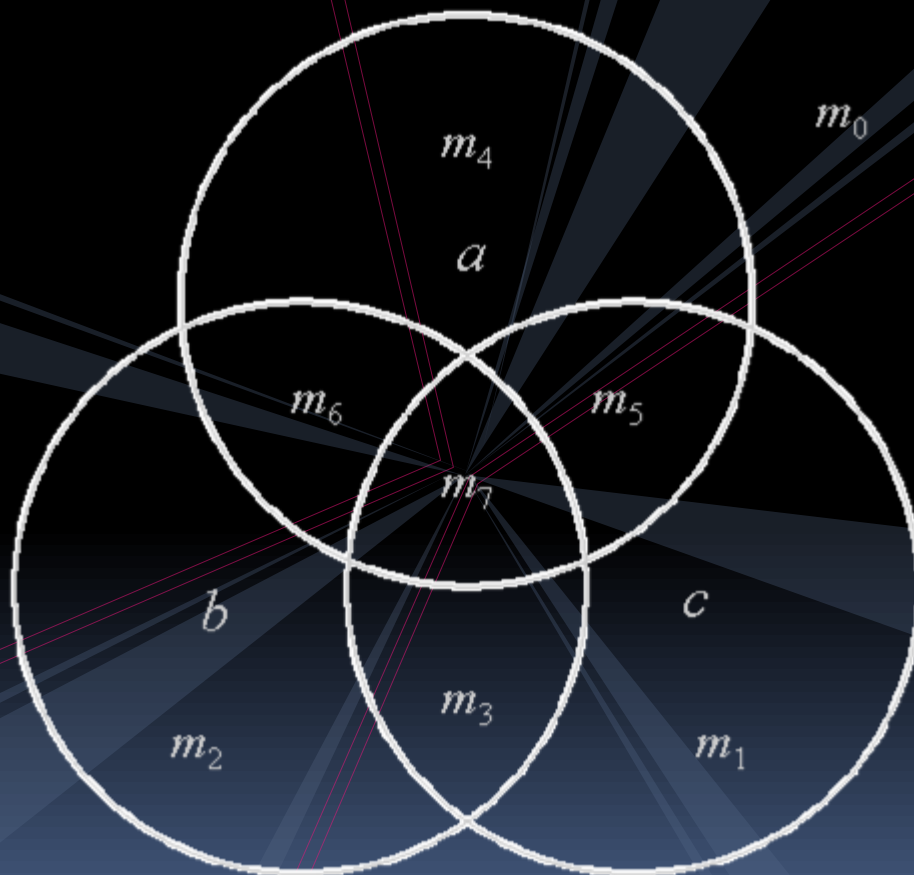


This point in the circuit is disconnected from both V_{cc} or GND.

If disconnected was the same as 0, then closing switch SW0 would cause the light to come on, because there would be a difference in potential between the light's two contacts!



Minterms



Example truth table

- Consider the following example:
 - *“Given three inputs A , B , and C , make output Y high whenever any of the inputs are low, except when all three are low or when A and C are high.”*
- This leads to the truth table on the right.
 - Is there a better way to describe the cases when the circuit's output is high?

A	B	C	Y
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0

Minterms

- An easier way to express logic function is to assume a standard truth table format, then list which rows cause high output.
- Each row of the truth table represents one **minterm**.

A	B	C	Y
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0



Minterm	Y
m_0	0
m_1	1
m_2	1
m_3	1
m_4	1
m_5	0
m_6	1
m_7	0

Formally speaking

- A more formal description:
 - **Minterm** = an AND expression with every input present in true or complemented form.
 - For example, given four inputs (A, B, C, D):

- Valid minterms:

$$A \cdot \bar{B} \cdot C \cdot D$$

$$\bar{A} \cdot B \cdot \bar{C} \cdot D$$

$$A \cdot B \cdot C \cdot D$$

- Not minterms:

$$A \cdot B + C \cdot D$$

$$A + B + C + \bar{D}$$

$$A \cdot B \cdot D$$

Numbering minterms

- Circuits are often described using minterms, as a form of logic shorthand.
 - Given n inputs, there are 2^n minterms possible (same as rows in a truth table).
 - Naming scheme:
 - Minterms are labeled as m_x
 - The x subscript indicates the row in the truth table.
 - x starts at 0 (when all inputs are low), and ends with $n-1$.
 - x corresponds to the binary number formed by the variables in the truth table, i.e. $000 \rightarrow m_0, 001 \rightarrow m_1, \dots, 111 \rightarrow m_7$
 - Example: Given 3 inputs
 - Minterms are $m_0 (\overline{A} \cdot \overline{B} \cdot \overline{C})$ to $m_7 (A \cdot B \cdot C)$

Using minterms

- What are minterms used for?
 - A single minterm indicates a set of inputs that will make the output go high.
 - Example: m_2
 - Output only goes high in third line of truth table.

A	B	C	D	m_2
0	0	0	0	0
0	0	0	1	0
0	0	1	0	1
0	0	1	1	0
0	1	0	0	0
0	1	0	1	0
0	1	1	0	0
0	1	1	1	0
1	0	0	0	0
1	0	0	1	0
1	0	1	0	0
1	0	1	1	0
1	1	0	0	0
1	1	0	1	0
1	1	1	0	0
1	1	1	1	0

Using minterms

- What happens when you OR two minterms?
 - Result is output that goes _____
 - For $m_2 + m_8$, both third and ninth lines of truth table result in high output.

A	B	C	D	m_2	m_8	$m_2 + m_8$
0	0	0	0	0	0	
0	0	0	1	0	0	
0	0	1	0	1	0	
0	0	1	1	0	0	
0	1	0	0	0	0	
0	1	0	1	0	0	
0	1	1	0	0	0	
0	1	1	1	0	0	
1	0	0	0	0	1	
1	0	0	1	0	0	
1	0	1	0	0	0	
1	0	1	1	0	0	
1	1	0	0	0	0	
1	1	0	1	0	0	
1	1	1	0	0	0	
1	1	1	1	0	0	

Creating Boolean expressions

- **Sum-of-Minterms (SOM):**
 - Since each minterm corresponds to a single high output in the truth table, the combined high outputs are a **union** of these minterm expressions.
 - Also known as _____.

$$Y = m_2 + m_6 + m_7 + m_{10} \quad (\text{SOM})$$

A	B	C	D	m_2	m_6	m_7	m_{10}	Y
0	0	0	0					
0	0	0	1					
0	0	1	0					
0	0	1	1					
0	1	0	0					
0	1	0	1					
0	1	1	0					
0	1	1	1					
1	0	0	0					
1	0	0	1					
1	0	1	0					
1	0	1	1					
1	1	0	0					
1	1	0	1					
1	1	1	0					
1	1	1	1					

Maxterms

- An alternative to minterms
 - Express logic function by specifying which rows cause **low** output.
 - Numbering of maxterms is the same as numbering of minterms (with a caveat)

A	B	C	Y
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0



Maxterm	Y
M₀	0
M ₁	1
M ₂	1
M ₃	1
M ₄	1
M₅	0
M ₆	1
M₇	0

Formally speaking

- A more formal description:
 - **Maxterm** = an OR expression with every input present in true or complemented form.
 - For example, given four inputs (A, B, C, D):
 - Valid maxterms:

$$A + \overline{B} + C + D$$

$$\overline{A} + B + \overline{C} + D$$

$$A + B + C + D$$

- Not maxterms:

$$A \cdot B + C \cdot D$$

$$A \cdot B \cdot C \cdot \overline{D}$$

$$A + B + D$$

Numbering maxterms

- Similar to minterms
 - Given n inputs, there are 2^n maxterms possible (same as rows in a truth table).
 - Maxterms are labeled as M_x
 - The x subscript indicates the row in the truth table.
 - x starts at 0 (when all inputs are low), and ends with $n-1$.
 - x corresponds to the binary number formed by the variables in the truth table, i.e. $000 \rightarrow M_0$, $001 \rightarrow M_1, \dots, 111 \rightarrow M_7$
 - Example: Given 3 inputs
 - Maxterms are $M_0 (A+B+C)$ to $M_7 (\bar{A}+\bar{B}+\bar{C})$

Using maxterms

- A single maxterm indicates a set of inputs that will make the output go low.
 - Example: M_2
 - Output only goes low in the third line of truth table.

□ $M_2 = A + B + \overline{C} + D$

A	B	C	D	M_2
0	0	0	0	1
0	0	0	1	1
0	0	1	0	0
0	0	1	1	1
0	1	0	0	1
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	1
1	0	0	1	1
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

Using maxterms

- What happens when you **AND** two maxterms?
 - Result is output that goes **low** in both maxterm cases.
 - For $M_2 \cdot M_8$, both third and ninth lines of truth table result in low output.

A	B	C	D	M_2	M_8	$M_2 \cdot M_8$
0	0	0	0	1	1	1
0	0	0	1	1	1	1
0	0	1	0	0	1	0
0	0	1	1	1	1	1
0	1	0	0	1	1	1
0	1	0	1	1	1	1
0	1	1	0	1	1	1
0	1	1	1	1	1	1
1	0	0	0	1	0	0
1	0	0	1	1	1	1
1	0	1	0	1	1	1
1	0	1	1	1	1	1
1	1	0	0	1	1	1
1	1	0	1	1	1	1
1	1	1	0	1	1	1
1	1	1	1	1	1	1

$$Y = M_3 \cdot M_5 \cdot M_7 \cdot M_{10} \cdot M_{14} \text{ (POM)}$$

A	B	C	D	M ₃	M ₅	M ₇	M ₁₀	M ₁₄	Y
0	0	0	0						
0	0	0	1						
0	0	1	0						
0	0	1	1						
0	1	0	0						
0	1	0	1						
0	1	1	0						
0	1	1	1						
1	0	0	0						
1	0	0	1						
1	0	1	0						
1	0	1	1						
1	1	0	0						
1	1	0	1						
1	1	1	0						
1	1	1	1						

Creating Boolean expressions

- Two canonical forms of Boolean expressions:
 - **Sum-of-Minterms (SOM):**
 - Since each minterm corresponds to a single high output in the truth table, the combined high outputs are a _____ of these minterm expressions.
 - Also known as: Sum-of-Products.
 - **Product-of-Maxterms (POM):**
 - Since each maxterm only produces a single low output in the truth table, the combined low outputs are an _____ of these maxterm expressions.
 - Also known as Product-of-Sums.

Using Minterms and Maxterms

- Using minterms or maxterms, logic functions can be expressed in a more compact way than showing entire truth tables
- Only works if ordering of columns in the truth table is agreed upon
- Sum-of-minterms are useful in cases _____

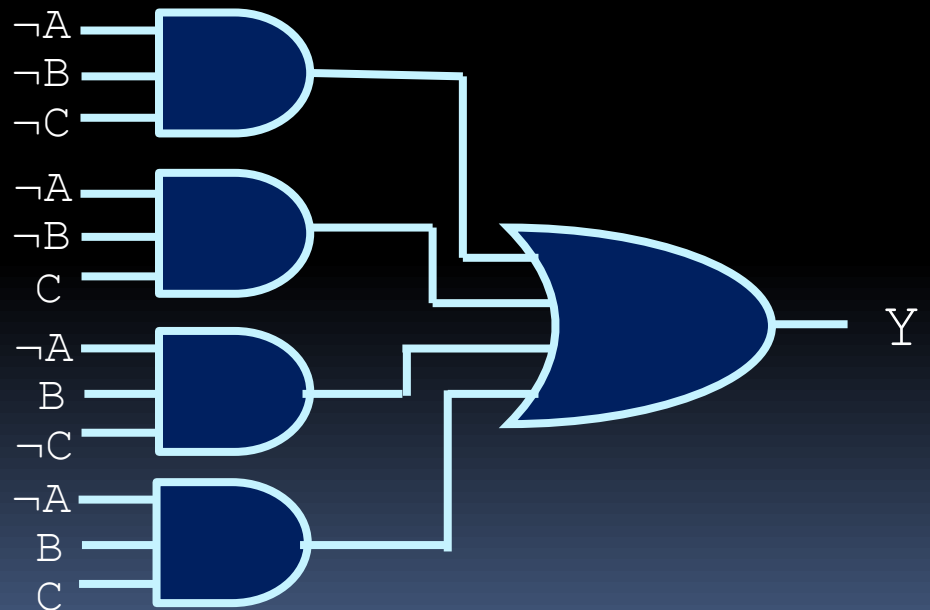
- Product-of-maxterms useful when expressing truth tables that have _____

Converting SOM to gates

- Once you have a Sum-of-Minterms expression, it is easy to convert this to the equivalent combination of gates:

$$m_0 + m_1 + m_2 + m_3 =$$

$$\bar{A} \cdot \bar{B} \cdot \bar{C} + \bar{A} \cdot \bar{B} \cdot C + \bar{A} \cdot B \cdot \bar{C} + \bar{A} \cdot B \cdot C =$$



Quick Exercises

- Given 4 inputs A, B, C and D write:

- $m_9 \Rightarrow$
- $m_{15} \Rightarrow$
- $m_{16} \Rightarrow$
- $M_2 \Rightarrow$

- Which minterm is this one?

- $A' \cdot B \cdot C' \cdot D'$

- Which maxterm is this one?

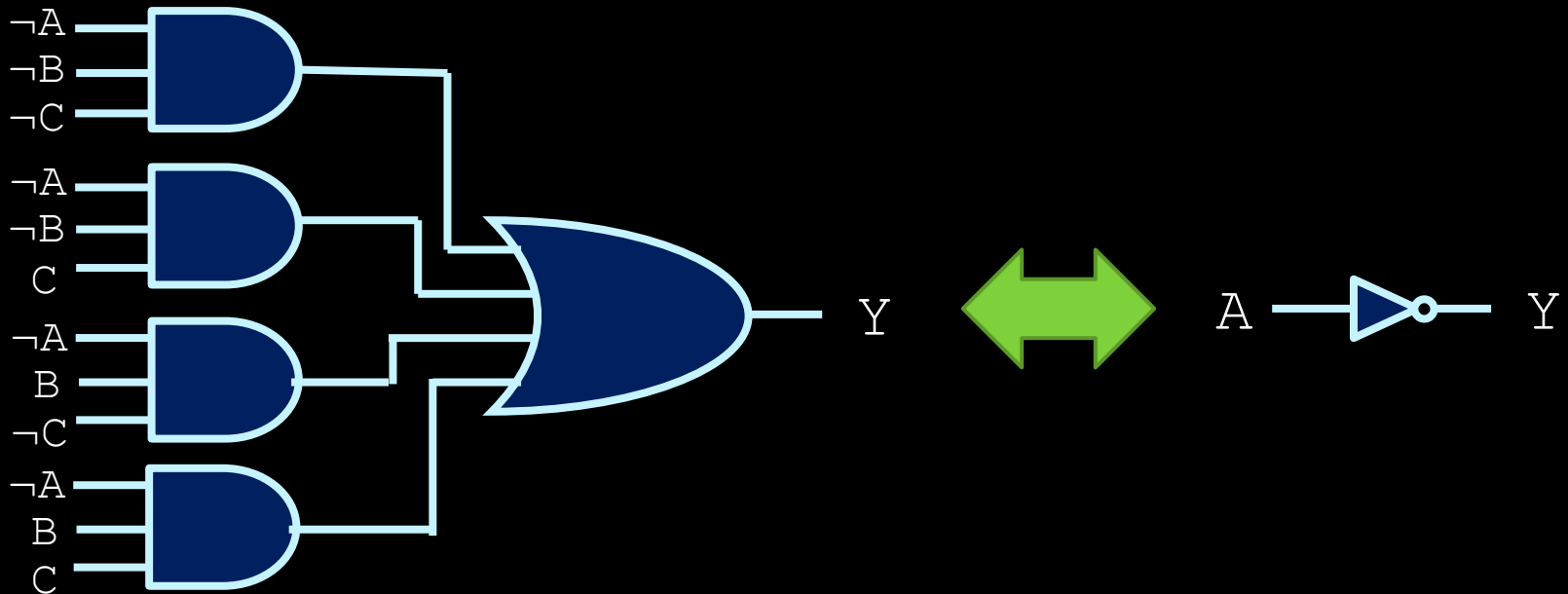
- $A + B + C + D'$



Reducing circuits



Reasons for reducing circuits



- Note example of Sum-of-Minterms circuit design.
- To minimize the number of gates, we want to reduce the Boolean expression as much as possible from a collection of minterms to something smaller.
- This is where CSC165 skills come in handy 😊

Boolean algebra review

- Axioms:

$$0 \cdot 0 = 0$$

$$1 \cdot 1 = 1$$

$$0 \cdot 1 = 1 \cdot 0 = 0$$

$$\text{if } x = 1, \overline{x} = 0$$

- From this, we can extrapolate:

If one input of a 2-input AND gate is 1, then the output is whatever value the other input is.

$$x \cdot 0 =$$

$$x \cdot 1 =$$

$$x \cdot x =$$

$$x \cdot \overline{x} =$$

$$\overline{\overline{x}} =$$

$$x + 1 =$$

$$x + 0 =$$

$$x + x =$$

$$x + \overline{x} =$$

If one input of a 2-input OR gate is 0, then the output is whatever value the other input is.

Other Boolean identities

- Commutative Law:

$$x \cdot y = y \cdot x \qquad x + y = y + x$$

- Associative Law:

$$\begin{aligned} x \cdot (y \cdot z) &= (x \cdot y) \cdot z \\ x + (y + z) &= (x + y) + z \end{aligned}$$

- Distributive Law:

$$\begin{aligned} x \cdot (y + z) &= x \cdot y + x \cdot z \\ x + (y \cdot z) &= (x + y) \cdot (x + z) \end{aligned}$$

This part is non-intuitive but true

Other Boolean identities

- Simplification Law:

$$x + (\bar{x} \cdot y) = x + y \qquad x \cdot (\bar{x} + y) = x \cdot y$$

- Consensus Law:

$$x \cdot y + \bar{x} \cdot z + y \cdot z = x \cdot y + \bar{x} \cdot z$$

- Proof by Venn diagram:



Consensus Law Proof -Venn diagram

- Consensus Law:

$$x \cdot y + \bar{x} \cdot z + y \cdot z = x \cdot y + \bar{x} \cdot z$$

- Proof by Venn diagram:

- $x \cdot y$
- $\bar{x} \cdot z$
- $y \cdot z$
 - Already covered!



Proving Boolean identities

- The most straightforward way to prove a Boolean identity is to write the truth table for left and right side, and show that they produce the same function
- Exercise:
 - Prove the “non-intuitive” distributive law
 - $x + (y \cdot z) = (x + y) \cdot (x + z)$

Other Boolean identities

- Absorption Law:

$$x \cdot (x + y) = x$$

$$x + (x \cdot y) = x$$

- De Morgan's Laws:

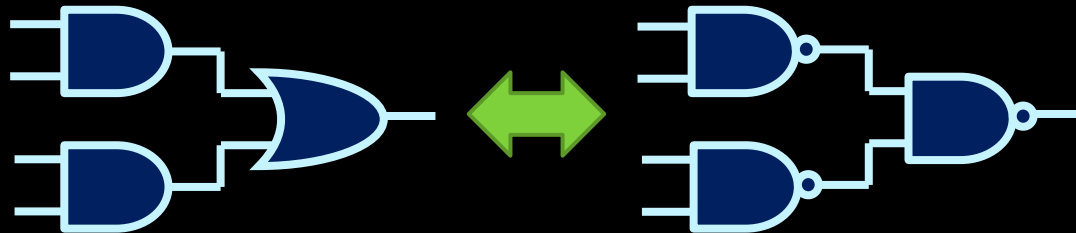
$$\overline{x} \cdot \overline{y} = \overline{x + y}$$

$$\overline{x + y} = \overline{x} \cdot \overline{y}$$

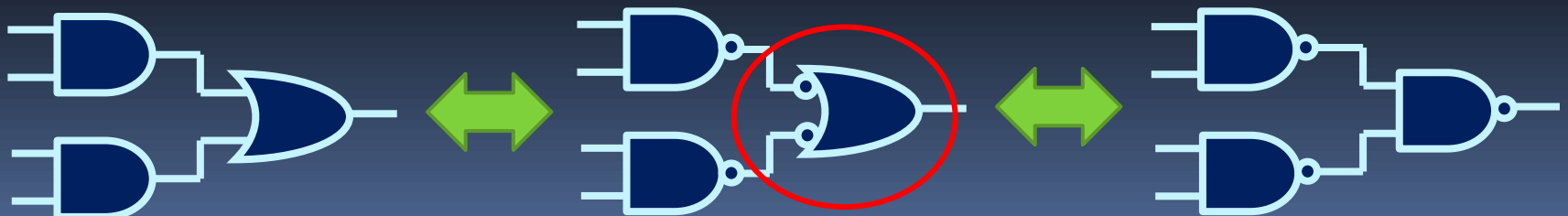


Converting to NAND gates

- De Morgan's Law is important because out of all two-input gates, NANDs are the cheapest to fabricate.
 - ▣ A Sum-of-Products circuit could be converted into an equivalent circuit of NAND gates:



- This is all based on de Morgan's Law:



Reducing Boolean expressions

A	B	C	Y
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

- Assuming function on the left, we get the following:

$$Y = \bar{A} \cdot B \cdot C + A \cdot \bar{B} \cdot \bar{C} + A \cdot B \cdot \bar{C} + A \cdot B \cdot C$$

- Now start combining terms, like the last two:

$$Y = \bar{A} \cdot B \cdot C + A \cdot \bar{B} \cdot \bar{C} + \mathbf{A \cdot B}$$

Reducing Boolean expressions

- Different final expressions possible, depending on what terms you combine.
- For instance, given the previous example:

$$Y = \bar{A} \cdot B \cdot C + A \cdot \bar{B} \cdot \bar{C} + A \cdot B \cdot \bar{C} + A \cdot B \cdot C$$

- If you combine the end and middle terms...

$$Y = B \cdot C + A \cdot \bar{C}$$

- Which reduces the number of gates and inputs!